

File Systems (1)

Essential requirements for long-term information storage:

- It must be possible to store a very large amount of information.
- The information must survive the termination of the process using it.
- Multiple processes must be able to access the information concurrently.

File Systems (2)

Think of a disk as a linear sequence of fixed-size blocks and supporting reading and writing of blocks. Questions that quickly arise:

- How do you find information?
- How do you keep one user from reading another's data?
- How do you know which blocks are free?

File Types

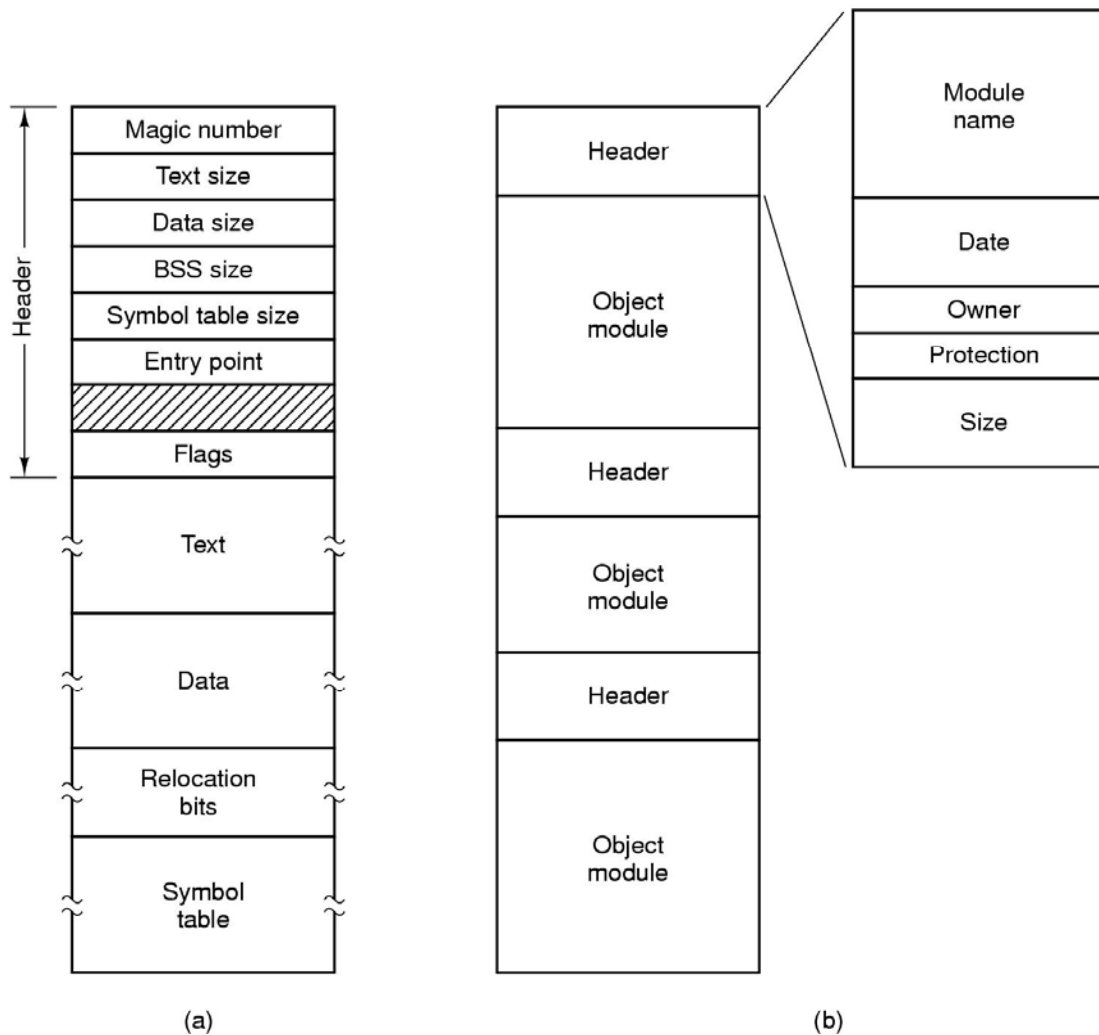


Figure 4-3. (a) An executable file. (b) An archive.

