

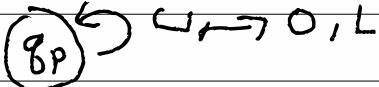
20-1 / $(Q, \Sigma, \Gamma, q_0, \delta: Q \times \Gamma \rightarrow Q \times \Gamma \times \{L, R\}, g_p)$

"machines" "languages / expression"

DFA _s	\Leftrightarrow	REX _s
PDA	\Leftrightarrow	CFG _s
TM		enumerator

$e = (Q, \Sigma, \Gamma, q_0, \delta: Q \times \Gamma \rightarrow Q \times \Gamma \times \{L, R\}, g_p)$

$w \in L(e)$ if $\epsilon[q_0] \epsilon \Rightarrow^* x[g_p] w$

\Downarrow  $\hookrightarrow O, L$ O^*

20-2/ Suppose I had a Turing decider and I wanted an enumerator...

enumerator := try all strings in lexicographic
and if the decider says yes,
then it prints

recognizer \Rightarrow enumerator := try all strings "in order"

enumerator \Rightarrow recognizer := check the output...
never say no

enumerator "in order" \Rightarrow decider := check the output...

say no, if somehow hyperlinked

20-3/ DFAs

↑

↔

REX

NFAs (more expressive)

TMs

↔

Enumerations

↓

NTM \leftrightarrow v, n

↙

NTM \leftrightarrow o

Zo-y) want: an extension of Turing machines

$\forall a \in \text{input}, \exists b \text{ output}, M(a) = M(b)$

$$a \xrightarrow{\text{compute}} b$$

$$\downarrow m \quad \swarrow m$$

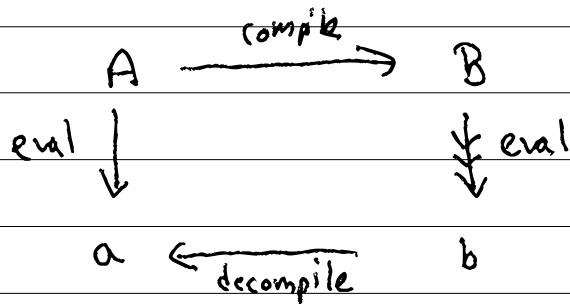
x

Galois connection on a bisimulation

70-5/

If a is input, $\exists b$ be output.

$\text{eval}_i(a)$ similar $\text{eval}_o(b)$



20-b/ TMs w/ a "stay" action

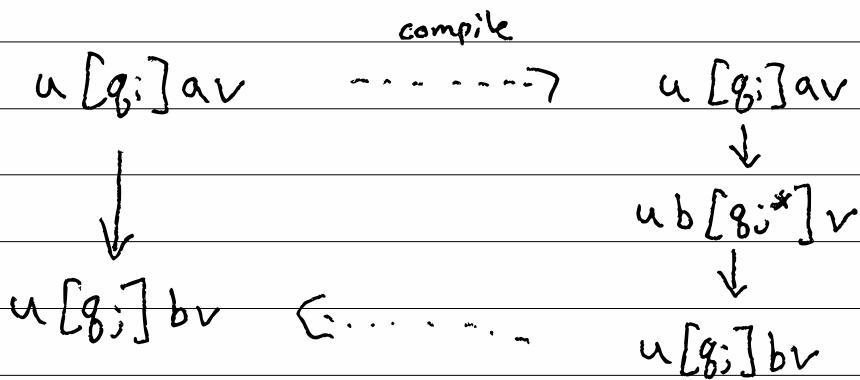
$$\delta: Q \times \Gamma \rightarrow Q \times \Gamma \times \{L, R\}$$

$$\delta_s: Q \times \Gamma \rightarrow Q \times \Gamma \times \{L, R, S\}$$

$$\underline{\delta(q_i, a) = (q_j, b, s)}$$

$$u[q_i]av \Rightarrow u[q_j]bv$$

20-7 / stay = input \longrightarrow normal Tm



$$\begin{aligned} S(g_i, a) &= (g_i, b, S) \\ \Rightarrow S(g_i, a) &= (g_i^*, b, R) \end{aligned}$$

$$\forall g_i \in Q, \quad \forall \gamma \in \Gamma. \quad S(g_i^*, \gamma) = (g_i, \gamma, L)$$