

6-1/

$$M = X \quad V = \lambda X.M$$

$$| \lambda X.M$$

$$| b$$

$$| (M, N)$$

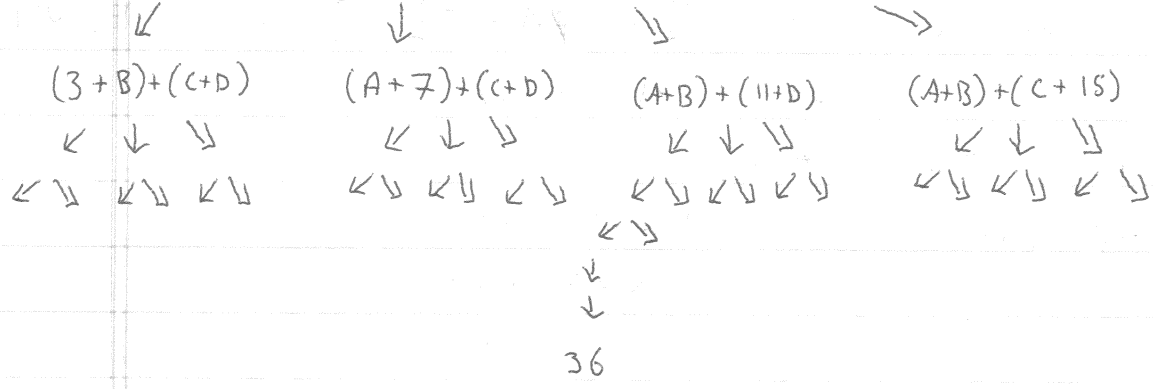
$$C = \lambda X.C \mid (C N) \mid (M C) \mid (0^n M \dots (M \dots))$$

$$| b$$

$$| (0^n M \dots)$$

$$\beta_v : (\lambda X.M) V \rightarrow M[X \leftarrow V]$$

$$\left( \left( (A+B) + (C+D) \right) + \left( (E+F) + (G+H) \right) \right)$$



$$\text{interp}^* : M \rightarrow M$$

$n^3$  where  $n$  is size of program

$$\text{interp}^* m = m \quad \text{if } \text{interp } m \text{ is false}$$

$$\text{o.w. } \text{interp}^* (\text{interp } m)$$

$$\text{interp} : M \Rightarrow M \text{ or } \# \text{false}$$

- interp  $m = \text{OR}$  (if  $m$  is  $(\lambda X.M) V$ ) then  $m[X \leftarrow V]$
- (if  $m$  is  $(\lambda X.M)$  and  $m' = \text{interp } M$ , then  $\lambda X.m'$ )
- (if  $m$  is  $(M^* N)$  and  $m' = \text{interp } M$ , then  $(m' N)$ )
- (if  $m$  is  $(M = N)$  and  $m' = \text{interp } N$ , then  $(M m')$ )
- (if  $m$  is  $(0^n N \dots M 0 \dots)$  and  $m' = \text{interp } M$ , then  $(0^n N \dots m' 0 \dots)$ )

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Evaluation Context - the place where work happens

$$E ::= \square \mid (E N) \mid (V E) \mid \dots \mid \text{or } \dots$$

$$C ::= \square \mid (\lambda X. C) \mid (C N) \mid (M C) \mid \dots \mid \text{or } \dots$$

$\mapsto_v$  (the standard reduction)

$$M \mapsto_v N \text{ iff } \exists E, M' \text{ s.t. } M = E[M'] \text{ and } N = E[N'] \text{ and } M' \mapsto_v N'$$

interp :  $M \rightarrow M$  or #false

interp (M N) = if  $M \in V$  then  
if  $N \in V$  then  
beta M N

o.w. (M (interp N))

o.w. (interp M) N

$$\text{interp } (\text{or } V_1 \dots M_0 M \dots) = (\text{or } V_1 \dots (\text{interp } M_0) M \dots)$$

Uniqueness of Eval Ctxts.

$\forall M. M = V$  or there exists a unique  $E$  s.t.

$$M = E[(V_1 V_2)] \text{ or } M = E[(\text{or } V_1 \dots V_n)]$$

Correctness

$$\forall M \text{ and } U. M \xrightarrow{\text{old way}} U \text{ iff } M \xrightarrow{\text{new way}} V \text{ and } V \xrightarrow{\text{new way}} U$$

Stuck :  $M$  is stuck if

- $M = (\text{or } b_1 \dots b_n)$  and  $\delta(\text{or } b_1 \dots b_n)$  is undefined
- or  $M = (\text{or } b_1 \dots (\lambda X. M) V \dots)$
- or  $M = (b V)$

Uniform Evaluation Theorem

$$M \xrightarrow{\text{new way}} V$$

If  $M$  is closed ( $FV(M) = \emptyset$ ), then  $M \xrightarrow{\text{new way}} N$  where  $N$  is stuck

$$(\forall N. M \xrightarrow{\text{new way}} N \Rightarrow \exists L. N \xrightarrow{\text{new way}} L)$$