File Attributes

Attribute	Meaning
Protection	Who can access the file and in what way
Password	Password needed to access the file
Creator	ID of the person who created the file
Owner	Current owner
Read-only flag	0 for read/write; 1 for read only
Hidden flag	0 for normal; 1 for do not display in listings
System flag	0 for normal files; 1 for system file
Archive flag	0 for has been backed up; 1 for needs to be backed up
ASCII/binary flag	0 for ASCII file; 1 for binary file
Random access flag	0 for sequential access only; 1 for random access
Temporary flag	0 for normal; 1 for delete file on process exit
Lock flags	0 for unlocked; nonzero for locked
Record length	Number of bytes in a record
Key position	Offset of the key within each record
Key length	Number of bytes in the key field
Creation time	Date and time the file was created
Time of last access	Date and time the file was last accessed
Time of last change	Date and time the file was last changed
Current size	Number of bytes in the file
Maximum size	Number of bytes the file may grow to

Figure 4-4a. Some possible file attributes.

File Operations

The most common system calls relating to files:

- Create
- Delete
- Open
- Close
- Read
- Write
- Append
- Seek
- Get Attributes
- Set Attributes
- Rename

Hierarchical Directory Systems (1)

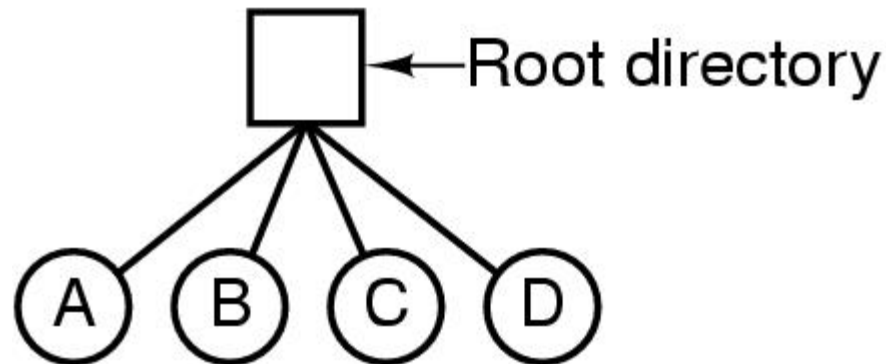


Figure 4-6. A single-level directory system containing four files.

Hierarchical Directory Systems (2)

