

$$A_{TM} = \{ \langle M, w \rangle \mid M \in TM \text{ and } M \text{ accepts } w \}$$

$$\langle \text{binary add}, 0110 + 1000 = 1110 \rangle \in A_{TM}$$

$$\langle \text{palindrome}, 0110 \rangle \notin A_{TM}$$

$$\langle \text{palindrome}, 1000 \rangle \notin A_{TM}$$

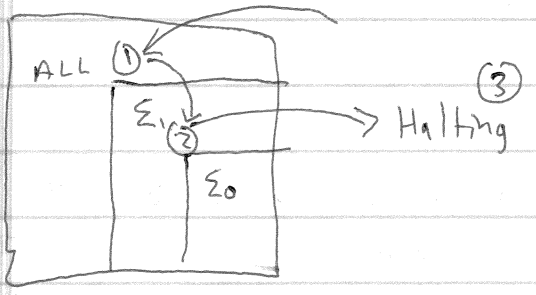
- $\Sigma_1$ : On input  $\langle M, w \rangle$ ,
  - use tape 2 as  $w$
  - use tape 3 for  $q$  (initialized  $q_0$ )
  - use tape 4 for  $\delta$
  - simulate the machine

$\Sigma_0$  + + + if  $M$  diverges on  $w$ , then reject  
 The Halting Problem

$X \notin \Sigma_0$  where  $X$  is the Halting Problem

"Turing Pumping Property"

$$X \notin REG \iff \neg RPP(X) \\ + \forall Y \in REG, RPP(Y)$$



$\exists x, x \notin ALL \wedge x \in \Sigma_1$   
 $\Sigma_1 \subseteq ALL$   
 Prove that ALL is bigger

$$|\{x, y, z\}| = 3 \\ |\{x, y\}| = 2$$

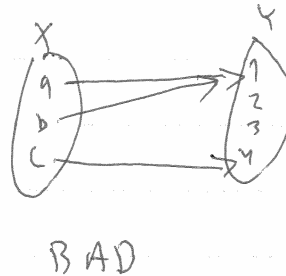
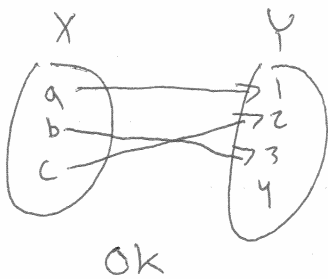
$|\cdot|$ : set  $\rightarrow$  natural  
 cardinality function

matching elements  $\Rightarrow$  same size  
 non-matching elements  $\Rightarrow$  diff (ie one is bigger)

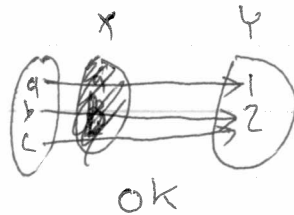
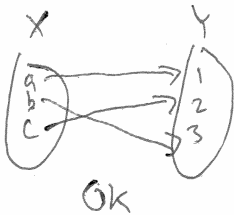
$$(X \leq Y) \wedge (X > Y \vee Y > X) \quad X = \Sigma_1 \\ \Rightarrow Y > X \quad Y = ALL$$

23-2 / Set X and Set Y have the same size if there exists a one-to-one and onto function,  $m$ , from X to Y.

> one-to-one:  $\forall a, b \in X, m(a) \neq m(b) \text{ if } a \neq b$   
 $m(a) = m(b) \Rightarrow a = b$



> onto:  $\forall b \in Y, \exists a \in X, m(a) = b$



The set of naturals  $N$  is the same size as the set of Even numbers,  $E$

$$N = 0, 1, 2, 3, 4, 5, 6, 7, \dots$$

$$E = 0, 2, 4, 6, 8, 10, 12, 14, \dots$$

$$m(n) = 2 \times n$$

$$\textcircled{1} m(a) = m(b) \Rightarrow a = b$$

$$2a = 2b \Rightarrow a = b$$

$$a = b \Rightarrow a = b$$

$$\textcircled{2} \forall b \in E, \exists a \in N, m(a) = b$$

$$\exists a \in N, m(a) = 2 \times n$$

$$\exists a, 2 \times a = 2 \times n$$

$$a = n$$

$\aleph_0$  or "countable"  
 or "same size as naturals"

$$N \times \{a, b\} = \{ (0, a) (0, b) (1, a) (1, b) \dots \}$$

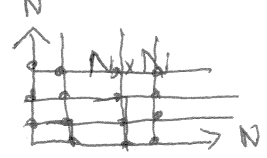
$$m(n, a \vee b) = 2n \quad \text{if } a$$

