

I-1

Theory of Computation

- What are computers? What can they do?
What can't they do?

Set Theory

$$\{1, 2, 3\}$$

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$$X \subseteq Y \text{ iff } \forall a, a \in X \rightarrow a \in Y$$

is a subset

\mathbb{N} - naturals

\mathbb{Z} - integers

$\mathbb{N} \subseteq \mathbb{Z}$

\mathbb{R} - reals

\mathbb{Q} - rationals

$$a \in X \cup Y \text{ iff } a \in X \text{ or } a \in Y$$

union

$$a \in X \cap Y \text{ iff } a \in X \text{ and } a \in Y$$

intersect

$$a \in X^c = \bar{X} \text{ iff } a \notin X \text{ and } a \in U$$

complement

not in

Tuples

(0, 1) is a two-tuple (or pair) of 0 and 1

X is a tuple, then $\pi_i(\bar{x})$ is the i th thing

$$\pi_0(0, 1) = 0$$

$$X \times Y = \{ (x, y) \mid x \in X \text{ and } y \in Y \}$$

cross product

such that

Relation R on X, Y , and Z is a subset of $X \times Y \times Z$

pluses 0 0 0

pluses : $N \times N \times N$ $\forall n.$ pluses 0 n n

$\forall m, \forall n, \forall x. (m+n) + x = m + (n+x) \Leftarrow \exists x. \text{pluses } m n x$

-3) A function F from $X \rightarrow Y$
 \Rightarrow a relation on $X \times Y$
 s.t.

$$\forall x, y_1, y_2, F x y_1 \wedge F x y_2 \Rightarrow y_1 = y_2$$

succ: { (0, 1), (1, 2), (2, 3), (99, 100) } 3

Reflexive: $\forall x, R x x$

Symmetric: $\forall x, y, R x y \Rightarrow R y x$

Transitive: $\forall x, y, z, R x y \wedge R y z \Rightarrow R x z$

Powerset: $P(X) = P(X) = 2^X$

$a \in P(X)$ iff $a \subseteq X$

$X = \{0, 1, 3\}$ $P(X) = \{ \emptyset, \{3\}, \{0\}, \{1\}, \{0, 1, 3\} \}$

Sequence of X is a function from $N \rightarrow X$

(a string if there's also a length N) $(N, N \rightarrow X)$

A string of Σ is a finite sequence of Σ

sigma, the alphabet $|x| = \text{length}$

0110

Jay

ϵ - epsilon, empty string

x^R - reverse, switch direction $f(x) = f(|str| - x)$

$x \circ y = xy$, concatenate

$f(i) = x(i)$ if $i < |x|$
 $y(i - |x|)$ o.w.

Strings can be ordered

$x^n = n$ copies of x

(lexicographically)

x^* = kleene star (not a string) a set of strings

$s = x^n$ of some n , then $s \in x^*$

Devangini

"Language of Σ " is a set of strings of Σ

1-3/

A computer "solves" a problem

A problem is a language

The addition problem is a set strings over $\Sigma = \{0, 1, +, 3\}$

$$= + + 3 \quad 3 + 3 = 6 \quad 3 + 3 = 18$$



"is just the ones that are true"

The even problem is the set of even-length strings over Σ

The C-factorial is the set of all C-programs that compute
the factorial of 25.

The job of the computer is "recognition"

given a string x

return Yes if it is in the set

No if it is NOT in the set

X A computer is a function from string $\rightarrow \Sigma$ Yes, No?

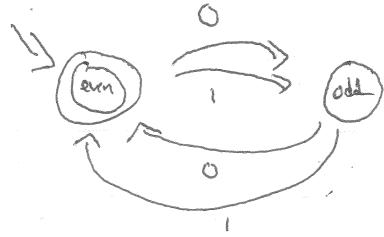
↳ not actually

A computer is a finite tuple

A computer semantics is a finite function $, x \Rightarrow \Sigma^*, N^3$

2-1 Deterministic Finite Automata

Even-ness:
over binary



$$= \{ s \mid s \in \{0, 1\}^* \text{ and } |s| = 2n \text{ for some natural } n \}$$

- is a node or a state
- ↳ ○ is the start state
- $a_0 \xrightarrow{c} a_b$ is a transition that goes FROM a TO b ON c
- is the accept state

$$\begin{array}{ccc} 2n & \rightarrow & 2n+1 \\ \text{even} & & \text{odd} \end{array} \rightarrow \begin{array}{ccc} 2n+2 & = & \cancel{2(n+1)} \\ \text{even} & & \end{array}$$

$$a_0 \xrightarrow{c,d} a_b \text{ (on } c \text{ and)}$$

$$a_0 \xrightarrow{\text{anything}} a_b \text{ (on anything)}$$

Machine with output

$$c \in \Sigma'$$

[c]

when the machine reaches this state, outputs c

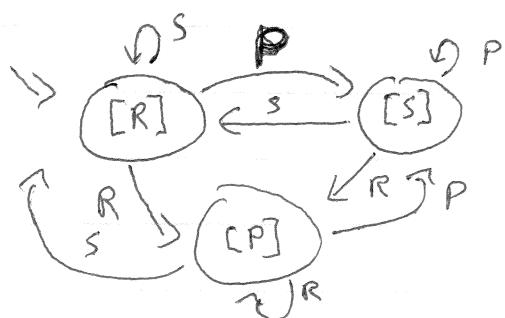
even[c]

Moore machine

Mealy machine

$$a[c]$$

outputs on transition



~~clocked~~ counter machine
negative float

