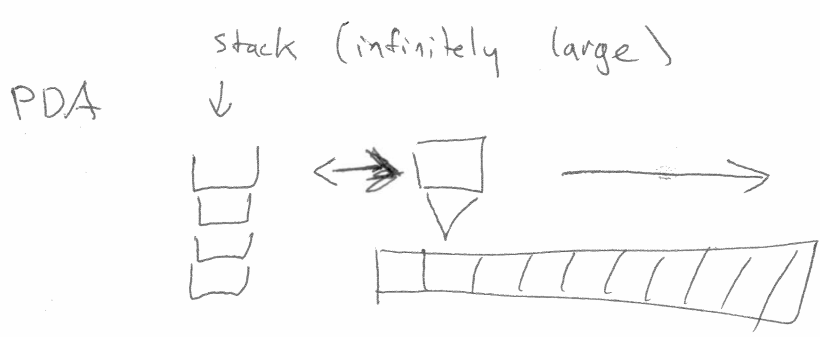
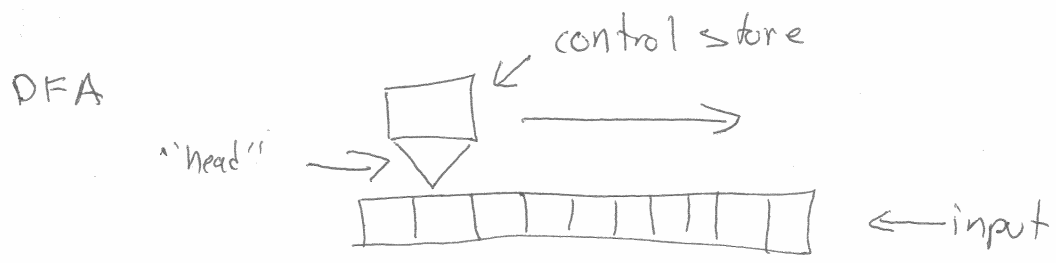


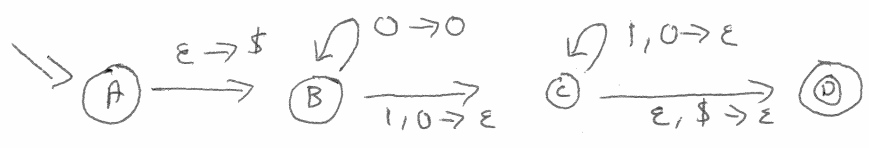
9-1/

DFA \approx REG \approx PDA \approx CFG

Push-down Automata



$0^n 1^n$ $S \rightarrow \epsilon \mid 0S1$



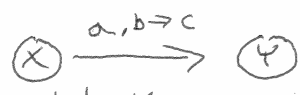
$a \rightarrow c = a, \epsilon \rightarrow c$

DFA



In state X, on an 'a' in input, goto Y

PDA



In state X, on an 'a' in input with 'b' on stack top, goto Y and replace stack top with 'c'

in: 000111 stack [state] input remaining

[A] 000111 \Rightarrow \$ [B] 000111 \Rightarrow \$ 0 [B] 00111 \Rightarrow \Rightarrow \$ 00 [B] 111 \Rightarrow \$ 00 [C] 11 \Rightarrow \Rightarrow \$ [C] \Rightarrow [D]

[A] 11 \Rightarrow \$ [B] 11

9-2/

$$PDA P = (Q, \Sigma, \Gamma, q_0, \delta, F)$$

Q = finite states

Σ = alphabet (input)