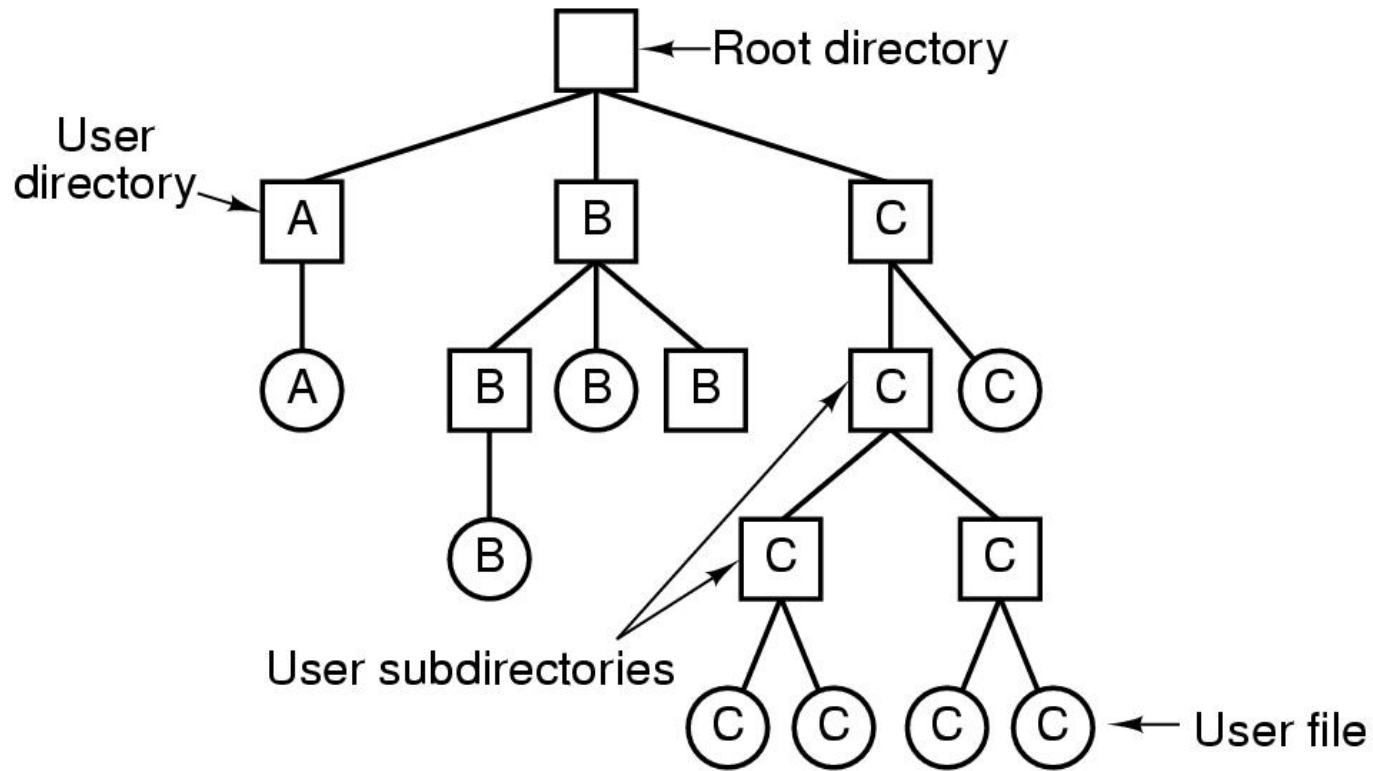


Figure 4-7. A hierarchical directory system.

Path Names

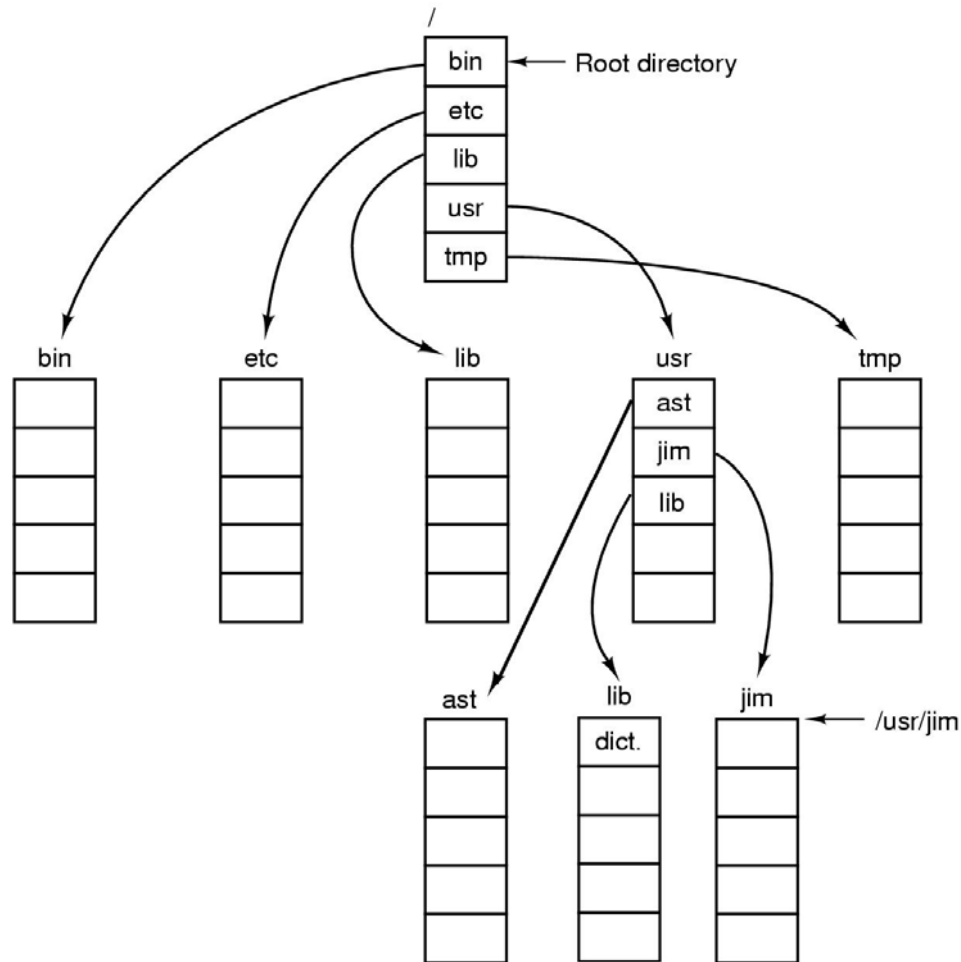


Figure 4-8. A UNIX directory tree.

