

2-1 \emptyset $\{x\}$ $A \cup B$ $A \cap B$ A^c $P(A)$

Alphabet Σ

$\{0, 1\}$

Universe - strings of Σ

Σ^*

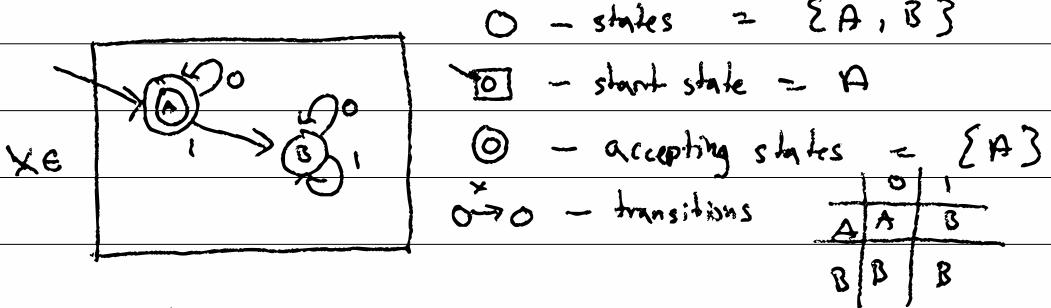
$$\text{ALL} = P(\Sigma^*)$$

- All strings made of 0s
- All strings that are QIF
- All Harry novels
- All algebra problems "x + 3 = 7"
- All the math formulae
"1+1=2" "forall x, x+0=x"

FIN
 $\{\text{fm, h}\}$ let's

2-2 / $\{ \text{pen, table} \}$ $\{\checkmark, \square\}$
 $= \{\text{pen}\} \cup \{\text{table}\}$

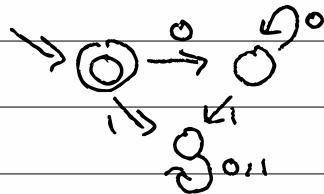
Only Zeros: DFA - Deterministic Finite Automata



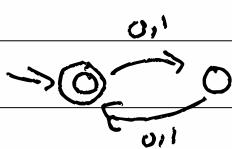
$\epsilon? \checkmark$ $0? \checkmark$ $10? X$ $000? \checkmark$

2-3/

only zeros



easily-long



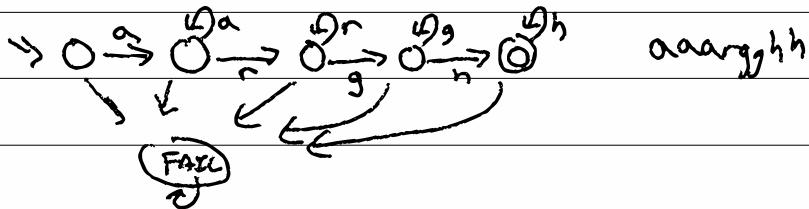
ε? ✓

0? X

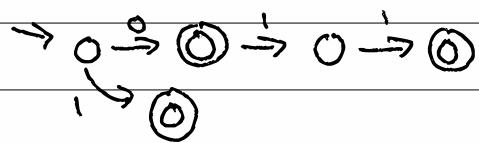
0!? ✓

10 ? ✓

Arg h



2-4) {0, 1, 0113}



2-5) $x \in \text{DFA} = (\mathbb{Q}, \Sigma, q_0, \delta, F)$

states' / alphabet $\Sigma \subseteq \mathbb{Q}$ | transitions $\delta: Q \times \Sigma \rightarrow Q$ | $F \subseteq Q$
 starting accepting states

DFA configuration = $\mathbb{Q} \times \Sigma^*$ = $[q]w$

state string

$[A]_{011}$

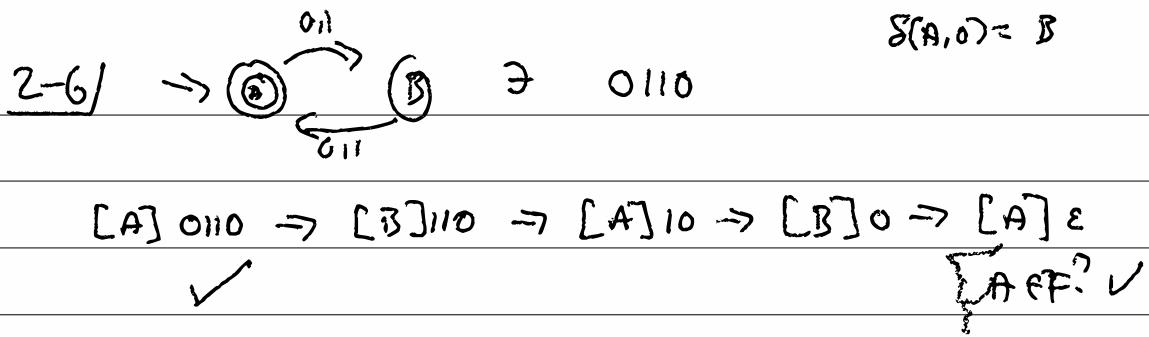
config update rule : $\xrightarrow{\text{DFA}}$ config $\xrightarrow{x} \text{config}$

$[q]w \rightarrow [q']w'$

$[q_i]cx \rightarrow [q_j]x \quad \text{iff} \quad \delta(q_i, c) = q_j$

$x \in \text{DFA} \quad \text{iff} \quad [q_0]x \xrightarrow{\delta} [q_f] \epsilon$

where $q_f \in F$



accepts : DFA \times string \Rightarrow bool

accepts $\stackrel{\text{def}}{=} (Q, \Sigma, q_0, \delta, F) \times = h \dashv q_0 \times$

h $\stackrel{\text{def}}{=} \text{DFA} \times \text{state} \times \text{string} = \text{bool}$

$h \stackrel{\text{def}}{=} (Q, \Sigma, q_0, \delta, F) \quad q_i \quad w =$

case w of $\epsilon \rightarrow q_i \in F?$

$c x \rightarrow h \dashv (\delta(q_i, c)) \times$

2-7/ class DFA \mathcal{C}

... $Q, \Sigma, q_0, \delta, F$

public bool accepts (String w) Σ

State $q_i = q_0;$

while ($\leftarrow 1 \text{ } w, \text{empty}() \right)$ {

$q_i = \delta(q_i, w, \text{first}())$

$w = w \text{rest}(); \}$

return ~~if~~ $F.i_n(q_i); \}$