stack

Γ = ~~finite~~ alphabet

$q_0 \in Q$

$= \{0, \epsilon, \$\}$

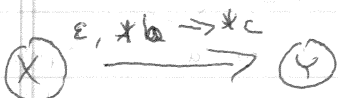
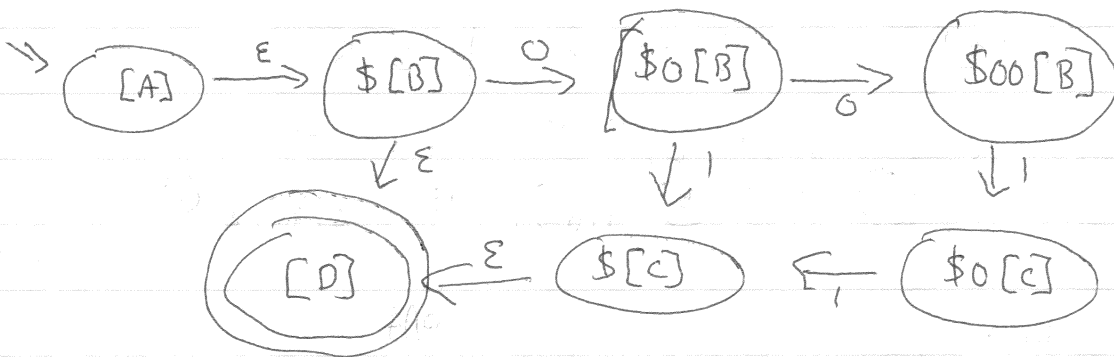
$F \subseteq Q$

$$\delta: Q \times \Sigma_{\epsilon} \times \Gamma_{\epsilon} \rightarrow P(Q \times \Gamma_{\epsilon})$$

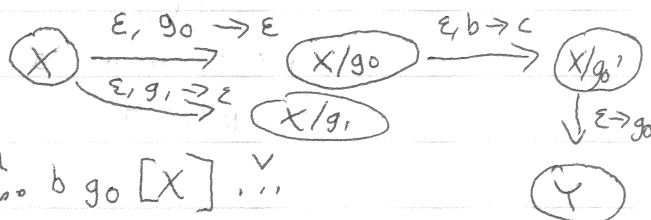
$$\delta(A, 0, \epsilon) = \{(B, 0)\}$$

In A, on 0, ignore stack, goto B, push 0

	<u>read input</u>		<u>ignore input</u>
push:	$a \rightarrow c$	\leftrightarrow	$\epsilon \rightarrow c$
pop:	$a, b \rightarrow \epsilon$		$\epsilon, b \rightarrow \epsilon$
replace:	$a, b \rightarrow c$		$\epsilon, b \rightarrow c$



On X, if 2nd stack is b,
replace with c, and goto Y



$$\begin{aligned}
 & u \dots b q_0 [X] \dots v \\
 \Rightarrow & u b [X/q_0] v \\
 \Rightarrow & u c [X/q_0'] v \\
 \Rightarrow & u c q_0 [Y] v
 \end{aligned}$$

9-3/

$$L(P) = \{ w \mid \exists [q_0] w \Rightarrow^* u [q_f] \epsilon \text{ s.t. } q_f \in F \}$$

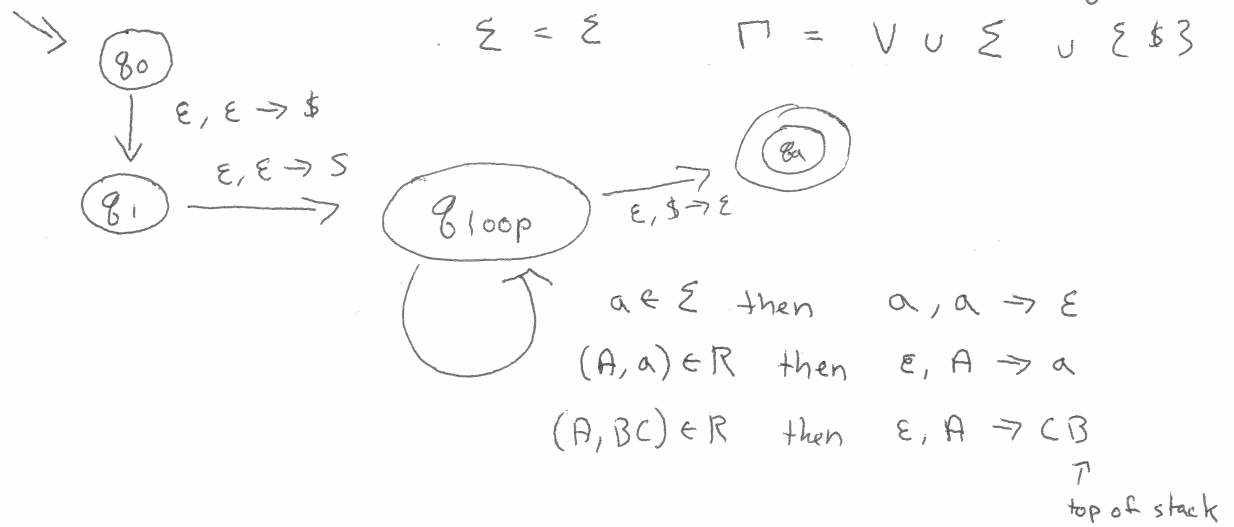
$$u b [q_i] a v \Rightarrow u c [q_j] v \text{ iff } (q_j, c) \in \delta(q_i, a, b)$$