Directory Operations

System calls for managing directories:

- Create
- Delete
- Opendir
- Closedir
- Readdir
- Rename
- Link
- Unlink

File System Layout

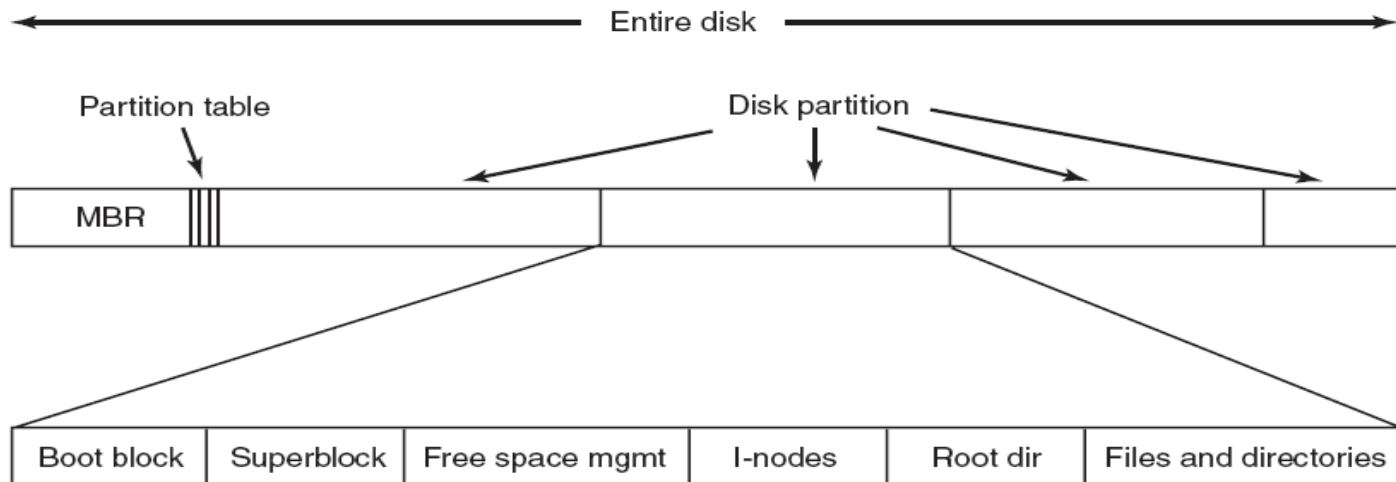


Figure 4-9. A possible file system layout.

