

# TY-1) functions

$ty := \dots \mid (\rightarrow (ty \dots) \overset{\text{dom}}{ty})$   
 $\neq$   
 $\text{rng}$

$R_3 \Rightarrow R_4$

$e := \dots \mid (\text{app } e \ e \dots)$   
 $\text{def} := (\text{define } (\text{var } [var:ty] \dots) : ty \ e)$

example:

```

(define (even? [x:SBY]) : B
  (if (== x 0) true
      (odd? (- x 1))))
(define (odd? [x:SBY]) : B
  (if (== x 0) false
      (even? (- x 1))))
(even? 9a)
  
```

~~program := (program info e)~~  
 (program info (def... ) e)

type-check  
 old:  $\Gamma \vdash e : ty \rightsquigarrow e'$   
 new:  $\Delta, \Gamma \vdash e : ty \rightsquigarrow e'$   
 $\uparrow$  top level funs       $\uparrow$  local vars + args

$$\frac{\lambda(x) = \emptyset \quad \Delta(x) = ty}{\Delta, \Gamma \vdash x : t \rightsquigarrow (\text{fun-ref } x) \neq}$$

$\Delta = \emptyset [dname \Rightarrow (\rightarrow \text{dom } \text{rng})] \dots$

$\Delta, \emptyset \vdash d \rightsquigarrow d' \dots$

$\Delta, \emptyset \vdash e \rightsquigarrow e'$

$\Delta, \Gamma \vdash \text{emtor} : (\rightarrow \text{doms } \text{rng})$

$\Delta, \Gamma \vdash \text{erands} : \text{doms } \dots$

$\Delta, \Gamma \vdash (\text{app } \text{emtor } \text{erands} \dots) : \text{rng}$

$\emptyset, \emptyset \vdash (\text{program } i \ (d \dots) \ e) \rightsquigarrow (\text{program } i \ [v \Rightarrow h \ \dots] \ (dr \dots) \ er)$

M-2) (define (f [x: S0Y]) : S0Y (+ x 1))  
 (define (g [x: S0Y]) : S0Y (\* x x))  
 (let z := read in let y := read in  
 @ (if (= z 0) f g) y))

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random generation

1/2 - no fns      1/2 - yes fns

if yes → random number ∈ [1, 12] }  
 how args?      0, [1, 6], [10, 14] } int → int

random result type (same colors for whole program)

old: mapping type → variable (n=0)  
 new: mapping type → fn that returns it (n=1+)  
 → one option is a function

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Optimization — inlining

error opt (fun-ref F)  
 $\Delta(F) = (\text{define } (F \text{ [args: ty]}))$

(app error e<sub>0</sub> ... e<sub>n</sub>)

⇒ if : my ebody)

(let\* arg<sub>0</sub> := e'<sub>0</sub> in  
 arg<sub>n</sub> := e'<sub>n</sub> in  
 ebody)

e<sub>0</sub> ... e<sub>n</sub> opt ⇒ e'<sub>0</sub> ... e'<sub>n</sub> }  
 e'<sub>0</sub> is simple ... }  
 e'<sub>n</sub> is simple . }  
 constant or variable

14-3 | (define (f [x: 504]) : 504

(if (= x 0) 42

(f (- x 1))))

(f 50)  $\xrightarrow{SB}$  42

$\xrightarrow{E}$  (f 40)

(f -1)  $\xrightarrow{E}$  (f -11)

$\xrightarrow{SB}$  (f -51)

$\xrightarrow{E}$  ⊥

Uniquity

→ rename

funcs

reveal - fun →

ensures that

fun-refs are

tagged

f ⇒ (funref f)

own / typec / uniqueness

new-pass : limit-fun

x86 : <sup>always</sup> funcs have 6 args (rdi rsi rdx rcx r8 r9)

R : funcs have no limit to args (f a b c ... x y z) <sup>old</sup>

⇒ (f a b c d e <sub>new</sub>

(vector f ... x y z))

(define (f [a: ta] [b: tb] ... [z: tz]) : rtype)

⇒

(define (f [a: ta] ... [e: te] [rest: (vector ref ...)])

: rtype (let\* (f := (vector-ref rest 0)

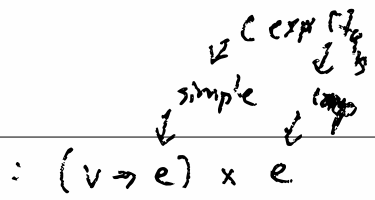
...

z := (vector-ref rest 20) in

ebody)

(A)!

14-4 extend rco pass



$(v \rightarrow e) \times e$

rco  $\sigma$  tail? (app rator rands) :=

rands-res := map rec rands

nv\_rator, tr\_rator := rec rator

all-nvs = nv\_rator ++ all nvs in rands-res

if tail? then

all-nvs, (rator' rands')

O.W.

all-nvs ++ [ans  $\mapsto$  (vator' rands')],  
ans

also ...

(program lets e)

$\Rightarrow$

(program defs ++ (main = e')  
(main))

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