

"P" =  $\epsilon \in x$

$\exists x \in \text{ALL}, x \notin \text{REG}$

$\forall d \in \text{DFA}, L(d) \neq x$

"P" =  $|x| \in \infty$  infinite ordinal  
i.e.  $\infty$ ; for some

lemma:  $\exists P, \forall x \in \text{REG}, P(x)$

burden:  $\exists x \in \text{ALL}, \neg P(x)$

$\Rightarrow x \notin \text{REG}$

$\exists P, \forall d \in \text{DFA}, P(L(d))$

P: LANG  $\rightarrow$  Prop

$\exists P', \forall d \in \text{DFA}, P'(d)$

P': DFA  $\rightarrow$  Prop

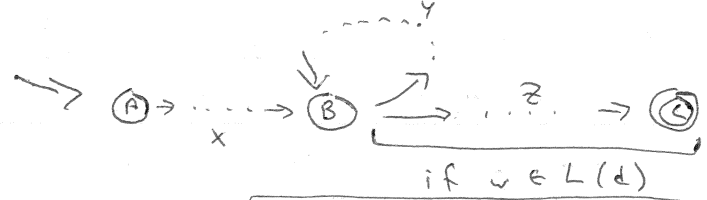
Suppose  $d$  accepts string  $w$  (i.e.  $w \in L(d)$ )

How many states could  $d$  visit during its run?

$[1, |w| + 1]$

Suppose  $d$  had 4 states and  $w = 01100$

$[1, \min(|Q|, |w| + 1)] = [1, 4]$



$x, y, z \in \Sigma^*$

$|x|$  could be 0

$|z|$  could be 0

$A=B \quad B=C$

①  $w = xyz$

If  $w \in L(d)$  exist, ②  $|xy| < |Q|$       ③  $|y| > 0$       = P'

and  $|w| > |Q| \Rightarrow$

④  $xz \in L(d), xyyz \in L(d)$

$\forall i \in \mathbb{N}, xy^i z \in L(d)$

$P = \exists p \in \mathbb{N}, P'$  (replace  $|Q|$  with  $p$ )

## Regular Pumping Property (RPP)

$\forall (A: \text{Language over } \Sigma) \rightarrow \text{Prop} :=$

$\exists p \in \mathbb{N}$ .

$\forall (w \in A \mid |w| > p)$

$\exists (x, y, z \in \Sigma^* \mid w = xyz \wedge |xy| \leq p \wedge |y| > 0)$ .

$\forall i \in \mathbb{N}$ .

$xy^i z \in A$ .

$(\forall A \in \text{REG. RPP}(A)) \leftarrow \text{previous page } (p = |Q|, xyz \text{ were based on DFA})$

$\neg \text{RPP}(A) :=$

$\forall p \in \mathbb{N}$ .

$\exists (w \in A \mid |w| > p)$

$\forall (xyz \in \Sigma^* \mid w = xyz \wedge |xy| \leq p \wedge |y| > 0)$

$\exists i \in \mathbb{N}$ .

$xy^i z \notin A$ .

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$A = 10^*01$  given:  $p$  choose:  $w = 10^p01$

given:  $x, y, z$   $w = xyz$   $|xy| \leq p$   $|y| > 0$

$xy = 10^a$   $z = 0^b01$   $a+b = p$

$1+a \leq p$   $x = 10^a$   $y = 0^d$   $a = c+d$

$d > 0$

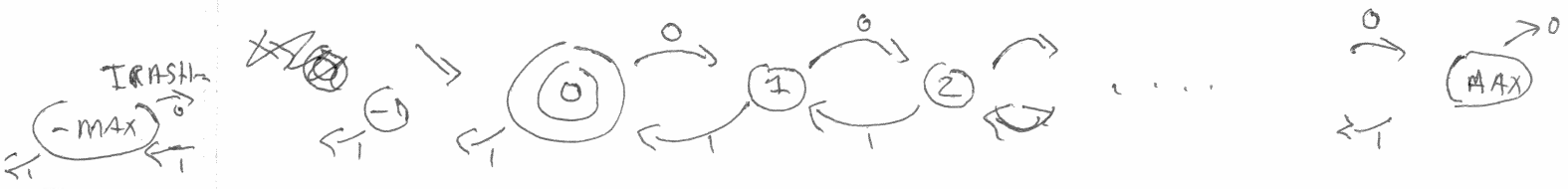
$xy^i z = 10^c 0^{di} 0^b 01$   $c+di+b \in \mathbb{N}$

7-3/

- $( ( ( ) ) ) \in \text{RACKET}$
- $( ( ) ( ) ( ( ) ) ) \in \text{RACKET} = \text{more complicated}$
- $( ) ) \in \text{TRASH} = \text{more complicated}$
- $( ( ) ) \in \text{BABY} = \binom{n}{n}$
- $00 11 \in \text{ASM} = 0^n 1^n$

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int sum = 0
while ((char c = getc()) != EOF)           = TRASH
{
  if (c) { sum++; } else { sum--; }
}
return (sum == 0);
  
```



ASM & REG.

Given:  $p$       Choose:  $w = 0^p 1^p$        $|w| = 2p$

Given:  $x y z$ ,  $w = x y z$ ,  $|x y| \leq p$        $|y| > 0$

$$x = 0^a \quad y = 0^b \quad z = 0^c 1^p$$

$$a + b \leq p \quad a + b + c = p \quad b > 0$$

$$x y^i z = 0^a 0^{bi} 0^c 1^p = 0^{a+bi+c} 1^p$$

$\in \text{ASM}$  if  $a + bi + c = p$

$$bi - b = 0 \quad b(i - 1) = 0$$

$$i = 1$$

choose:  $i \neq 1$

$\Rightarrow$  ASM & REG  
and REG  $\neq$  ALL

*[Faint, illegible handwriting on lined paper]*