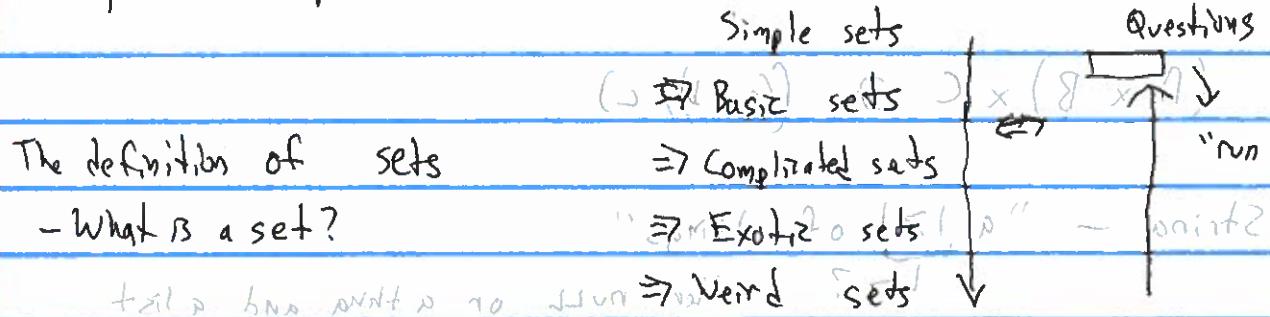


Foundation of Comp Sci

S-1

Theory of Computation



$\{a, b, c\}$ or $\{a, b, c\}$ \checkmark most sets not \emptyset

$\{b, a, c\}$	$\begin{array}{ c c } \hline b & \checkmark \\ \hline c & \checkmark \\ \hline a & \times \\ \hline e & \times \\ \hline f & \times \\ \hline \end{array}$	$(U \cup V)$ - the universe
		$(\exists x \in U, \forall y, y \in x \Rightarrow x = y)$

$\{\emptyset, \{a\}, \{\emptyset, a\}\}$ to discuss | extension to sets in part A
too long/0

Finite sets - can be written as list of elements

\emptyset - empty set

$$\forall x, x \notin \emptyset \quad \boxed{x \notin \emptyset}$$

$a \in \{a, b, c\} \checkmark$

$a \in \emptyset \quad \times$

not in

\in - set membership

\cup - union operation : Set \times Set \rightarrow Set

$\forall x, x \in A \cup B$ iff $x \in A$ or $x \in B$

\cap - intersection : $\forall x, x \in A \cap B$ iff $x \in A$ and $x \in B$

\subset - subset : Set \times Set \rightarrow Prop. $A \subset B$ iff $\forall x, x \in A \Rightarrow x \in B$

P - power set ($P(A)$, $P(A)$, 2^A) $P(\emptyset, \{a\}) =$

$\forall x, x \in P(A)$ iff $x \subset A$ $\Rightarrow \{x\}, \{\{a\}\}, \{\emptyset, \{a\}\}, \{\emptyset, \emptyset\}$

$P(\emptyset) = \{\emptyset\}$ $P(P(\emptyset)) = \{\emptyset, \{\emptyset\}\}$ $P^n(\emptyset) = \text{a set with } n \text{ things}$

relations and functions mult b by cartesian product

$f : X \rightarrow Y$ $f \subseteq X \times Y$ $x \in A \times B$ iff $\{x\} = (a, b)$

dom range

$f(x) = y \Leftrightarrow (x, y) \in f$

and $a \in A$

$r : X \times Y \rightarrow r \subseteq X \times Y$

$\forall x, y_1, y_2$

and $b \in B$

$f(x) = y \wedge f(x) = y_2$

$\Rightarrow y_1 = y_2$

1-2) Pairs like $(a, b) \in A \times B$ (≥ 2 no) \Rightarrow NS, $\{ab, ba\}$

are a special case of tuples, 2-tuples (no) \Rightarrow pairs

$$(A \times B) \times C \ni ((a, b), c)$$

String - "a list of things"

\hookrightarrow ? either null or a thing and a list

a function from N (natural numbers) to things

abc $\xrightarrow{\text{separating}}$ (and a N)

$$:= (3, 0 \mapsto a, 1 \mapsto b, 2 \mapsto c)$$

$\begin{array}{ c c } \hline d & \\ \hline \end{array}$	$\{a, b, c\}$
$\begin{array}{ c c } \hline d & \\ \hline \end{array}$	$\{a, b, c\}$
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$\begin{array}{ c c } \hline d & \\ \hline \end{array}$	$\{a, b, c\}$
$\begin{array}{ c c } \hline d & \\ \hline \end{array}$	$\{a, b, c\}$

\nwarrow sigma

A string is a list of characters (elements of set Σ)
= alphabet

ϵ - lower case greek epsilon = the string of length 0, for $\epsilon \in \Sigma$

x_i - the i th character in x for strings - ϵ

$|x|$ - the length of string x

$x \cdot y, xy$ - the concatenation

$(xy)_i = x_i$ if $i < |x| + |y|$: notations now - ϵ

y_j where $j = i - |x|$ $\forall x, y \in \Sigma^*$

lexicographic ordering of strings Σ^* $\Sigma = \{0, 1\}$ natural - ϵ

$0 \in \Sigma, 01, 00, 010, 011, 000, 001, 0100, 0110, \dots$ factue - ϵ

$= (\Sigma^*, \leq)$ $(\Sigma^*, (A)^*, (A)^*)$ for revs - ϵ

language of an alphabet is a set of strings Σ^*

definition $\{0, 1\}^* (\emptyset)^* \{0, 1\}^* = ((\emptyset)^*)^* \{0, 1\}^* = (\emptyset)^*$

"All binary strings that represent odd numbers"

$\{0, 1\} \cup \{10, 11\}^* = \{01, 01, 11\}^*$

$Y \subseteq X : ?$

$X \cup Y \cap \text{odd}(A) = ?$ "odd numbers"?

$\Sigma = \{0, 1\}$

$Y \times X \subseteq \Sigma^* \quad Y \times X \subseteq \Sigma^*$

$SX = (x) \wedge Y = (x)^2$

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