

$$\begin{array}{c} \text{19-1} \\ \diagdown \quad \nearrow \\ 0^a 1 0^b 1 0^{a+b} \end{array} \rightarrow \begin{array}{c} \downarrow \\ 0^{a-1} 1 0^b 1 0^{a+b} \end{array}$$
$$0^a 0^b 1 0^{a+b} = 0^{a+b} 1 0^{a+b}$$
$$\downarrow$$
$$w \# w \quad w \in 0^\times$$
$$\# =)$$

1a-2) A TM + = $(Q, \Sigma, \Gamma, \delta, q_0, q_a, q_r)$

Q - a finite set of states

Σ - an alphabet

Γ - an alphabet $\Sigma \subset \Gamma$

$w \in \Gamma$ $w \notin \Sigma$

$q_0, q_a, q_r \in Q$ - start, accept, reject

$\delta : Q \times \Gamma \rightarrow Q \times \Gamma \times \{\text{L}, \text{R}\}$
 $Q - \{q_a, q_r\}$

config: $\underset{\text{left}}{\Gamma^*} \times \underset{Q}{\text{Q}} \times \underset{\text{right}}{\Gamma^*} = c$

19-3) $w \in L(+) \text{ iff}$

$$\epsilon [g_0] w \Rightarrow^* x[g_a] y \quad x, y \in \Gamma^*$$

$$x[g_i] y \Rightarrow w x[g_i] y w$$

$$\frac{\delta(g_i, a) = (g_j, b, R)}{x[g_i] ay \Rightarrow x b [g_j] y}$$

$$\frac{\delta(g_i, a) = (g_j, b, L)}{x c [g_i] ay \Rightarrow x [g_j] cb y}$$

step : config \Rightarrow config

$$\text{step } (+l, g_i, +r) = (+l', g_j, +r')$$

$$\text{where } (a, y) = \text{look } tr$$

$$(g_i, b, d) = \delta(g_i, a)$$

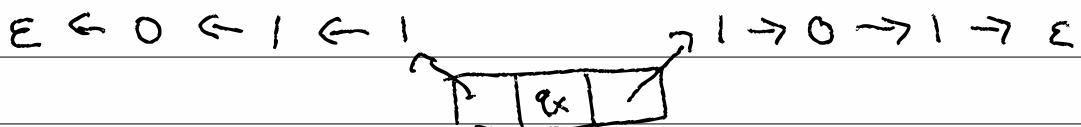
$$(+l', +r') = \begin{cases} \text{case } d \text{ of } R & ((b: +l), y) \\ & L \Rightarrow (x, c: b: y) \end{cases}$$

$$\text{look} : \text{list char} \Rightarrow (\text{char}, \text{list char}) \quad \text{where } (\text{B}, x) = \text{look } l$$

$$\text{look } \epsilon = (w, \epsilon)$$

$$\text{look } (a: x) = (a, x)$$

$$G \amalg [g_x] \amalg$$



19-21 / A computable function

f

is a Turing machine.

And

$$f(x) = y$$

iff

$$\varepsilon [g_0]_x \Rightarrow^* w [g_a]_y$$

$$\text{add 1} \quad 0 \leftarrow 1 \quad f(0^x + 0^y) = 0^{x+y}$$

$$1 \leftarrow 10$$

$$10 \leftarrow 11$$

...

19-5/ When we run a DFA on input w ,
how long could it take to accept
or reject?

$|w|$ steps \rightarrow Yes or
No

When we run a PDA?

$2^{|w|}$ steps \rightarrow Yes or No
or running forever
 \Rightarrow No

19-6 How long does a TM take?

Accept $\varepsilon [q_0]w \xrightarrow{*} \dots \xrightarrow{*} x [q_a]y$

Reject $\varepsilon [q_0]w \xrightarrow{*} \dots \xrightarrow{*} x' [q_r]y'$

Loop $\varepsilon [q_0]w \xrightarrow{*} \dots \xrightarrow{*} u [q_i]v \xrightarrow{*} \dots \xrightarrow{*} u [q_i]v$

=

Diverging $\forall x, q_i, y. \quad \varepsilon [q_0]w \xrightarrow{*} x [q_i]y$
implies $x [q_i]y \xrightarrow{*} x' [q_j]y'$
 $\exists x', q_j, y'. \quad s.t. \quad q_j \notin \{q_a, q_r\}$

19-7/ A TM is either

recognizer — may LOOP on
some input

decider — never LOOP

a language, A , is T-recognizable (Σ_1)
 \exists m recognizer, $L(m) = A$

T-decidable (Σ_0)
 $L(m) = B$

19-8 /

