

16-1 Macro systems

Our language is minimal and we implement things in the design as far as possible.

Threads were just in the stdlib.

... could have been made by user

$$\begin{array}{c} 2 \\ \downarrow \\ (+\ 2\ 2) \end{array} \rightarrow \left(\begin{array}{l} \text{let } x = 2 \\ (+\ x\ x) \end{array} \right) + \left(\begin{array}{l} \text{let } x = 3 \\ (+\ x\ x) \end{array} \right)$$

let double $x = x + x$

double 2 + double 6

16-2) desugar ["let", x, ":", xe, "in", b]
 $\Rightarrow (\lambda (x) b) xe$

$(\lambda (x) (+ x x)) 2$ rather let $x = 2$ in
 $x + x$

(define (let x xe b)
 ($\lambda (x) (xe) b$))

(let x 2 (+ x x)) \leftarrow not a program
 \hookrightarrow unbound \hookrightarrow unbound (free vars)

functions cannot introduce binding forms

16-3/ (extend-the-designer!

(let x xe b)

((lambda (x) b) xe))

C textual subst

#define MAX(x,y) x>y?x:y

MAX(3,4) \Rightarrow 3>4?3:4

#define TEMP(x) auto tmp = x

TEMP(3);

return tmp+2;

\Rightarrow

auto tmp = 3;

return tmp+2;

16-4/ #define LET(x, x-e) auto x = x-e
LET(x, 3); auto x = 3;
LET(y, 4); \Rightarrow auto y = 4;
return x+y; return x+y;

2* MAX(1+2, 3) $\rightarrow (2 \times 1) + (2 > 3) ? (1+2) : 3$

#define MAX(x, y) ((x) > (y) ? (x) : (y))

MAX(i++, 3) $\rightarrow ((i++) > (3) ? (i++) : (3))$

#define MAX(x, y) { auto tmpx = (x); auto tmpy = (y);
return x > y ? x : y; }

16-5 / Lisp / Scheme / Racket macros solve these problems.

define-syntax-rules

(define-syntax-rules

[(let ([x xe] ...) be)

(λ (x ...) be) xe ...]])

(let ([x 2] [y 3]) (+ x y))

\Rightarrow (λ (x y) (+ x y)) 2 3)

16-6 / (define-syntax-rules
 [(or) false] ← template
 [(or x y ...)] ← pattern
 (let ([tmp x])
 (if tmp tmp (or y ...))))

(define-syntax-rules

[(and) true]

Pattern → [(and x) x] ← macro name

case → [(and x y ...)]

clauses (if x (and y ...) false))

16-7 define-syntax-rules : id x List(pair (pattern
, template))

letm := "let", [< (let ([x xe] ...) be),
(λ (x ...) be) xe ... >]

desugar will track a database of macros
ie a mapping from id to
List(pair (pattern, template))

desugar M [m, more ...] where $m \in \mathcal{M} :=$
expand M[m] [m, more ...]

16-8/

$\text{expand } [\langle \text{let } ([x \ xe] \dots) \text{ be} \rangle, \text{ ← def}$
 $\text{template } \rightarrow [(\lambda (x \dots) \text{ be}) \text{ xe} \dots]$
 $(\text{let } ([z \ 2]) (\text{+ } z \ z)) \text{ ← use}$
 $= ((\lambda (z) (\text{+ } z \ z)) \ 2)$

match : pattern x sexpr → maybe (env)

env : id → sexpr

transcribe : template x env → sexpr

16-9/ expand [< (let ([x xe] ...) be), ← def
 template → ((λ (x ...) be) xe ...) >
 (let ([z 2]) (+ z z)) ← use
 = ((λ (z) (+ z z)) 2)

match (let ([x xe] ...) be)
 (let ([z 2]) (+ z z))
 = [x ↦ (z), xe ↦ (2), be ↦ (+ z z)]

transcribe ((λ (x ...) be) xe ...) ↑
 = ((λ (z) (+ z z)) 2)

$$m \text{ '() '() } = \emptyset$$

16-10/ $m \text{ (cons pa pd) (cons ua ud) } =$
 $m \text{ pa ua } \oplus m \text{ pd ud}$

$$m \text{ var(x) u} = [x \mapsto u]$$

$$m \text{ const(n) n} = \emptyset$$

$$m \text{ (list part ...) u} = \text{merge (map (m part) u)}$$

$$\text{tr '() env} = \text{'()}$$

$$\text{tr (cons ta td) env} = \text{cons (tr ta env) (tr td env)}$$

$$\text{tr var(x) env} = \text{env(x)}$$

$$\text{tr const(n) env} = \text{const(n)}$$

$$\text{tr (list fmp ...) env} = \text{map (tr fmp) (unmerge env)}$$

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Hygiene
Macro

16-11 / (define-syntax-rules

[(or) false]
[(or x) x]
[(or x y ...)]

(let ([tmp₁ x])
 (if tmp₁ tmp₁
 (or y ...))))

capture!



timestamps

↓ colors

→ sets of
scopes

(let ([tmp₀ 4])
 (or false tmp₀))
 ↓
 4

⇒

(let ([tmp₀ 4])
 (let ([tmp₁ false])
 (if tmp₁ tmp₁
 tmp₀)))
 ⇒ false