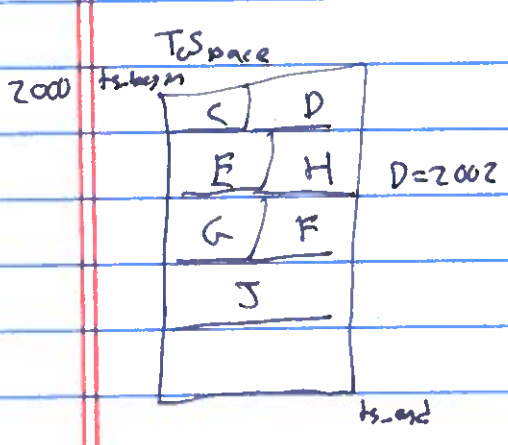
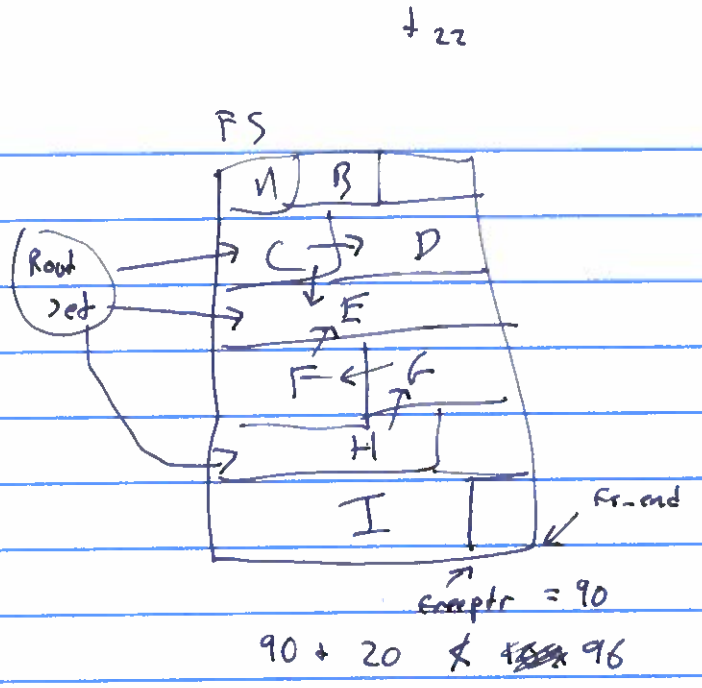
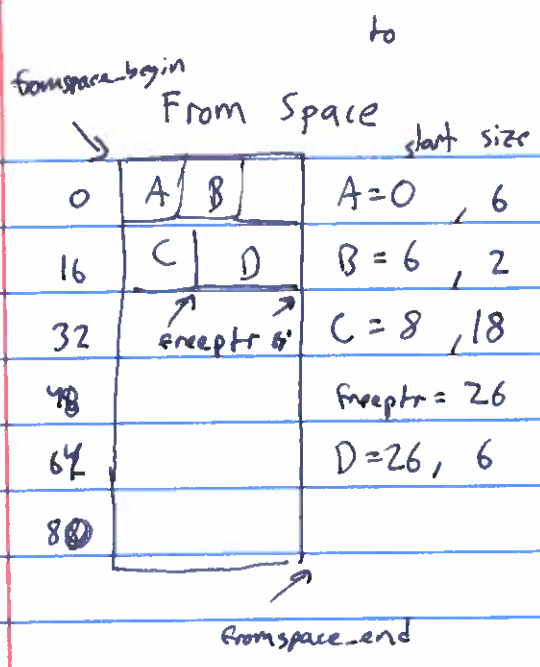


3-1/

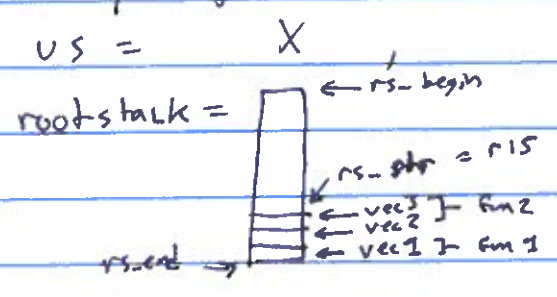
Stop + Copy



- switch the role of from & to
- Implementation issues
- What's the root set?
 - What do objects point to?
 - How do update refs
 - Record already moved
 - check if succ for ref
 - how to track seen/unseen

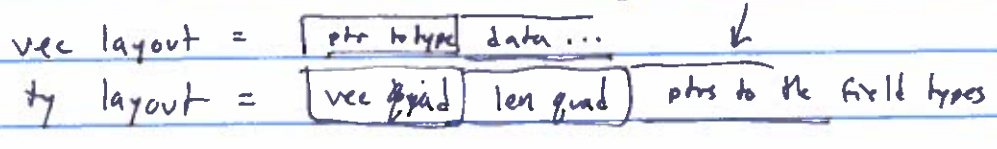
Rootset

theory = global vars, ptrs in the program src, local vars (things on stack) registers, objects w/ ptrs are on rootset



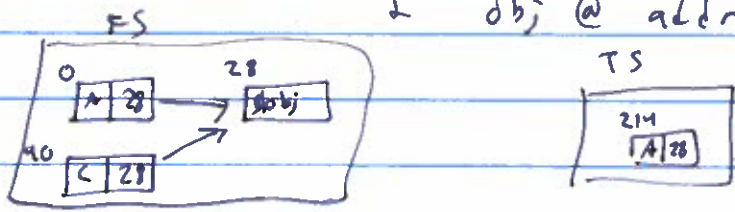
collect (curr_rs_ptr, alloc_req)

What do objects point to?

