

3-4 / typedef enum { XL, YM, YW, ZW } state\_t;

```
int machine ( ) {
```

```
    state_t st = XL;
```

```
    while ( char c = getc ( ) ) {
```

```
        switch (st) {
```

```
            case XL: switch (c) {
```

```
                case '0': st = YM; break;
```

```
                case '1': st = ZW; }
```

```
            case YM: st = YW; break;
```

```
            case YW: st = YW; break;
```

```
            case ZW: st = ZW; break; } }
```

```
    return st == YW;
```

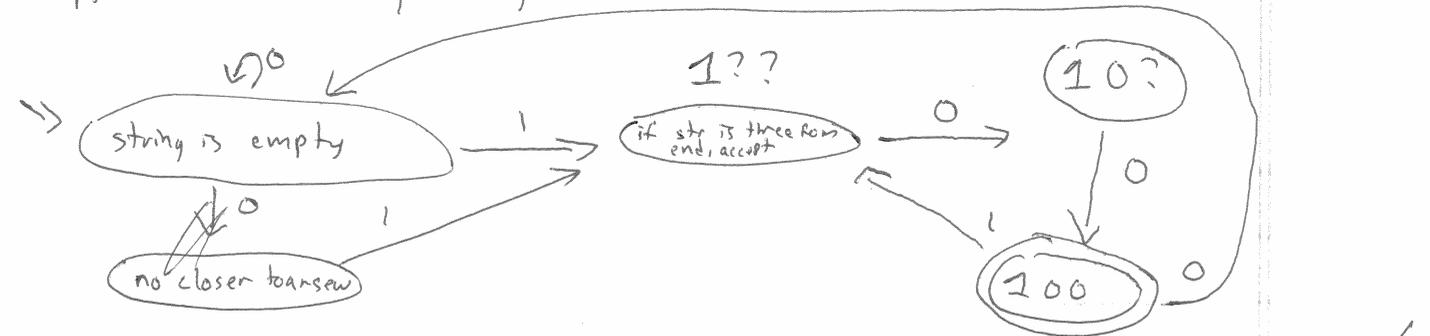
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Regexp :  $\underbrace{\text{courses}}_x / \underbrace{*}_{\text{any}^*} / \underbrace{\text{grades}}_y / \underbrace{*}_{\text{any}^*} \underbrace{\text{fail}}_z$

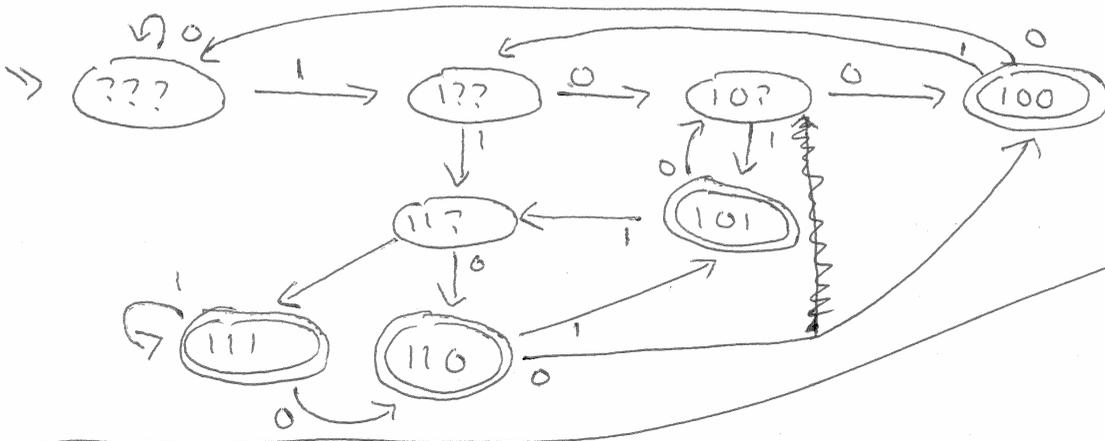
re =  $x \circ \Sigma^* \circ y \circ \Sigma^* \circ z$

1/0

1-1/ A = { all binary strings where third to-last character is 1 }



0100 ✓  
1000 X  
11100 ✓



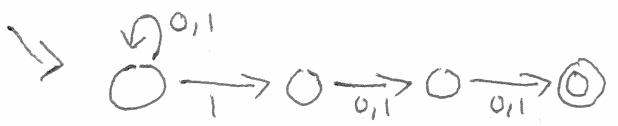
NFA -

Non-deterministic Finite Automata

diff 1: states don't have all transitions

diff 2: states may have > 1

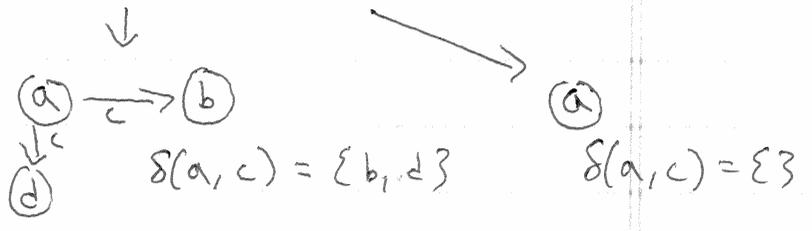
states can have any number of trans. per character



