

Subtyping

$$D \rightarrow R$$

we want it to work on things

"like" D but not exactly D

$$\text{nat} \rightarrow \text{nat}$$

$$(\text{nat} \rightarrow \text{nat}) \rightarrow \text{nat}$$

$$\Leftarrow \lambda f. \text{S}$$

: $\forall \alpha, \alpha \rightarrow \text{num}$

$$\wedge f. f(\text{S})$$

most general polymorphic type (principal types)

type inference provably discovers $\lambda f. \text{S}$

$$(\text{nat} \times \text{nat}) \rightarrow \text{nat}$$

$$\Leftarrow \lambda p. (\text{fst } p) + (\text{snd } p)$$

$$\forall \alpha. (\text{nat} \times \alpha) \rightarrow \text{nat}$$

$$\Leftarrow \lambda p. \text{S} + (\text{fst } p)$$

$$\text{Cat} = \text{"Animal"} \times \text{CatStuff}$$

$$\text{Animal} = \text{"Animal"} \times \text{Nothing}$$

"Duck Typing" \approx "A type is what you can do to it"

~~in Python~~ in Lua/Ruby, everything is a dictionary, capabilities

= keys in dictionary

Subtyping so far =

Structural Subtyping

Objects have fields (methods are function-valued fields)

One object is compatible with another if shared fields are compatible

Obj $\mathbb{O} = \{a, b, c, d, e, f\}$

Interface $X = \{a, b, c\}$

$Y = \{c, d, e\}$

$Z = \{d, e, f\}$

$\mathbb{O} \leq X$

Y

Z

OCaml, Haskell, Typed Racket, ie egghead academic languages

Java = Nominal Sub-typing

An object implements interfaces (fields) (method sets)

One object \mathbb{O} is compatible with interface I if \mathbb{O} implements I

$\mathbb{O} = \{f, \dots\} \times \text{List Int}$

$\text{Int} = \{m, \dots\}$

$\mathbb{O} = \{a, b, c, d, e, f\} \times [X, Z]$

$\mathbb{O} \leq X$

$\mathbb{O} \leq Z$

$\mathbb{O} \not\leq Y$