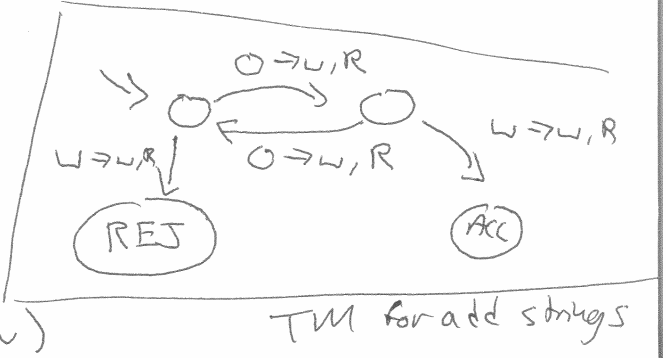


REG \subset ~~TM~~ Σ_0
 \downarrow \downarrow
 a language that has a DFA a language that has a decider



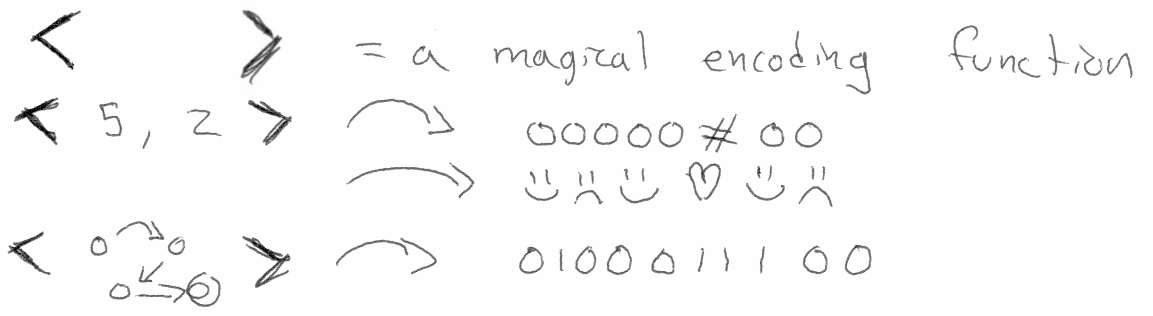
convert a DFA into a TM decider
 1. move right (write blank on whatever)
 2. on blank, look at F
 and go to ACC/REJ



$$\delta(q_i, w) = (ACC, w, R) \text{ if } q_i \in F$$

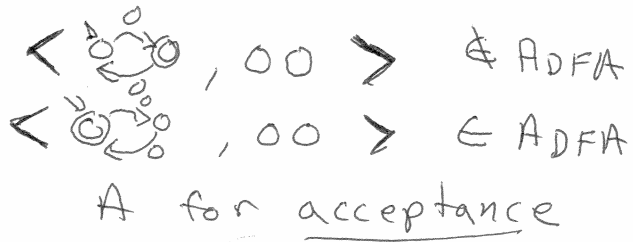
$$\delta(q_i, w) = (REJ, w, R) \text{ if } q_i \notin F$$

Cooler: 1 TM that is all REGs (interpreter)



$$A_{DFA} = \{ \langle D, w \rangle \mid D \in DFA, w \in D, \Sigma^* \}$$

D accepts w



- $\langle \text{DFA}, 000 \rangle \in A_{DFA}$
1. Decode input
 2. Look at first char in input tape (ie w)
 3. Consult δ and change the state tape
 4. If the char is w, consult F

Tapes: D, w, q_i
 \uparrow
 state tape

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$$A_{NFA} = \{ \langle N, w \rangle \mid N \in NFA, w \in \Sigma^*, N \text{ accepts } w \}$$

Option 1: use non-det TTM

Option 2: Compile N to D, then use ADFA
TM transducer

$$A_{REX} = \{ \langle R, w \rangle \mid R \in REX, w \in \Sigma^*, R \text{ accepts } w \}$$

option: use compiler to NFA, then use ADFA

ADFA, NFA, REX are all decidable ($\in \Sigma_0$)

$X \in \Sigma_0$ (compiler)

$A_X \in \Sigma_0$ (interpreter exists, is correct, and a decider)

$$E_{DFA} = \{ \langle D \rangle \mid D \in DFA \text{ and } L(D) = \emptyset \}$$

- Emptiness
1. If F is empty
 2. Look at stuff in F

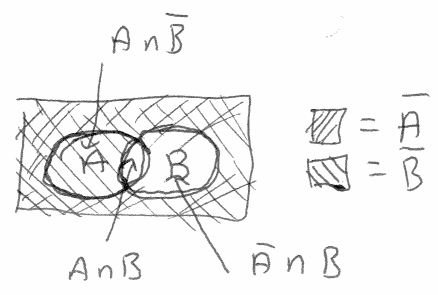


3. Is there a path from q_0 to $q_f \in F$?
Do a DFS on the graph (Q, δ)
from q_0

$$E_{EQ_{DFA}} = \{ \langle A, B \rangle \mid A \in DFA, B \in DFA, L(A) \neq L(B) \}$$

$$\left(\begin{array}{l} (A \cap \bar{B}) \\ (\bar{A} \cap B) \end{array} \right) \cup C = \emptyset \quad L(C) = \emptyset$$

$E_{DFA}(C)$



- Model-checking
- A = Program / Model
 - B = Specification