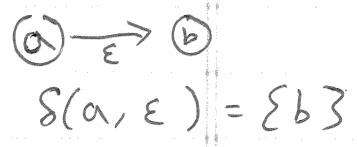
Feature telling  
backtracking  
for k()-ing  
mystery good one



DFA  $\delta = \langle \Sigma, Q, q_0 \in Q, \delta: Q \times \Sigma \rightarrow Q, F \subset Q \rangle$   
 NFA  $\delta = \langle \Sigma, Q, q_0 \in Q, \delta: Q \times \Sigma \rightarrow P(Q), F \subset Q \rangle$



$\Sigma_\epsilon = \Sigma \cup \{ \epsilon \}$   
epsilon



4-2/

$$L(\text{NFA } n) = \{ x \mid x \text{ is accepted by NFA } n \}$$

A string  $x$  is accepted by NFA  $n$  iff

$$q_0 \xRightarrow{x}^* q_i \text{ s.t. } q_i \in F \quad \text{[exactly the same as DFA]}$$

An NFA  $n$  runs from  $q_i$  to  $q_j$  on  $x$  ( $q_i \xRightarrow{x}^* q_j$ )

$$q_i \xRightarrow{\epsilon}^* q_i \quad q_i \xRightarrow{ax}^* q_k \text{ iff } q_i \xRightarrow{a} q_j \text{ where } q_j \in Q$$

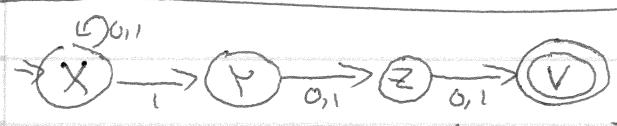
$$q_j \xRightarrow{x}^* q_k \quad a \in \Sigma_\epsilon$$

only difference from DFA

An NFA  $n$  steps from  $q_i$  to  $q_j$  on  $a$  ( $q_i \xrightarrow{a} q_j$ )

$\Sigma_\epsilon$  (for DFA =  $\Sigma$ )

$$q_j \in \delta(q_i, a) = \text{for DFA}$$



$$1 \Rightarrow [Y]100 \Rightarrow [Z]00 \Rightarrow [V]00$$

~~$$[X]11100 \xRightarrow{1} [X]1100 \Rightarrow [X]100 \Rightarrow [X]00 \Rightarrow [X]0 \Rightarrow [X]_0$$

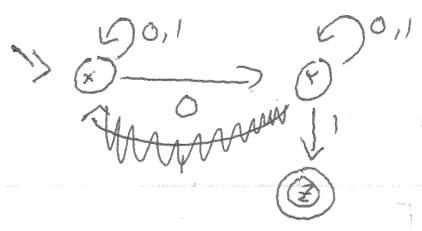
$$\Rightarrow [Y]1100 \Rightarrow [Z]100 \Rightarrow [V]00$$~~

$$Y \in \delta(X, 1) = \{X, Y\}$$

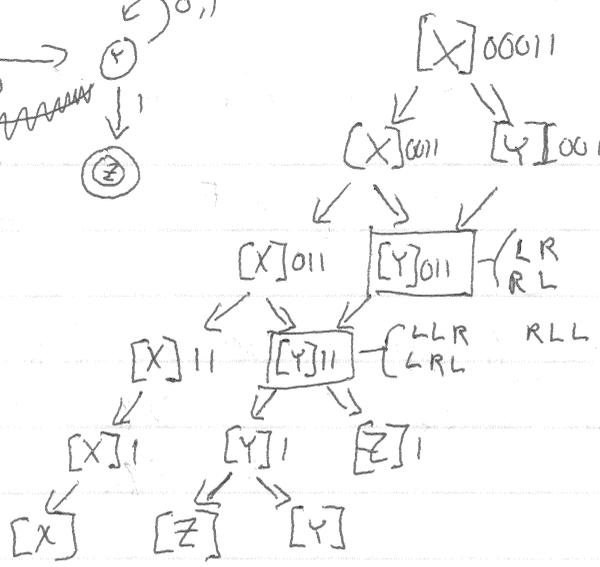
$$X \xRightarrow{11100} V \quad ? \checkmark \quad 11100 \in L(n)$$

back-tracking = DFS      oracle = cheating  
 funk()-ing = BFS

1-3/



000 11



4