Contiguous Allocation

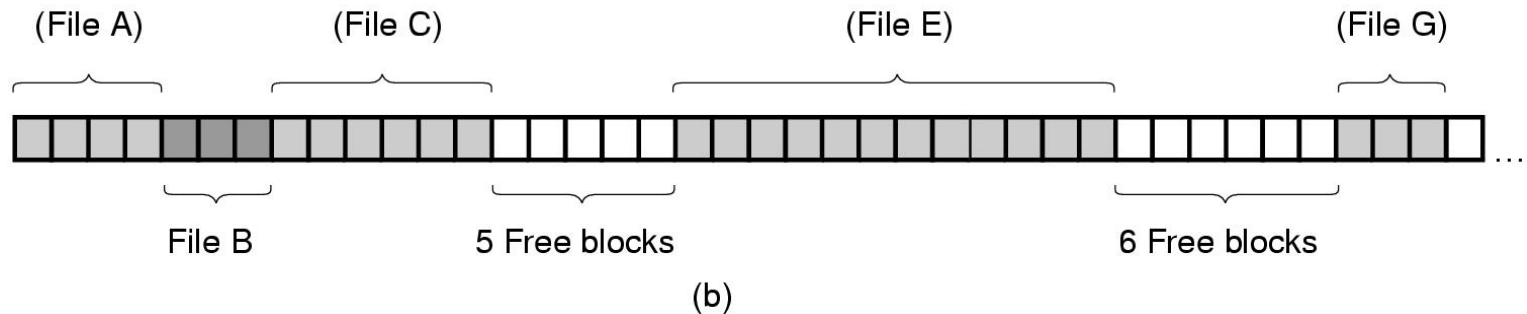
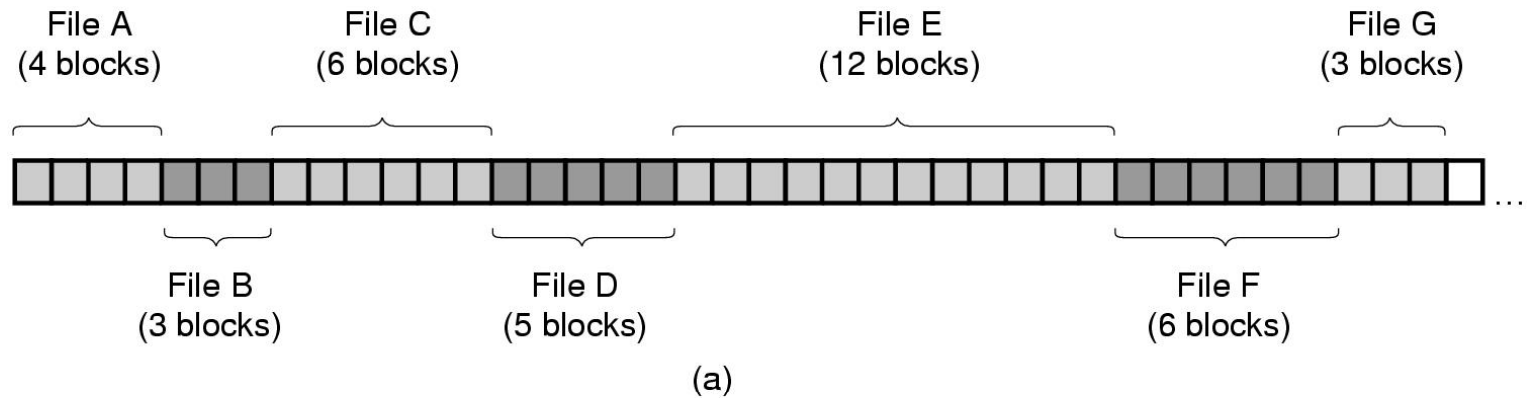


Figure 4-10. (a) Contiguous allocation of disk space for 7 files. (b) The state of the disk after files D and F have been removed.

