3-1)

R_0 $interp_r : R_0 \Rightarrow ans$

$opt : R_0 \Rightarrow R_0$

X_0 $interp_x : X_0 \Rightarrow ans$

R_1 ..

..

compile : $R_1 \Rightarrow X_0$

R_1 : tree-shaped & recursive, expr-oriented structure

(+ 1 (+ 2 (- (+ 3 4))))

infinite variable (let x := 2 in ....)

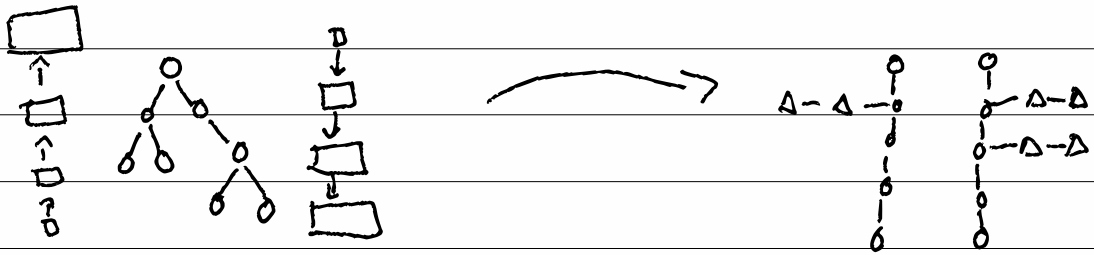
variables are scoped (+ (let x ...) (let x ...))

3-2)

X_0 : linear, heavily structured

fixed #/set of registers

registers/mem are global

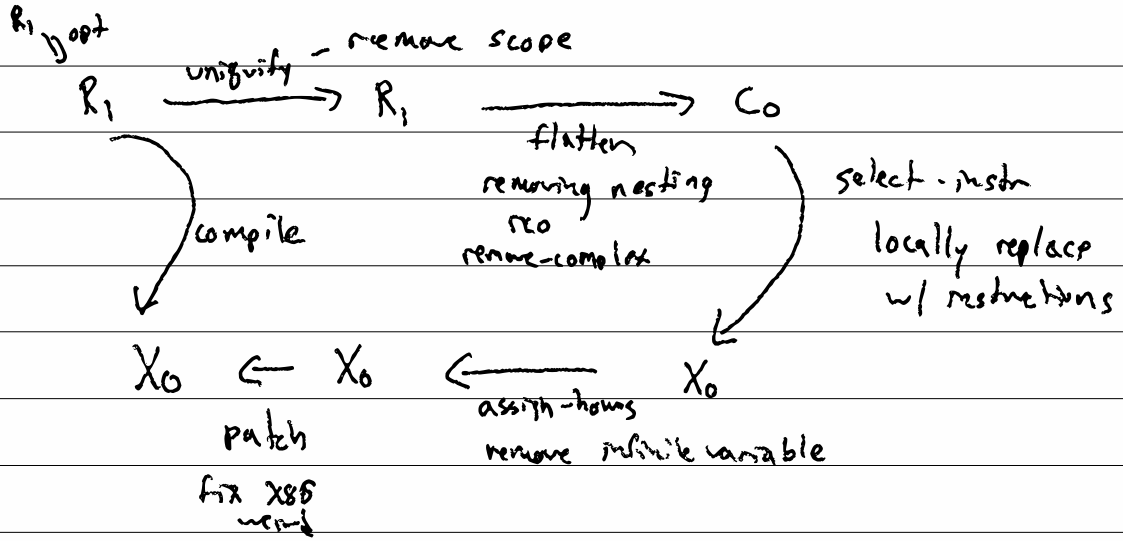


locally: ○ ⇝ ○-△-△

top-down fixed-point/cyclic

bottom up

3-3/ compilation w/ intermediate languages



3-4/

C0 p ::= (program info [label => tail] ...)

tail ::= (return arg) | (seq stmt tail)

stmt ::= (set! var expr)

expr ::= arg | (read) | (- arg) | (+ arg arg)

arg ::= num | var

cip : p ⇒ ang cip (program _ l=>tm) =

cit ∅ (l=>tm ("main"))

cit env (return arg) = cia env arg

(seq stmt t) = cit (cig env stmt) t

3-5

cis env (set! v e) = env[v ↦ cie env e]

cie env a = cia env a

(read) = read from the user

(Neg a) = -1 * cia env a

(Add l r) = cia env l + cia env r

cia env (Num n) = n

(Var v) = env[v]