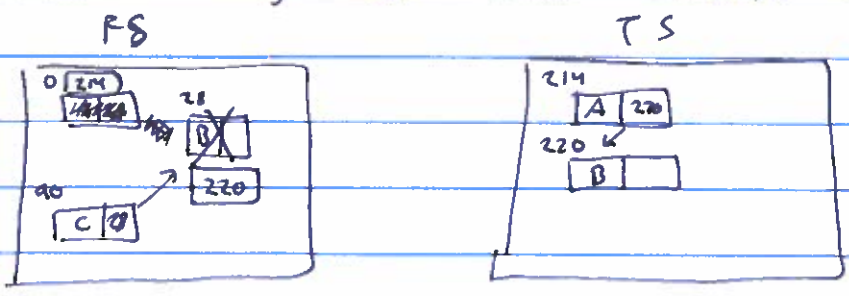


something already copied

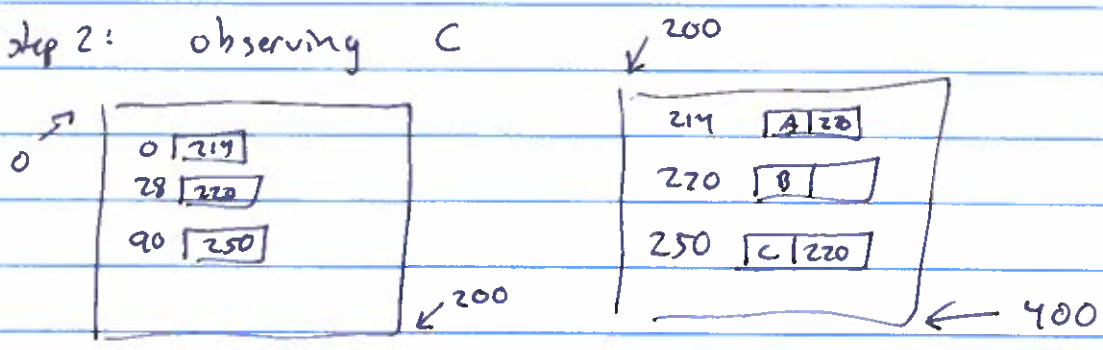
before: obj @ addr 28
pointed by obj @ addr 0
+ obj @ addr 90



step 1: observing obj @ 0 and doing copy



step 2: observing C



$A = \text{addr } 0 =$ [ptr to type] [field 1] ... [field N]
64 bits

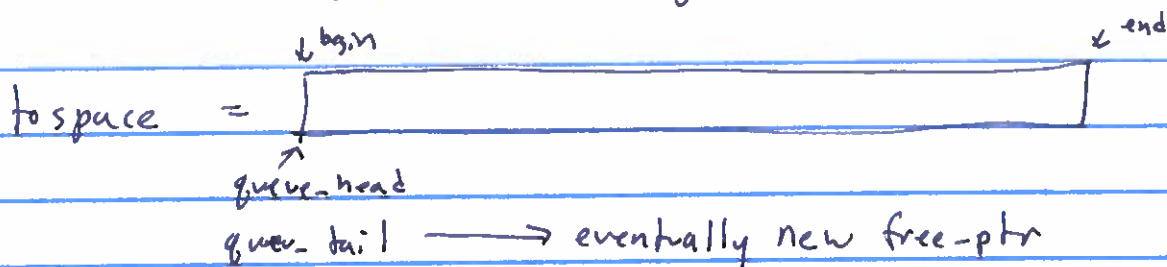
$A' = \text{addr } 0 =$ [214]

[28]

3-3

How to do DFS w/o memory?

- use the tospace as a queue



step 1: enqueue all elements on root stack
head ≠ tail (if there were objects)

step 2: while (head ≠ tail)
scan the head object

enqueue (int64_t ** objref_ptr)

```

- new_loc = queue_tail - obj = *objref_ptr
- tag = obj[0] (either type defn ptr OR fwd)
if tag is fwd
    size = 0, new_loc = tag
o.w. do the copy
    size = tag[1] + 1
    for (...) { g + [i] = obj[i]; }
queue_tail += size
update ptr => *objref_ptr = new_loc
install fwd => *obj = new_loc

```

13-9

scan()

read queue-head (obj = queue-head)

tag = obj[0]

for (i = 0, ..., tag[1]) {

elem = obj[1+i]

elem tag = tag[2+i]

if elem tag \neq vector

~~copy~~ ignore

o.w. enqueue (obj[1+i])

queue-head += size (= tag[1] + 1)