Linked List Allocation

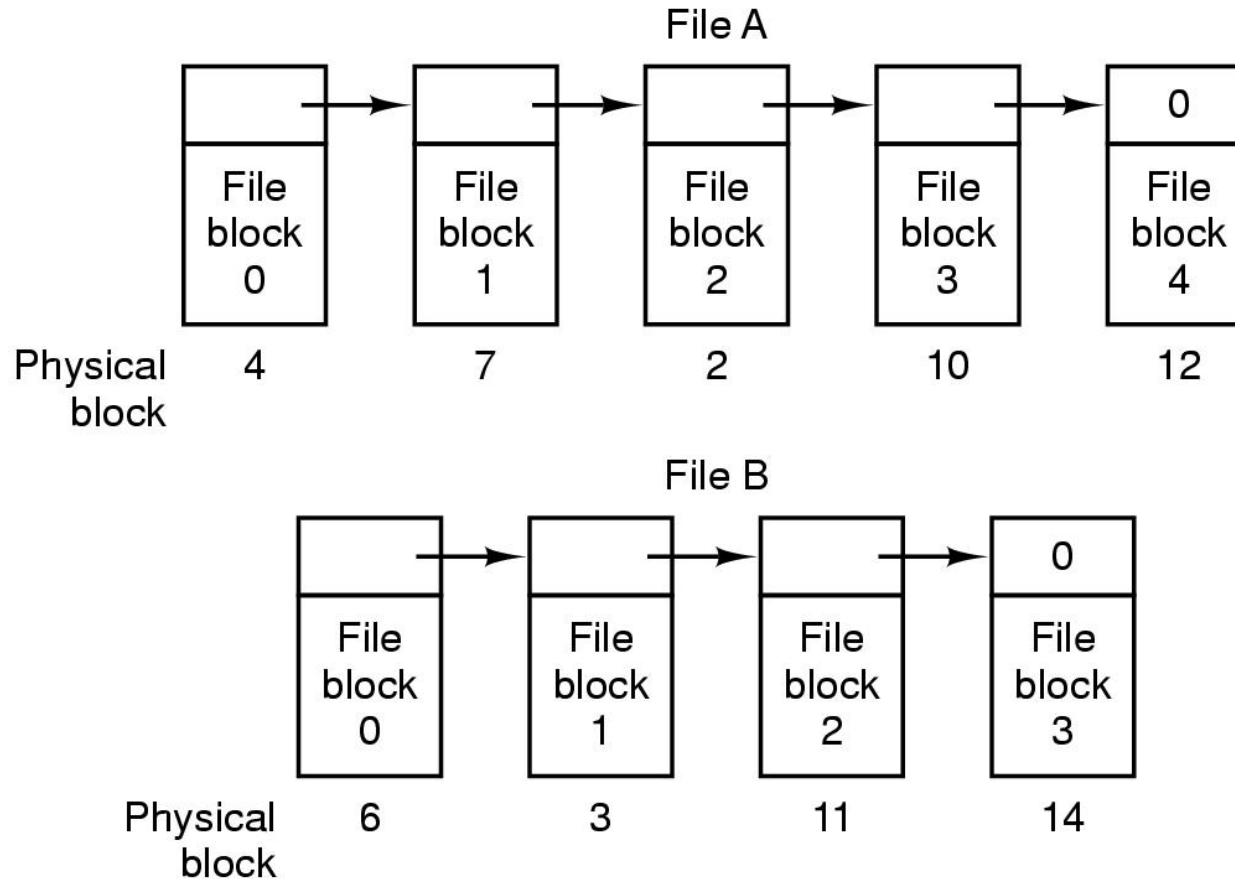


Figure 4-11. Storing a file as a linked list of disk blocks.

Linked List Allocation Using a Table in Memory

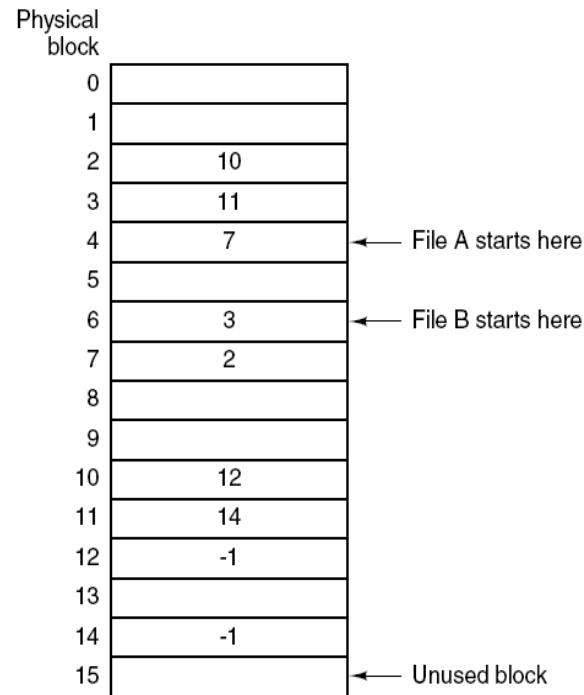


Figure 4-12. Linked list allocation using a file allocation table in main memory.