$$a \in \Sigma_\epsilon \quad b, c \in \Gamma_\epsilon$$

- (1) CFG \Rightarrow PDA \quad (2) PDA \Rightarrow CFG
 $= (V, \Sigma, R, S) \quad = (Q, \Sigma, \Gamma, q_0, \delta, F)$

(1) Assume that the CFG is in CNF

If " $S \Rightarrow \epsilon$ " $\in R$, then $q_0 \in F$
 $\Sigma = \Sigma \quad \Gamma = V \cup \Sigma \cup \{ \epsilon, \$ \}$



$$[q_0] 000111 \Rightarrow \$S' [q_{loop}] 000111 \Rightarrow \$R [q_e] 000111$$
~~$$\$X [q_e] 000111$$~~

$$\Rightarrow \$R0 [q_e] 000111$$

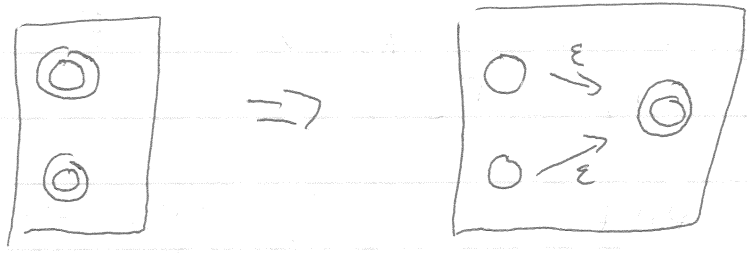
$$\$YRX [q_e] 00111 \leftarrow \$YS [q_e] 00111 \leftarrow \$R [q_e] 00111$$

$$\Rightarrow \$YR0 [q_e] 00111 \Rightarrow \$YR [q_e] 0111$$

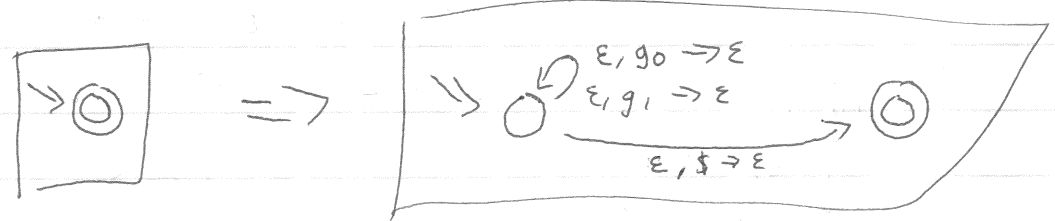
9-4)

Turn a PDA $p = (Q, \Sigma, \Gamma, q_0, \delta, F)$ into a CFG $g = (V, \Sigma, R, S)$

Assume that p has one accept state ($F = \{q_f\}$)



Assume that p empties the stack before accepting



Assume that every transition pushes OR pops

X $a, \epsilon \rightarrow \epsilon$ Y $X \xrightarrow[a, \epsilon \rightarrow c]{[push]} X' \xrightarrow[\epsilon, c \rightarrow \epsilon]{[pop]} Y$
 $a, b \rightarrow \epsilon$ [pop]
 $a, \epsilon \rightarrow c$ [push]

X $a, b \rightarrow c$ Y $X \xrightarrow[a, b \rightarrow \epsilon]{[pop]} X' \xrightarrow[\epsilon, \epsilon \rightarrow c]{[push]} Y$

$$V = \{ A_{p,q} \mid p, q \in Q \}$$

$$A_{p,q} \Rightarrow^* w \quad \text{if} \quad [p]w \Rightarrow^* [q]$$

To be continued...