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$0^n 1 0^m 1 0^{n \times m}$ - multiplication machine

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 $\langle n \times m = 0 \rangle$ s.t. $0 = n \times m$

$0^n 1 0^m 1 0^0$

$0^n * 0^m = 0^0$

n is bin \times m is bin $=$ 0 in binary

n in roman

n in arabic

A_{DFA} is the "acceptance" problem on DFAs

$\langle D, w \rangle \in A_{DFA}$ iff w is accepted by the DFA D .

A_{DFA} is DFA interpreter

$\forall D \in DFA. \exists m \in TM. L(m) = L(D)$

in: $(Q, \Sigma, q_0, \delta, F)$

DFAs

out: $(Q', \Sigma, \Gamma, q'_0, \delta', q_a, q_r)$

compiler

$Q' = Q \cup \{q_a, q_r\}$

$\Gamma = \Sigma \cup \{\omega\}$

$q'_0 = q_0$

$\delta'(q_i, a) = (\delta(q_i, a), \omega, R)$ $a \in \Sigma$

$\delta'(q_i, \omega) = q_a$ if $q_i \in F$

q_r o.w.

tape 1: $\langle D, w \rangle = \langle (Q, \Sigma, q_0, \delta, F), w \rangle$

tape 2: $\langle q_0 \rangle \rightarrow \langle q_1 \rangle \rightarrow \langle q_2 \rangle \rightarrow$

tape 3: w

Meta-program - a program about programs

interpreter : $Prog \rightarrow Ans$

compiler : $Prog_A \rightarrow Prog_B$

analysis : $Prog \rightarrow Bool$

generator : $Nvm \rightarrow Prog$

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$$A_{NFA} = \{ \langle N, w \rangle \mid w \in L(N) \}$$

$$M_{ANFA} = \text{COMPILER}_{NFA, DFA} \circ M_{ADFA}$$

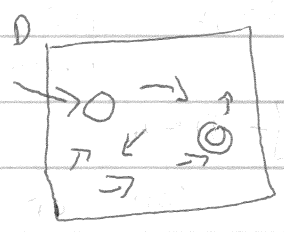
$$\text{Trans NFA DFA} = \{ \langle N, D \rangle \mid N \in \text{NFA}, D \in \text{DFA}, L(N) = L(D) \}$$

$$N \rightarrow D \text{ st. } L(D) = L(N)$$

A_{DFA} A_{NFA} A_{REG} $\in \Sigma_0?$ $\in \Sigma_1?$

✓

$$E_{DFA} \quad \langle D \rangle \in E_{DFA} \text{ iff } L(D) = \emptyset$$



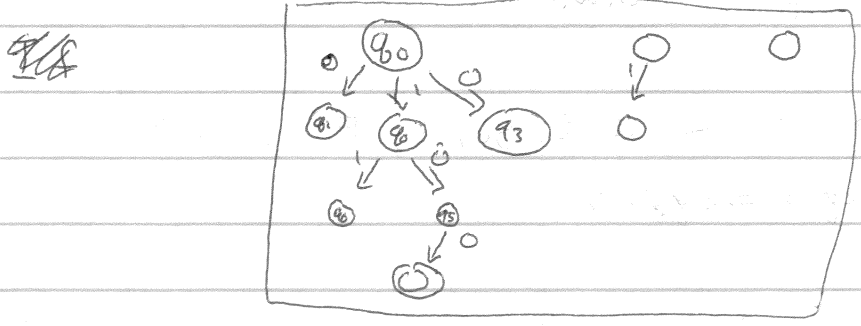
Is there a path from q_0 to something in F ?

$V = Q \quad (x, y) \in E \text{ iff } \delta(x, a) = y \text{ for some } a$

What is the spanning from q_0 ? What is reachable from q_0 ?

$$R(A, Q \cup \{q_i\}) = R(A \cup \{q_i\}, Q \cup \{q_j \mid \exists a. \delta(q_i, a) = q_j\})$$

$$R(A, \emptyset) = A \quad R(\emptyset, \{q_0\}) \cap F = \emptyset?$$



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$$EQ_{DFA} = \{ \langle D_1, D_2 \rangle \mid D_x \leftarrow DFA \text{ and } L(D_1) = L(D_2) \}$$

$$E_{DFA}(\langle C \rangle) = Y \text{ iff } L(D_1) = L(D_2)$$

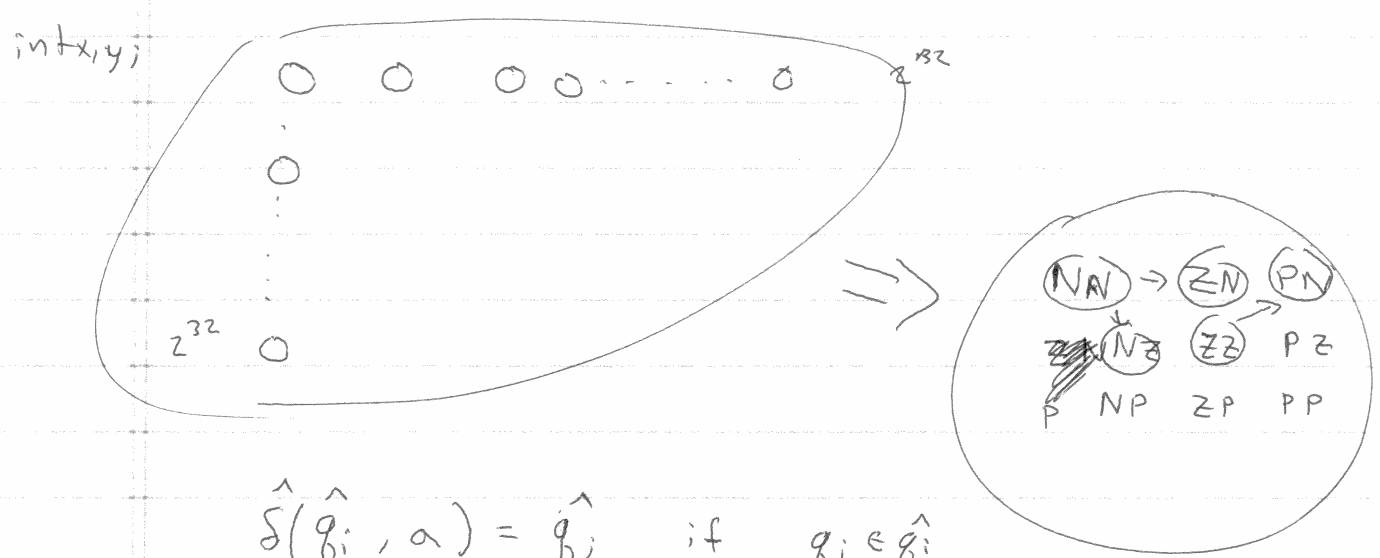
$C = D_1 \cap D_2$ — nothing in common — want everything

$C = \bar{D}_1 \cap D_2$ — stuff that D_2 does but D_1 doesn't
— want empty

$C = (\bar{D}_1 \cap D_2) \cup (D_1 \cap \bar{D}_2)$ — all mismatches → want empty

$$(N \times N) \times (N \times N) = N^4$$

Model Checking



$$\hat{\delta}(\hat{q}_i, a) = \hat{q}_j \text{ if } q_i \in \hat{q}_i$$

$$q_j \in \delta(q_i, a)$$