I-nodes

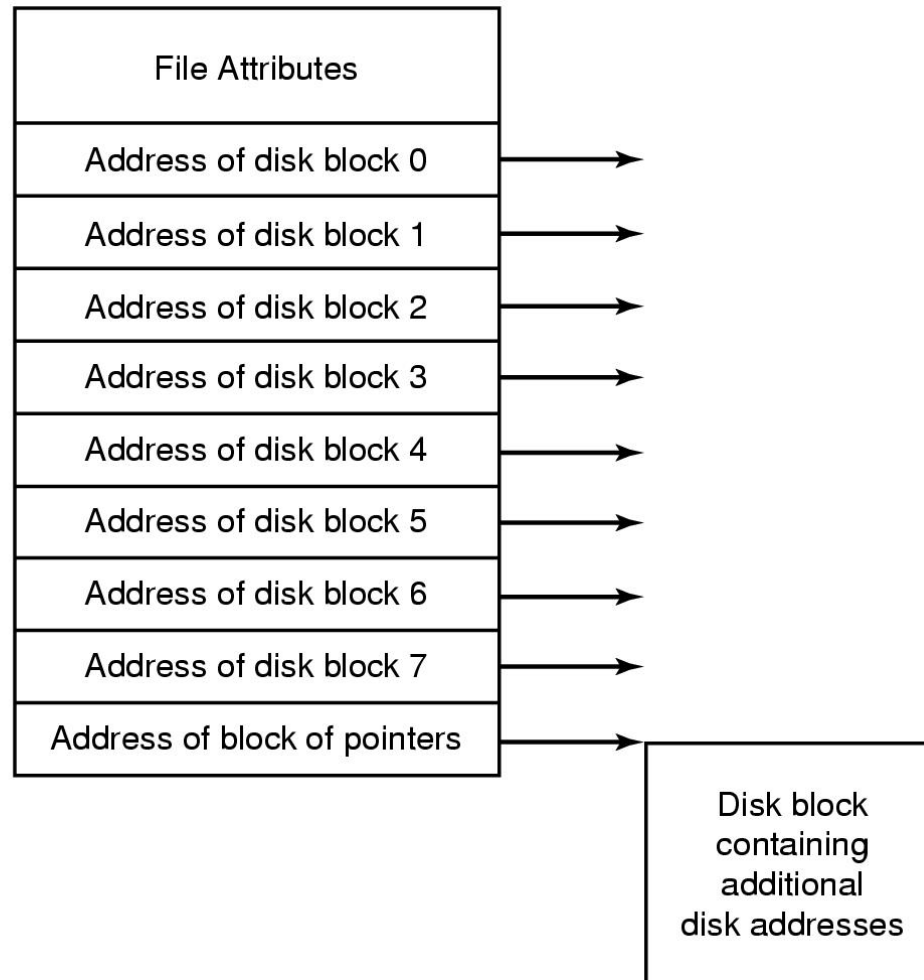


Figure 4-13. An example i-node.

Implementing Directories (1)

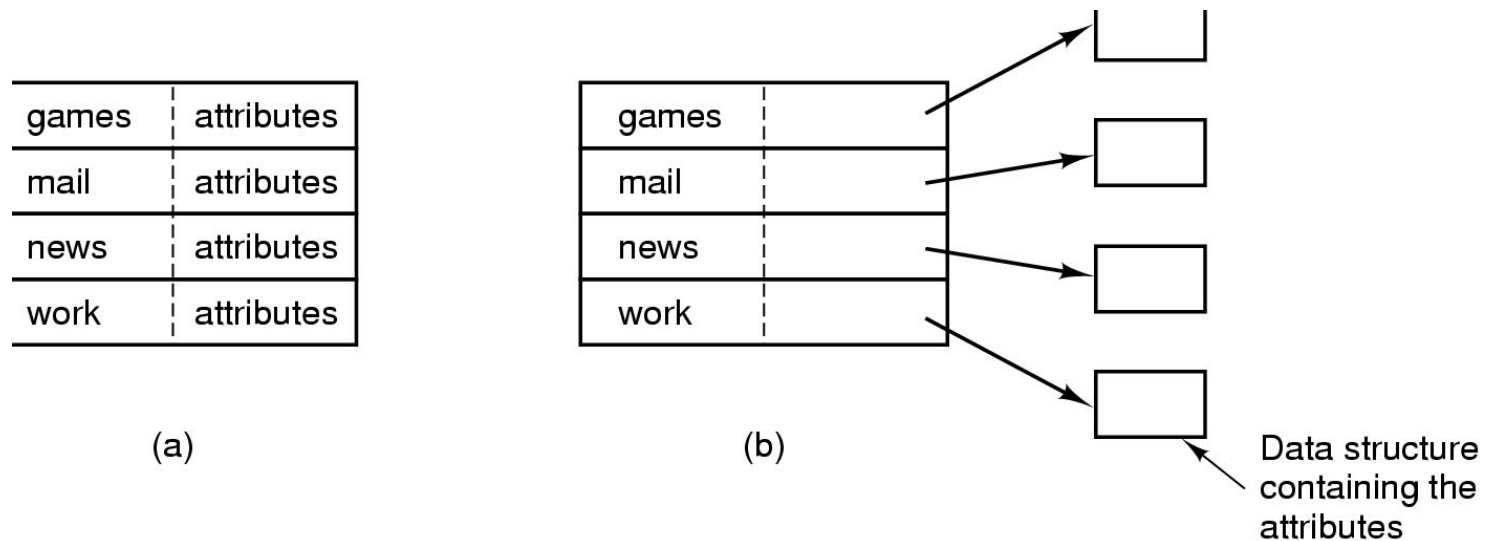


Figure 4-14. (a) A simple directory containing fixed-size entries with the disk addresses and attributes in the directory entry. (b) A directory in which each entry just refers to an i-node.

