

20-2/ Suppose I had a Turing decider and
I wanted an enumerator...

enumerator := try all strings in lexicographic
and if the decider says yes,
then it prints

recognizer \Rightarrow enumerator := try all strings "in parallel"

enumerator \Rightarrow recognizer := check the output...
never say no

enumerator "in order" \Rightarrow decider := check the output...
say no, if something is generated

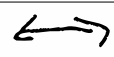
20-3/

DFA's



NFA's

(more expressive)



REG

TM's



MT ↔ U, N



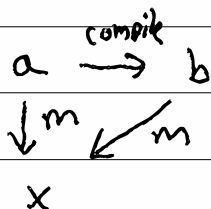
Enumerators



NTM ↔ ∅

20-4 want: an extension of Turing machines

$\forall a \text{ input. } \exists b \text{ output. } M(a) = M(b)$

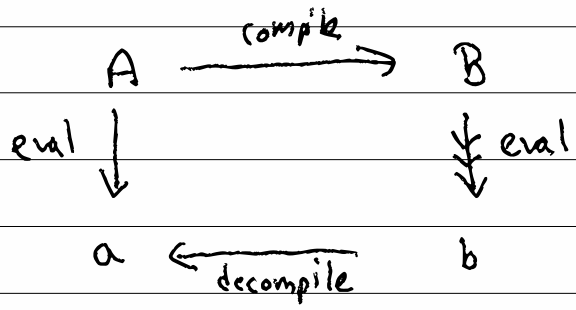


Galois connection on a bisimulation

20-5/

$\forall a \in \text{input}, \exists b \in \text{output}.$

$\text{eval}_i(a)$ similar $\text{eval}_o(b)$



20-b TMs w/ a "stay" action

$$\delta: Q \times \Gamma \rightarrow Q \times \Gamma \times \{L, R\}$$

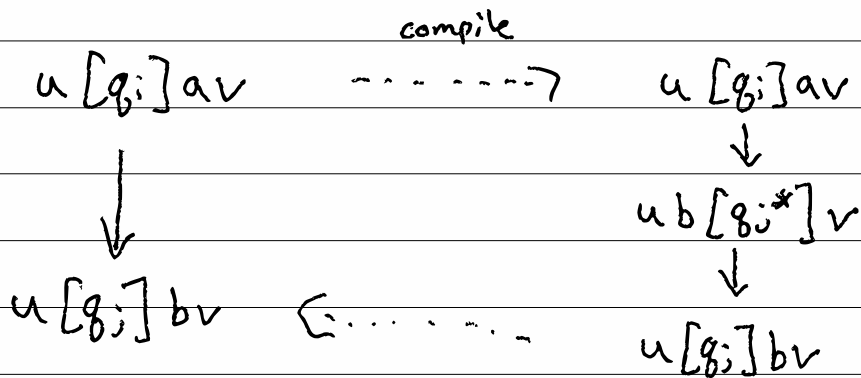
$$\delta_s: Q \times \Gamma \rightarrow Q \times \Gamma \times \{L, R, S\}$$

$$\delta(q_i, a) = (q_j, b, s)$$

$$u[q_i]av \Rightarrow u[q_j]bv$$

20-7 / stay = input

→ normal tm



$$\delta(q_i, a) = (q_i, b, S)$$

$$\Rightarrow \delta(q_i, a) = (q_i^*, b, R)$$

$$\forall q_i \in Q, \delta(q_i^*, \sigma) = (q_i, \sigma, L)$$

$\sigma \in \Gamma.$