$$2n + 1 \quad \text{if } b$$

$\mathbb{Z}$  are countable

$N \times F$  where  $F$  is finite, then it is countable

$$\mathbb{N} \times \mathbb{N} = \{ (0,0), (1,1), (0,1), (2,3), \dots \}$$



$\mathbb{N} \times \mathbb{N}$  is the same size as  $\mathbb{N}$   
 plane line

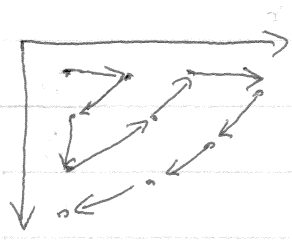
$$m((a,b)) = m((x,y)) = \text{one number}$$

$$m(x,0) = x \quad \text{BAD} \quad \text{not 1-1}$$

$$m(0,y) = y \quad \text{BAD}$$

$$m(x,y) = x+y \quad \text{BAD}$$

$$m(x,y) = 2^x \cdot 3^y \quad \text{BAD} \quad \text{not onto}$$



$$m(x,y) = \frac{1}{2}(x+y)(x+y+1) + y$$

Q (~~the~~ rationals) are countable

$$\mathbb{N} \times \mathbb{N} \cong \mathbb{N}$$

$$A \cong \mathbb{N} \wedge B \cong \mathbb{N}$$

$$\mathbb{N} \times (\mathbb{N} \times \mathbb{N}) \cong \mathbb{N} \times \mathbb{N} \cong \mathbb{N}$$

$$\Rightarrow A \cup B \cong \mathbb{N}$$

$$\mathbb{N}^k \cong \mathbb{N}$$

The set of Turing machines over  $\Sigma = \{0,1\}$ ,  $\Gamma = \{0,1\}$

$$\langle Q, \Sigma, \Gamma, q_0, \delta : Q \times \Gamma \rightarrow Q \times \Gamma \times \{L,R\}, q_a, q_r \rangle$$

$$Q \quad |Q| \quad (Q \times \Gamma) \times (Q \times \Gamma \times \{L,R\}) \quad |Q| \quad |Q|$$

$$n = |Q| \quad n \times (n \times 3) \times (n \times 3 \times 2) \times n \times n = 18 \times n^5$$

( $\mathbb{N} \rightarrow$  Finite set)

$$\text{TMs} \cong \mathbb{N}$$

$$0 \rightarrow \emptyset$$

$$\Sigma_0 \cong \mathbb{N}$$

$$1 \rightarrow \{0\}$$

$$\Sigma_1 \cong \mathbb{N}$$

23-4) An infinite binary sequence,  $B$ , is a function from  $\mathbb{N}$  to 0 or 1.

$$B_0(n) = 0$$

000000...

$$B_1(n) = 1$$

1111...

$$B_{10}(n) = 1 \text{ if } n \text{ is odd}$$

0, if  $n$  is even

0101010101...

( $X \preceq Y$  then  $Y \preceq X$ )

Let's assume that  $B \preceq \mathbb{N}$ .

$\exists m: B \rightarrow \mathbb{N}$  where  $m$  is 1-1 and onto

$$\textcircled{1} \forall a, b \in B. m(a) = m(b) \Rightarrow a = b$$

$$\textcircled{2} \forall n \in \mathbb{N}. \exists a. m(a) = n$$

$$\neg \textcircled{2} = \exists n \in \mathbb{N}. \forall a. m(a) \neq n$$

$$\neg \textcircled{1} = \exists a, b \in B. m(a) = m(b) \Rightarrow a \neq b$$

$$m(B_0) = 12 \quad m(B_1) = 13$$

$$m(B_?) = 0 \quad m(B_{??}) = 1$$

$\mathbb{N}$	$B$
0	01110001110...
1	110011000000...
2	00 <u>0</u> 1 <u>1</u> 1 <u>1</u> 1 <u>4</u>
3	001 <u>1</u> ~
4	~
5	~

Assume  $\mathbb{N} \preceq B$

$$m(12) = B_0$$

$\exists m: \mathbb{N} \rightarrow B$

$$m(13) = B_1$$

$$\textcircled{1} \forall a, b \in \mathbb{N}. m(a) = m(b) \Rightarrow a = b$$

$$\textcircled{2} \forall n \in B. \exists a. m(a) = n$$

$$\neg \textcircled{2} \exists n \in B. \forall a \in \mathbb{N}. m(a) \neq n$$

$n =$  the missing binary string  $= (\text{fun } i \Rightarrow \neg m(i)(i))$

$m(i) =$  the  $i$ th binary string

$m(i)(i) =$  the  $i$ th digit of the  $i$ th binary string

$$= 0101 \dots$$

$$\neg(m(i)(i)) = 1010 \dots$$

$$\forall a \in \mathbb{N}. m(a) \neq (\text{fun } i \Rightarrow \neg m(i)(i))$$

$$\text{take any } a \vdash m(a) \neq (\text{fun } i \Rightarrow \neg m(i)(i))$$

~~$m(a)(a) \neq$~~

$$m(a)(a) \neq (\text{fun } i \Rightarrow \neg m(i)(i))(a)$$

$$m(a)(a) \neq \neg m(a)(a)$$

$$0 \neq \neg 0$$

$$1 \neq \neg 1$$