

try M_1 , finally M_2

\Rightarrow catch M_1 with $(\lambda x. \text{seq}(M_2, \text{throw } x))$

$\text{seq}(M_1, M_2) = (\lambda x. M_2) M_1$
where $x \notin \text{FV}(M_2)$

Resume - ISWIM

$M := \dots \mid \text{throw } M \mid \text{catch } M_1 \text{ with } (\lambda x. \lambda R. M_2)$ ↗ resumer
 $E := \dots \mid \text{throw } E$ ↓ thrown value

$E[\text{catch } F[\text{throw } v] \text{ with } (\lambda x. \lambda R. M_2)]$
 $\mapsto E[(M_2 \ v) (\lambda z. F[z])]$
(+ 1 (R 3)) no longer forgotten, turned into a function "resified context"

$K := \dots \mid \text{catch}(E, \lambda x. \lambda R. M_2, K) \mid \text{throw}(k)$
 $V := b \mid \text{clo}(\lambda x. M, E) \mid \text{kont}(k)$

~~$\leftarrow \text{throw}$~~ $\langle V, E, \text{throw}(k) \rangle \quad \langle E', k', M_2, R \rangle$
 $\mapsto \langle (M_2 \ v) \ R, E', k' \rangle \quad := \text{search}(K, \text{ret})$

- $\text{search}(\text{fun}(V, k_1) \equiv, k_2) = \text{search}(k_1, \text{fun}(V, k_2))$
- $\text{search}(\text{arg}(E, M, k_1), k_2) = \text{search}(k_1, \text{arg}(E, M, k_2))$
- $\text{search}(\text{ret}, k_2) = \langle \emptyset, \text{ret}, \lambda x. \lambda R. x, \text{kont}(k_2) \rangle$
- $\text{search}(\text{catch}(E, M, k_1), k_2) = \langle E, k_1, M, \text{kont}(k_2) \rangle$

$\langle V, E, \text{fun}(\text{kont}(k_2), k_1) \rangle \mapsto \langle V, E, k_2 \rangle$
 $M := \dots \mid \text{abort } M$

$\mapsto E[(M_2 \ v) (\lambda z. \text{abort } E[F[z]])]$
 $E[\text{abort } M] \mapsto M$

First-Class Continuations

- Continuations are first class in the language

~~M = V | X | (M N) | (o^n M ...) | callcc M~~

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V = b | (λX. M) | K E

E = ~~λ~~ | (E N) | (V E) | (o^n V ... E M ...) | callcc E

- $E[(\lambda X. M) V] \mapsto E[M[X \leftarrow V]]$
- $E[(o^n b \dots)] \mapsto E[S(o^n, b \dots)]$
- $E[\text{callcc } V] \mapsto E[V(K E)]$ ← copied
- $E[(K E') V] \mapsto E'[V]$ ← forgotten

$\Gamma \text{catch } M_1 \text{ with } (\lambda X. M_2) \Uparrow$ $\Gamma \text{throw } M \Uparrow$
 \Rightarrow ~~callcc (λK. ...)~~ $\Rightarrow \text{THROW } M$
 $\text{callcc } (\lambda K.$
 $\quad \text{let } \text{THROW} = (\lambda X. (K M_2)) \text{ in}$
 $\quad M_1)$

callcc = call with current continuation