

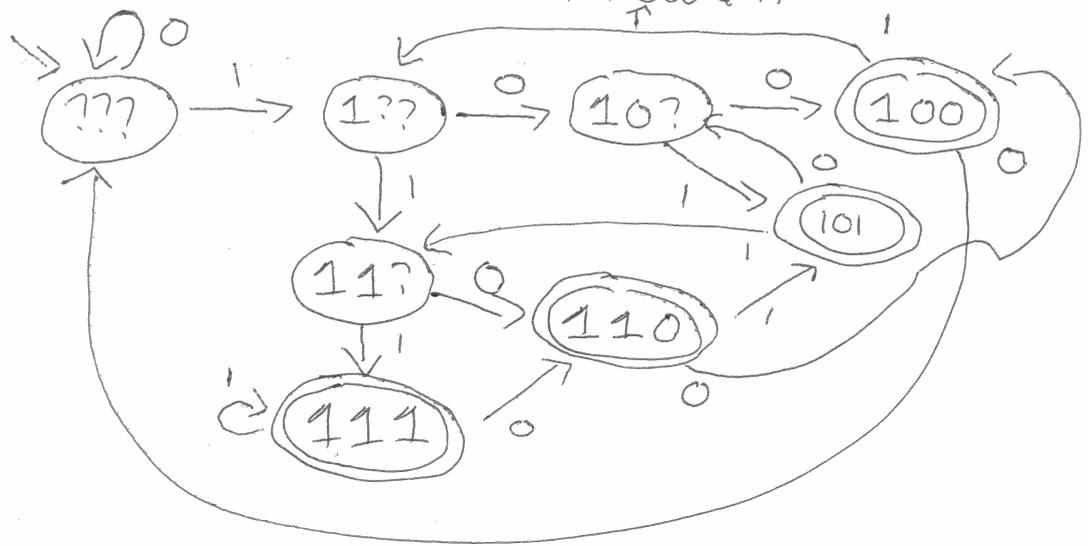
4-1/

Even machine (number is even, last char = 0)



3rd to last is 1 = A    100 ∈ A    0111 ∈ A

111000 ∉ A



~~enum~~ enum states { A, B };

state st = A;

while (chrc = getch()); {

  switch (st) {

    case A: switch (c) {

      case 0: st = B; bri;

      case 1: st = A; bri;

    }

  }

  case B: switch (c) {

    case 0: st = B; bri;

    case 1: st = A; bri; } }

return (st == B);

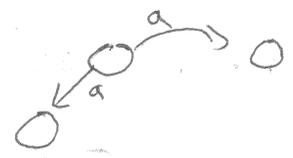
UML statecharts

1-2/

# NFA - non-deterministic finite automata

differences: 1) every input doesn't need a next state

2) there may be multiple transitions per input char



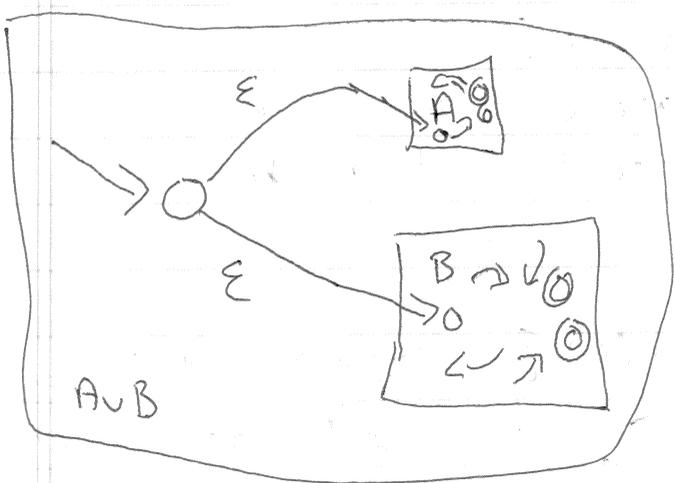
3) you may transition w/o reading



interpretations:

- 1) the NFA sees the future (rest of input) and chooses what works
- 2) the NFA Fork()s and does both
- 3) the NFA tries the first, if it doesn't work, go back and try second (a long time) — backtracking
- 4) math & cleverness

$$L(\rightarrow \circ \overset{1}{\downarrow} \rightarrow \circ \rightarrow \circ \rightarrow \odot) = (\text{3rd to last is } 1)$$



4-3/

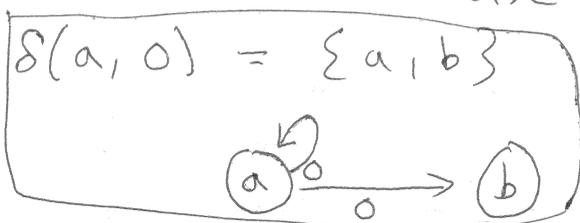
DFA  $d = (Q, \Sigma, q_0, \delta: Q \times \Sigma \rightarrow Q, F \subseteq Q)$

NFA  $n = (Q, \Sigma, q_0, \delta: Q \times \Sigma_\epsilon \rightarrow P(Q), F \subseteq Q)$

$$\Sigma_\epsilon = \Sigma \cup \{\epsilon\}$$

(input alphabet ~~plus~~ plus epsilon)

$Q = \{a, b, c\}$   $P(Q) = \{\emptyset, \{a\}, \{b\}, \{c\}, \{a, b\}, \{a, c\}, \{b, c\}, \{a, b, c\}\}$



DFA  $\delta: Q \times \Sigma \rightarrow Q$

NFA  $\delta: Q \times \Sigma_\epsilon \rightarrow P(Q)$

$L(\text{NFA } n) = \text{some set} = \{x \in \Sigma^* \mid \text{such that } x \text{ is accepted by } n\}$

A string  $x$  is accepted by NFA  $n$  iff

$$q_0 \xrightarrow{x}^* q_f \quad \text{sit. } q_f \in F$$

An NFA  $n$  transitions (or runs) from  $q_i$  to  $q_k$  via string

$x$  ( $q_i \xrightarrow{x}^* q_k$ ) iff

$$q_i \xrightarrow{\epsilon}^* q_i$$

$$q_i \xrightarrow{a} q_l$$

$$a \in \Sigma_\epsilon$$

and  $q_l \xrightarrow{y}^* q_k$

$$y \in \Sigma^*$$

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$$q_i \xrightarrow{ay}^* q_k$$

An NFA  $n$  steps from  $q_i$  to  $q_l$  on  $a$  iff.

$$q_l \in \delta(q_i, a)$$

