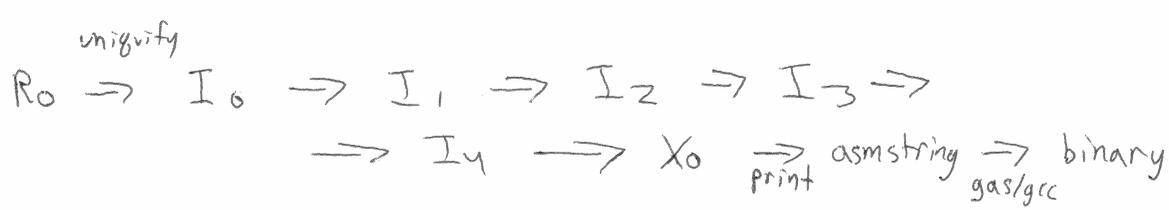




| R0  | X0  |
|---|---|
| operations take expressions<br>(+ (+ 1 1)<br>(+ 2 3)) | inst take 1 or 2 <u>arg</u><br>and arg are dest or src<br>(ie atomic)<br>inst take dest and maybe src |
| expr evaluate to a num<br>Ans = num                   | x expr evaluates to an effect (a change in memory)<br>Ans = void                                      |
| infinite variables<br>(map to mem?)                   | finite registers + finite memory  |
| tree-shaped   | sequential  |

shadowing  
 (let x 5  
 (let x 6  
 (x)))

register + memory are globally named



X0 → asm-string — print  
 job: turn a tree into a string easy

R0 → I0 — uniquify I0 = R0 but no dupe vars  
 job: remove shadowing

(let ([x 10]) (+ x (let ([x 20]) x)))

↓                    ↓                    ↓                    ↓  
 x1                    x1                    x2                    x2

uniquify: (dupe var → no-dupe var) e

2-3/

(+ (let ([x 7]) 5) x)

↓  
x0

↓  
consult mapping  
but it's empty  
so we error

$I_0 \Rightarrow I_1$  (  $R_0$  w/o dupe vars  $\rightarrow C_0$  )  
job: remove nested expressions — flatten

$C_0 = (\text{program } (\text{var } *) \text{ stmt } + )$   
 $\text{stmt} = ( := \text{ var } \text{ exp} ) \mid ( \text{ret } \text{ arg} )$   
 $\text{exp} = \text{ arg } \mid ( \text{read} ) \mid ( - \text{ arg} ) \mid ( + \text{ arg } \text{ arg} )$   
 $\text{arg} = \text{ int } \mid \text{ var}$

(+ (+ 1 2) (+ 3 4))  $\Rightarrow$   $x = 1 + 2$        $\text{add\_rhs } q1 = 1 + 2$   
 $y = 3 + 4$        $\text{add\_rhs } q2 = 3 + 4$   
 $z = x + y$        $\text{ans\_to\_prog } q3 = \dots$   
 $\text{ret } z$        $\text{ret } \dots$

continue next  
time ...

