compiler : A ⇒ B

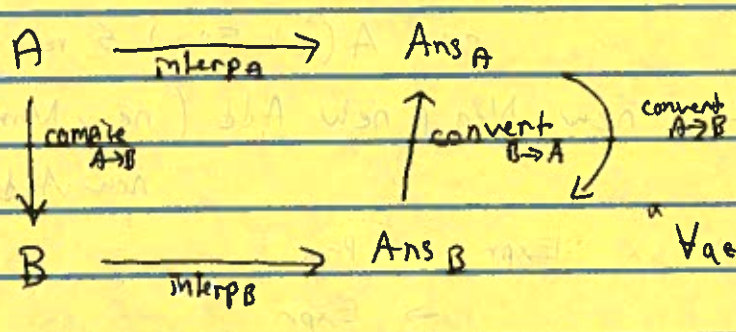
A is a program in lang X

B is a program in lang Y

gcc : C-prog ⇒ X-86-prog

interp : A ⇒ answer

When is a compiler right?



~~tagram~~ "commute 5"

Cog

"$\forall a \in A.\ \text{interp}_A(a) = \text{interp}_B(\text{compile}_{A\to B}(a))$"

When is a compiler good?

some other metric (time, space) m (measures goodness)

$m(A) < m(\text{compile}(A))$

large gap ⟶ good compiler

"optimizations" ways of improving m

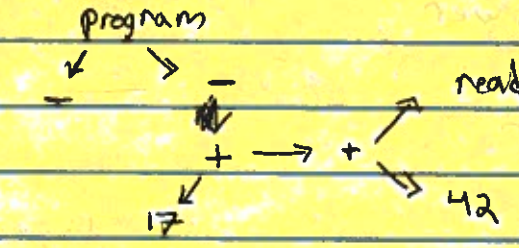
1-2/ R_0 : p = program info e

e = number | (- e) | (+ e e) | (read)

"return - 17 + read() + 42;"

→

* (program - (- (+ 17 (+ (read) 42))))



B* A(B*, B*) { return new Add(1, 2); };

new Program (NULL, new Neg (new Add (new Num(17),
new Add (new Read(),
new Num(42)))));

Program : info-ptr x Expr → Prog

Neg : Expr → Expr

Add : Expr x Expr → Expr

Read : → Expr

Num : int → Expr

interp (program : e) = interp (e) interp : R_0 → int

interp (Num n) = n

interp (Neg e) = -1 * interp (e)

interp (Add e_1 e_2) = interp (e_1) + interp (e_2) ←

interp (Read) = ask the user for a number

(+ (read) (- (read)))

class Expr {
virtual int interp() = 0; }

class Add {
Expr *lhs, *rhs;
int interp() {
return lhs->interp() + rhs->interp(); }

3)

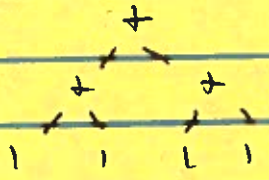
$2^n : int \to R_0$

interp $(2^n\ m) = 2^m$

$2^n\ 0 = 1$

$2^n\ (1+m) = (+\ (2^n\ m)\ (2^n\ m))$

$2^n\ 2 =$



randp : int $\to R_0$

randp 0 = flip a coin head → (read)
tail → number ∈ {0, 1024}

randp (1+n) = flip a coin head → − (randp n)
tail → + (randp n) (randp n)

opt : $R_0 \to R_0$

opt (+ 2 3) = 5

opt (− (− 7)) = 7

(− (− e)) = opt e

opt (+ 7 (+ read 8)) = (+ 15 read)

opt n = n

opt read = read

opt (+ e_1 e_2) = e'_1 = opt e_1 , e'_2 = opt e_2

cases e'_1, e'_2 ~~40% done~~

n, m → n+m

n, (+ m e) → (+ (n+m) e)

opt (− e_a) = e'_a = opt e_a 80%

cases e'_a

(+ e −e) → (+ −n (− e)) | (− e) → e