

ALL = Σ_1 — FALSE

$A \in \Sigma_0$ iff $A \in \Sigma_1$ and $\neg A \in \Sigma_1$
 \Rightarrow — obvious ^{normal} \neg negated
 \Leftarrow — not obvious

assume X is the TM for $A \in \Sigma_1$
 Y is the TM for $\neg A \in \Sigma_1$

???

Halting Problem :=

construct Z , the TM for $A \in \Sigma_0$

$U := \text{Atm} \in \Sigma_1 \quad \# \text{Atm} \notin \Sigma_0$

$Z(w) :=$

Assume $\neg \text{Atm} \in \Sigma_1$ is false
 $\therefore \text{Atm} \in \Sigma_0 \Rightarrow \text{False}$

\rightarrow take a step in $X(w)$
 \rightarrow take a step in $Y(w)$
 repeat but if X says $Y \rightarrow Y$

$\neg \text{Atm} = \{ \langle M, w \rangle \mid M \in \text{TM and } w \notin L(M) \}$

Y says $Y \rightarrow N$

When are two sets the same size?

cardinality : set \rightarrow num $|A| = |B|$ iff $\text{card}(A) = \text{card}(B)$

cardinality (0^*) = ? \aleph_0 Georg Cantor

Set A is the same size as Set B ($A \approx B$) iff

exists $f : A \rightarrow B$ such that $\forall b \in B, \exists a \in A, f(a) = b$

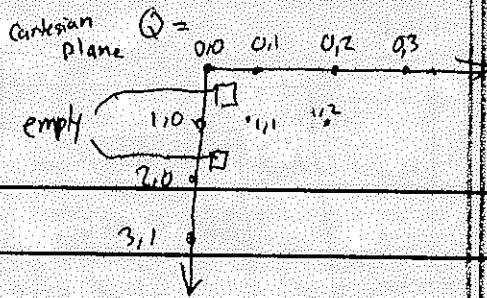
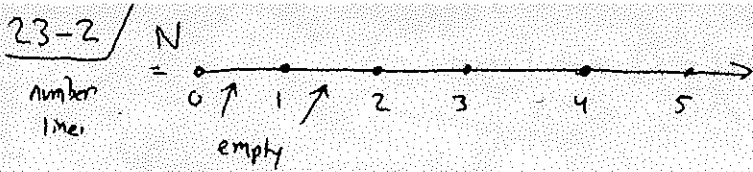
and $\forall a \in A, \exists b \in B, f(a) = b$

$\forall a, a' \in A, \forall b, f(a) = b \wedge f(a') = b \Rightarrow a = a'$

$\text{Nat}(\mathbb{N}) \not\approx \text{Evens}$ (wrong)

$A = \mathbb{N} \quad B = \text{Evens} \quad f(n) = 2 \times n \Rightarrow \mathbb{N} \approx \text{Evens}$

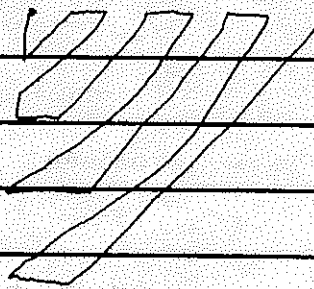
B is countable $\Rightarrow \mathbb{N} \approx B \quad |B| = \aleph_0$



$N \not\cong \mathbb{Q} = N \times N$
wrong!

$N \cong N \times N$

$$f(x, y) = \frac{1}{2}(x+y)(x+y+1) + y$$

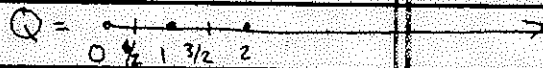
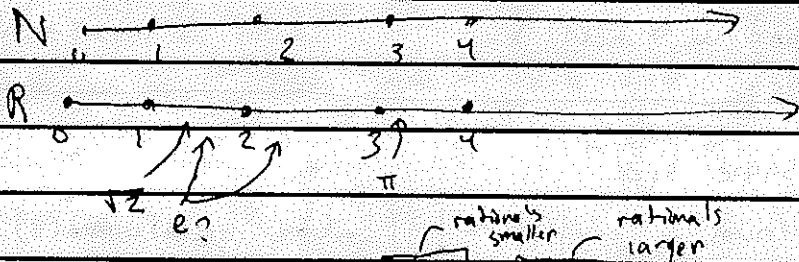


$\mathbb{Q} = N \times N$

$N \cong N \times N \times N$

$$g(x, y, z) = f(x, f(y, z))$$

$N \cong N^k$ for all $k \in \mathbb{N}$



Dedekind Cut $R = P(\mathbb{Q}) \times P(\mathbb{Q})$

Cauchy Sequence $R = N \rightarrow \mathbb{Q}$ s.t. $f(n)$ approaches $l \in R$ as $n \rightarrow \infty$

Infinite Digit Sequences $R = N \rightarrow D$ (where $D = [0, 9]$)

Infinite Binary Sequence $N \rightarrow \{0, 1\}$ $R_{[0, 1)}$ $x \in R_{[0, 1)}$ $x = 0.01010\dots$

$\pi/4 =$ some fun

$$f(n) = \begin{cases} 1 & (n \equiv 1) \pmod{4} \\ 0 & \text{otherwise} \end{cases}$$

$N \not\cong R_{[0, 1)}$ (IDS)

Assume $N \cong \text{IDS}$, $\Rightarrow f(n: N) : \text{IDS}$ (and f is onto / one-to-one)

$\forall b \in \text{IDS}, \exists a \in N, f(a) = b$ - onto

$$\neg (\forall b \in D, \exists a \in N, f(a) = b) = \exists b \in \text{IDS}, \forall a \in N, \boxed{f(a) \neq b} = \exists i \in N, f(a)(i) \neq b(i)$$

$$b(i) := \neg f(i)(i) \Rightarrow N \not\cong \text{IDS} \Rightarrow N \not\cong R$$

Diagonalization Argument

3-3/

ALL
 \approx
 ISS
 \approx
 R

\neq

Σ_1

\approx
 TM

← every lang in Σ_1
 is some TM (by defn)

\approx

← by encoding 7-tuple

N^k for some k

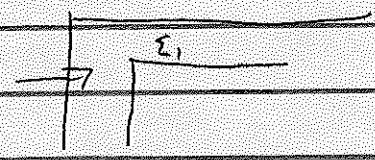
\neq

\approx

N

ALL

infinite
 shell



$ALL = P(\Sigma^*)$

$\Sigma^* = \epsilon, 0, 1, 00, 01, 10, 11, 000, 001, 010, 011, 100, 101, \dots$

$\Sigma^* \setminus 1 = A \in ALL = \epsilon, 1, 01, 11, 001, 011, 100$

$A(i) = \underset{\substack{\underbrace{}_0 \\ \underbrace{}_1 \\ \underbrace{}_2 \\ \underbrace{}_3 \\ \underbrace{}_4 \\ \underbrace{}_5 \\ \underbrace{}_6}}{1 \ 0 \ 1 \ 0 \ 1 \ 0 \ 1 \ 0 \ 1 \ 0 \ 1 \ 0 \ 1 \ \dots}$

