

$$A_{CFG} = \{ \langle G, w \rangle \mid G \in CFG, w \in \Sigma^*, \text{ and } G \text{ derives } w \}$$

$\langle S \rightarrow OS1 \mid \epsilon, 0011 \rangle \in A_{CFG}$

- 1 Starting from S
- 2 Non-deterministically apply rules
- 3 if you generate w, then accept

Tape 1: G
 Tape 2: w = 0011
 Tape 3: S → OS1 →
00S11 → ...11

definitely in Σ_1 (accepter)
 not in Σ_0 (decider) because it keeps trying

idea 1: can't terminals in tape 3, if $> |w|$, reject (stop on this path)
 breaks on: $S \rightarrow SS$

idea 2: assume / translate G into Chomsky Normal Form
 $S \rightarrow \epsilon \quad A \rightarrow BC \quad (B \& C \neq S) \quad A \rightarrow +$

there's never more than $2^{|w|}$ rule applications to generate |w|
 $\in \Sigma_0$

$$REG \subset \text{Deterministic CFL} \subset CFL$$

$$P \subset \Sigma_0 \quad EXP \subset \Sigma_0$$

$$NP \subset \Sigma_0$$

$$REG \subset \Sigma_0 \text{ (compiler)}$$

$$(CFL \subset \Sigma_0) \Leftrightarrow$$

$$(\forall G \in CFG, \exists T \in \Sigma_0)$$

Pf. $T = 0^n$ input w, run $A_{CFG} \langle G, w \rangle$

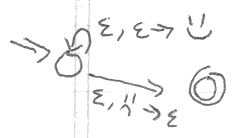
E_{CFG} = is a CFG's language empty?

idea: check which vars go to terminals, back-propagate this

info, then check S

(n^2 in number of rules)

(vars going terminals, = (\emptyset, V)
 other vars)



24-2/

$$A_{TM} = \{ \langle T, w \rangle \mid T \in TM \text{ and } w \in \Sigma^* \text{ and } T \text{ accepts } w \}$$

A_{TM} is solved by U

$U =$ "On input $\langle T, w \rangle$, Interpret
simulate T on w "

$A_{TM} \in \Sigma_1$ ✓

$A_{TM} \in \Sigma_0$? X $\langle T, w \rangle$ where $w \in L(T)$ it rejects
 U doesn't decide

PROMISE: Prove that it's not

(Σ_0) SQL, Datalog, 1d linker-script, Linux/Cisco firewall rules,
regexps (but not Perl), Coq, Agda, ML module linker

Georg Cantor

N (naturals = $0, 1, \dots, 42, \dots, 1978, \dots$)
are infinite

E (evens = $0, 2, 4, \dots$)
are infinite

Z (integers = $\dots, -1, 0, 1, \dots$)
infinite

Q (rationals = $0/1, 0/2, -1/3, 1/4, \dots$)
infinite

R (reals = $0, 1, 2, \pi, e, \sqrt{2}, \dots$)

R_{∞} (extended reals) = $R \cup \{-\infty, +\infty\}$

4-3
Two sets are the same size ($|X| = |Y|$)
if we can turn elements of
one into the other and back
and not lose information.
(same size (X, Y))

\exists a mapping function, f with 2 props:
one-to-one = $\forall a' \in A, b \in B,$
 $f(a) = b \rightarrow$
 $f(a') = b \rightarrow$
 $a = a',$

onto = $\forall b \in B, \exists a \in A, f(a) = b$

N is the same size as E
 $f(n) = 2 \cdot n$

1-to-1: Given a, a' , and b
 $2 \cdot a = b$
 $2 \cdot a' = b$
 $\rightarrow a = a' \quad \checkmark \quad \checkmark$

onto: $\forall b, \exists a, f(a) = b$
 $b = 2 \cdot n$
find a , s.t. $2 \cdot a = 2 \cdot n$
 $a = n$

A set X is countable if X is the same size as N
 $|X| = \aleph_0$

Next ... Z ? Q ? R ?

Foreshadowing m m m

Is TM the same size as ALL?

