

Machine

Language Category

ALL
 \supset
 REG
 \supset
 FIN

DFA $\xrightarrow[n]{L}$ REG (regular)
 (5-tuples) (languages = sets of strings)

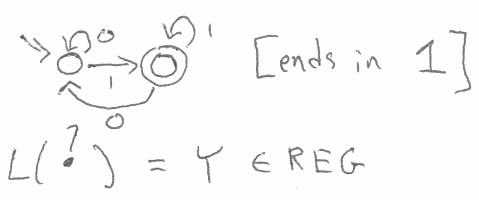
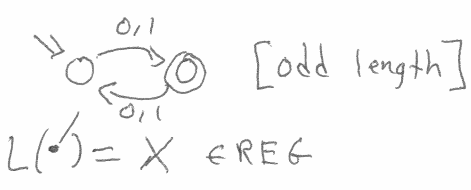
$\xleftarrow{1}$ Find the smallest DFA whose language matches

$$\{a, b\} \cup \{b, c\} = \{a, b, c\}$$

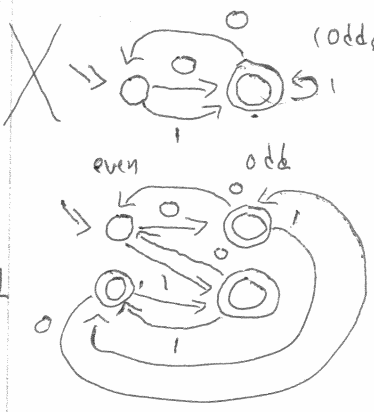
$$\{a, b\} \cup \{b, c\} = \{(0, a), (0, b), (1, b), (1, c)\}$$

Union

$a \in X \cup Y$ iff $a \in X$ or $a \in Y$



$\exists d_3. L(d_3) = X \cup Y?$ If so, then $X \cup Y \in \text{REG}$



REG is closed union
 $\forall X \in \text{REG}, Y \in \text{REG}, \exists A \in \text{REG}, A = X \cup Y$

$X \in \text{REG}$ iff $\exists d_X \in \text{DFA}, L(d_X) = X$

$\forall d_X \in \text{DFA}, d_Y \in \text{DFA}, \exists d_A \in \text{DFA}, L(d_A) = L(d_X) \cup L(d_Y)$

Input: $d_X : \text{DFA} = (Q_X, \Sigma, q_{0X}, \delta_X, F_X)$

$\exists a \in U, P(a)$

$d_Y : \text{DFA} = (Q_Y, \Sigma, q_{0Y}, \delta_Y, F_Y)$

constructive: give a , Prove $P(a)$

Output: $d_A : \text{DFA} = (Q_A, \Sigma, q_{0A}, \delta_A, F_A)$

$\neg \forall a \in U, \neg P(a)$

$Q_A = Q_X \times Q_Y$

$\delta_A((q_X, q_Y), a) =$

$q_{0A} = (q_{0X}, q_{0Y})$

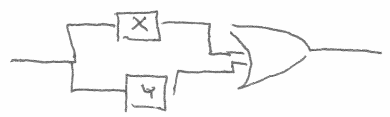
$(\delta_X(q_X, a),$

$\delta_A : Q_A \times \Sigma \rightarrow Q_A = (Q_X \times Q_Y) \times \Sigma \rightarrow (Q_X \times Q_Y)$

$\delta_Y(q_Y, a))$

$F_A = \{ (q_X, q_Y) \mid q_X \in F_X \text{ or } q_Y \in F_Y \}$

$(F_X \times Q_Y) \cup (Q_X \times F_Y)$
 X accepts Y accepts

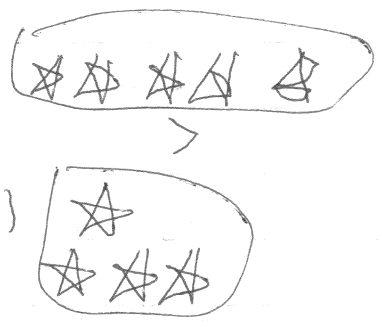


Regular Languages are closed under the Regular Operations

- Union : \cup (we proved)
 - Intersect: \cap (we proved $= F_x \times F_y$)
 - Complement: $\bar{}$ ($F_A = F_x^c$, we proved)
 - Reversal : $\bar{}$
 - Star : $*$
 - Concat : \circ
- } — hard proofs

Road map:

- concat is hard on DFAs
- Define BlahFA
- prove concat on BlahFA (very easy)
- Show that BlahFA = DFA (medium)
- We like BlahFA!
- Use BlahFA to make DFA-PL



Talk about smallest Union