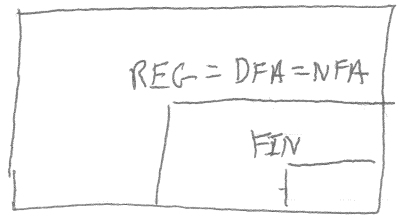


6-1/

ALL



Categories of Languages	Computers i.e. ϵ	Specifications or programs i.e. $=$
ALL		
FIN	list of values	
REG	DFA, NFA	???

$\epsilon : \Sigma^* \rightarrow Y/N$ — identifies members
 $= : \rightarrow P(\Sigma^*)$ — produces the entire set

A regular expression over Σ is either:

ϵ	$r_1 \cup r_2$	where r_1, r_2 are regexp
\emptyset	$r_1 \circ r_2$	
$c \in \Sigma$	r_1^*	

$r_1 r_2$	r_1^R
$\overline{r_1}$	Optional

A regexp is a program written using the regular operators = (operations the regular languages are closed under)

$L : \text{regexp} \rightarrow P(\Sigma^*)$

$L(\epsilon) = \{\epsilon\}$ $L(r_1 \cup r_2) = L(r_1) \cup L(r_2)$
 $L(\emptyset) = \emptyset$ $L(r_1 \circ r_2) = L(r_1) \circ L(r_2)$
 $L(c) = \{c\}$ $L(r_1^*) = L(r_1)^*$

Think from end is 1:

$(1 \cup 0)^* \circ 1 \circ (1 \cup 0) \circ (1 \cup 0)$

unix: $*1??$

perl regexp:

$\cdot * 1 \dots$

$((aob) \cup (cod)) \circ e \circ f \circ g$
 $(0 \cup 1 \cup 2 \cup 3 \cup 4 \cup 5 \cup 6 \cup 7 \cup 8 \cup 9)$

$(ab|cd)ef \setminus 1$
 $\setminus d$

6-2/

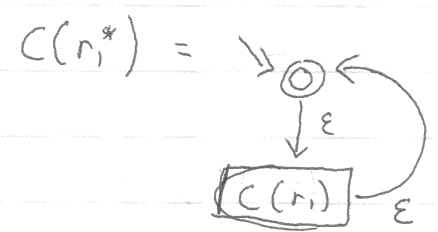
DFA = REG

⇒ decompiles a DFA into a REG (hard)

⇐ compiles a REG into a DFA (easy)

REG → NFA → DFA

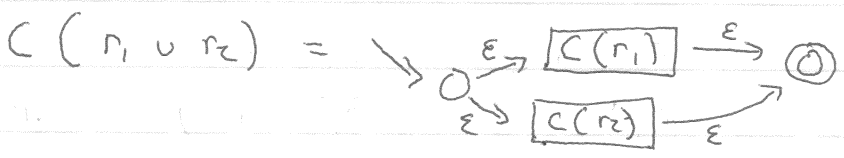
$C : REG \rightarrow NFA$



$C(\epsilon) = \rightarrow \odot$

$C(\emptyset) = \rightarrow \emptyset$

$C(c \in \Sigma) = \rightarrow \circ \xrightarrow{c} \odot$



$C(r_1 \circ r_2) = \rightarrow [C(r_1)] \xrightarrow{\epsilon} [C(r_2)] \xrightarrow{\epsilon} \odot$

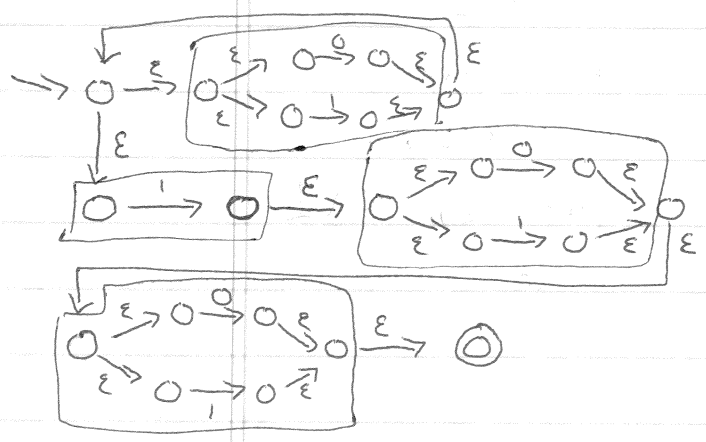
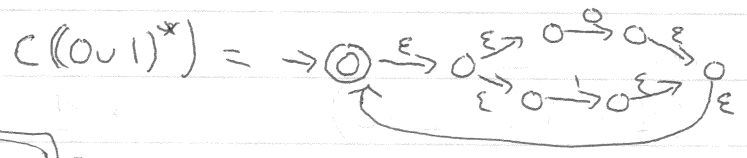
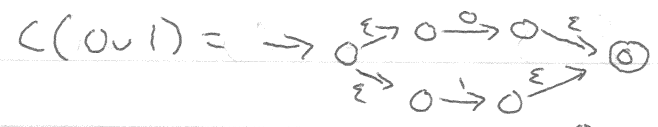
$C((0 \cup 1)^* \circ (1 \circ ((0 \cup 1) \circ (0 \cup 1)))) =$

$\rightarrow [C((0 \cup 1)^*)] \rightarrow [C(1 \circ ((0 \cup 1) \circ (0 \cup 1)))] \xrightarrow{\epsilon} \odot$

$C(0) = \rightarrow \circ \xrightarrow{0} \odot$

$C(1) = \rightarrow \circ \xrightarrow{1} \odot$

~~Diagram~~



$= \rightarrow \odot \xrightarrow{0,1} \circ \xrightarrow{0,1} \circ \xrightarrow{0,1} \odot$

* 1??