3-b)

uniquify : $R_1 \Rightarrow R_1$

job : remove scopedness from R

```
(+ (let x := 7 in x)          (+ (let v0 := 7 in v0)
   (let x = 8 in      ⇝          (let v1 = 8 in
   (let x = 1 + x in               (let v2 = 1 + v1 in
      (+ x x))))                      (+ v2 v2))))
```

uni : $(\text{var} \Rightarrow \text{var})^{\sigma} \Rightarrow e \Rightarrow e$

uni σ (Var v) = Var (σ v)

uni σ (Let x xe be) = Let x' (uni σ xe) (uni σ' be)

where x' = is the new variable $\sigma' = \sigma[x \mapsto x']$

3-7

```
uni e   (Add l r) = Add (uni e l) (uni e r)
        (Neg e)   = Neg (uni e e)
        (Num n)   = Num n

for i = 0 to 1024
    p = randp 6
    check (interp p)    (interp (uni p))
    check (interp p)    (interp (opt p))
check-all-eq p , (uni p) , (opt (uni p))
```

3-8 (assume that all vars are unique)

rco : $R_1 \Rightarrow R_0^*$ flatten = econ • rco

p : (program info e)

e ::= arg | (let x = c in e)

c ::= (read) | (- arg) | (+ arg arg)

arg ::= num | var

in : (+ (+ 2 3) (let x = read in (+ x x)))

out : (let v_0 = (+ 2 3) in
(let x = (read) in
(let v_1 = (+ x x) in
(let v_2 = (+ v_0 v_1) in v_2))))

3-9/ rco : (xvar ⇒ e) x e ⇒ (first (xvar x e)) x e

rco (program ; e) = let $x_0 = c_0$ in ... let $x_n = c_n$ in ans

$[x_0, c_0] \ldots [x_n, c_n]$, ans) = rco ∅ e

rco σ (Num n) = (∅, (Num n))

(var x) = (∅, σ(x))

(Add e_L e_R) = (nv_L ++ nv_R ++ [(x, Add a_L a_R)], x)

where (nv_L, a_L) = rco σ e_L

(nv_R, a_R) = rco σ e_R

x is a fresh variable

rco σ (Neg e_L) = (nv_L ++ [(x, (Neg a_L))], x)

where (nv_L, a_L) = rco σ e_L

rco σ (Read) = ([(x, (Read))], x)

rco σ (Let x xe be) = (nv_x ++ nv_b, a_b)

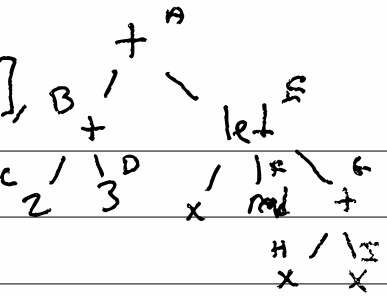
where (nv_x, a_x) = rco σ xe

(nv_b, a_b) = rco σ' be

σ' = σ [x ↦ a_x]

3-10

rco	$\emptyset$	$A = [(v_0, (+\ 2\ 3)), v_1, v_2, (v_3, (+\ v_0\ v_2))]$, v_2	
rco	$\emptyset$	$B = [(v_0, (+\ 2\ 3))]$, v_0	
rco	$\emptyset$	$C = []$, 2	
rco	$\emptyset$	$D = []$, 3	
rco	$\emptyset$	$E = [(v_1, (\text{rand})), (v_2, (+\ v_1\ v_1))]$, v_2	let $v_0 = (+\ 2\ 3)$
rco	$\emptyset$	$F = [(v_1, (\text{rand}))]$, v_1	let $v_1 = (\text{rand})$
rco	$[x \mapsto v_1]$	$G = [(v_2, (+\ v_1\ v_1))]$, v_2	let $v_2 = (+\ v_1\ v_1)$
rco	"	$H = []$, v_1	let $v_3 = (+\ v_0\ v_2)$
rco	"	$I = []$, v_1	v_3



3-III → explicate-control

econ : R_1 (rco-stk) → C_0

econ (Program : e) = (Program : [main → t]
  where t = econ_e e

econ_e arg = (return arg)

econ_e (let x = xe in be) = (seq (get! x xe)
  (econ_e be))