

Theory of Computation

- What are computers? What can they do?
What can't they do?

Set Theory

$$\{1, 2, 3\} \quad \{1, 1, 2, 3\}$$

$$X \subseteq Y \quad \text{iff} \quad \forall a, a \in X \rightarrow a \in Y$$

\Rightarrow
is a subset

N - naturals Z - integers $N \subseteq Z$
 R - reals Q - rationals

$$a \in X \cup Y \quad \text{iff} \quad a \in X \quad \text{or} \quad a \in Y$$

\Rightarrow
union

$$a \in X \cap Y \quad \text{iff} \quad a \in X \quad \text{and} \quad a \in Y$$

\Rightarrow
intersect

$$a \in X^c = \bar{X} \quad \text{iff} \quad a \notin X \quad \text{and} \quad a \in U$$

\Rightarrow
complement not in

Tuples

$(0, 1)$ is a two-tuple (or pair) of 0 and 1

X is a tuple, then $\pi_i(\bar{i})$ is the i th thing

$$\pi_0(0, 1) = 0$$

$$X \times Y = \{ (x, y) \mid x \in X \text{ and } y \in Y \}$$

\Rightarrow
cross product such that

Relation R on $X, Y, \text{ and } Z$ is a subset of $X \times Y \times Z$

pluses 0 0 0

pluses : $N \times N \times N$ $\forall n$, pluses 0 n n

$\forall n, m$ pluses $(1+n) m (1+x) \Leftarrow \exists x$, pluses $n m x$

-2] A function F from $X \rightarrow Y$
 \Rightarrow a relation on $X \times Y$
 s.t.

succ: $\{ (0,1)$
 $(1,2)$
 $(3,4)$
 $(99,100)$
 \dots $\}$

$$\forall x, y_1, y_2, \quad F x y_1 \wedge F x y_2 \Rightarrow y_1 = y_2$$

Reflexive: $\forall x. R x x$

Symmetric: $\forall x, y. R x y \Rightarrow R y x$

Transitive: $\forall x, y, z. R x y \wedge R y z \Rightarrow R x z$

Powerset: $\mathcal{P}(X) = P(X) = 2^X$

$a \in P(X)$ iff $a \subseteq X$

$X = \{0,1\}$ $P(X) = \{ \emptyset, \{0\}, \{1\}, \{0,1\} \}$
 $\begin{matrix} \emptyset & \{0\} & \{1\} & \{0,1\} \\ 00 & 10 & 01 & 11 \end{matrix}$

Sequence of X is a function from $\mathbb{N} \rightarrow X$

(a string of X here's also a length N) $(\mathbb{N}, \mathbb{N} \rightarrow X)$

A string of Σ is a finite sequence of Σ
 Σ the alphabet $|x| = \text{length}$

ϵ - epsilon, empty string

x^R - reverse, switch direction $f(x) = f(|\text{str}| - x)$

$x \circ y = xy$, concatenate

$$f(i) = \begin{cases} x(i) & \text{if } i < |x| \\ y(i - |x|) & \text{o.w.} \end{cases}$$

strings can be ordered (lexicographically)

~~$x^n = n$ copies of x~~

x^* = kleene star (not a string) a set of strings

$s = x^n$ of some n , then $s \in x^*$

Derivation

"Language of Σ " is a set of strings of Σ

1-3/

A computer "solves" a problem

A problem is a language

The addition problem is a set strings over $\Sigma = \{0-9, +\}^*$

$2+2=3$ $3+3=6$ $3+3=18$



"is just the ones that are true"

The even problem is the set of even-length strings over Σ

The C-factorial is the set of all C-programs that compute the factorial of 25.

The job of the computer is "recognition"

given a string x

return Yes if in the set

No if NOT in the set

X A computer is a function from string $\rightarrow \{Yes, No\}$
↳ not actually

A computer is a finite tuple

A computer semantics is a finite function $x \Rightarrow \{Y, N\}$

2-1 / Deterministic Finite Automata

Even-ness over binary :



$$= \{ s \mid s \in \{0,1\}^* \text{ and } |s| = 2n \text{ for some natural } n \}$$

- is a node or a state
- ⇒ ○ is the start state
- $a \circ \xrightarrow{c} \circ_b$ is a transition that goes FROM a TO b ON c
- ⊙ is the accept state

$$\begin{matrix} 2n & \rightarrow & 2n+1 & \rightarrow & 2n+2 & = & 2(n+1) \\ \text{even} & & \text{odd} & & \text{even} & & \end{matrix}$$

$$a \circ \xrightarrow{c/d} \circ_b \text{ (on } c \text{ and } d) \qquad a \circ \rightarrow \circ_b \text{ (on anything)}$$

Machine with output

$$c \in \Sigma'$$

Ⓚ when the machine reaches this state, outputs c

even[0] Moore machine

Mealy machine $a[c] \rightarrow$ outputs on transition