Implementing Directories (2)

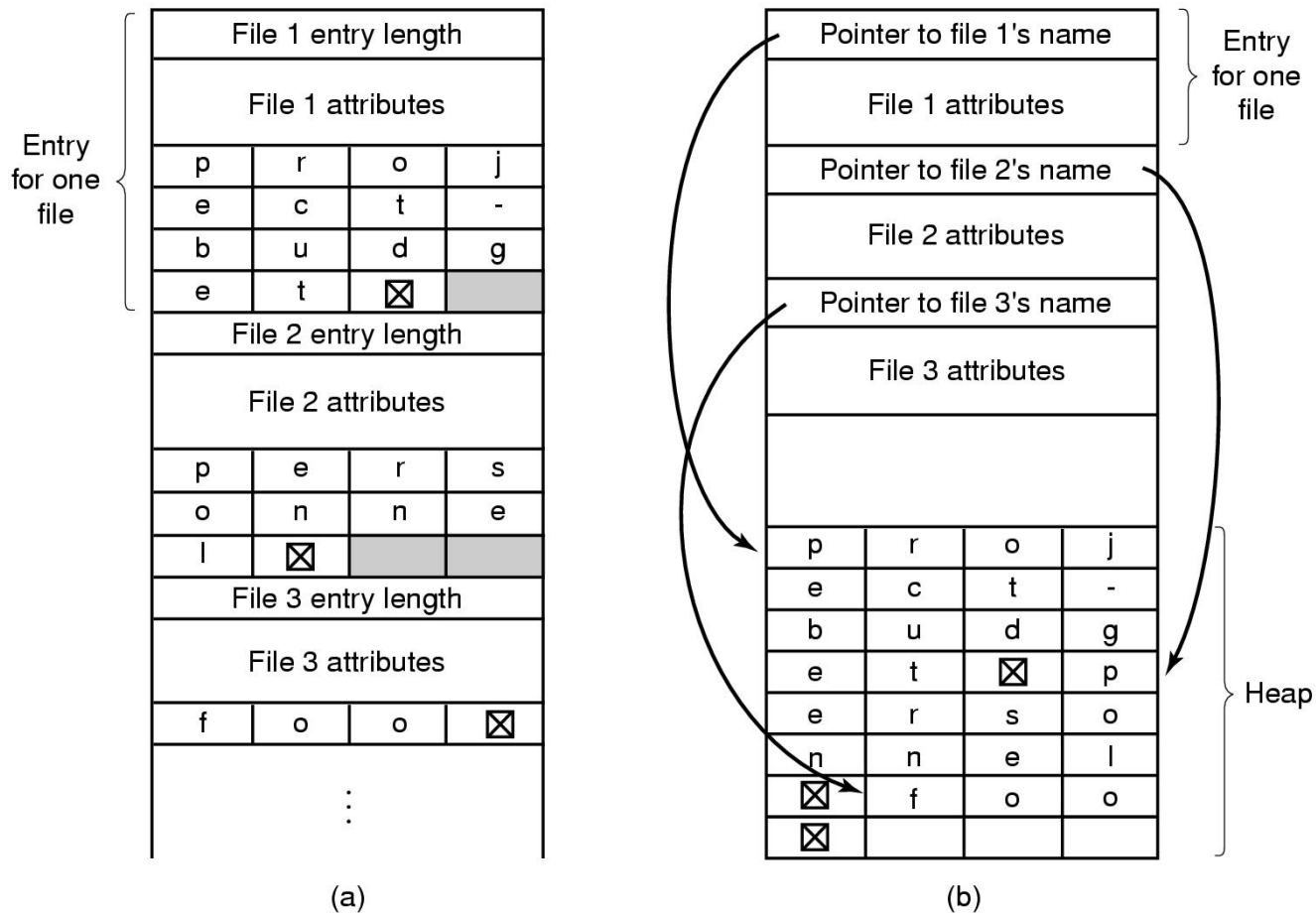


Figure 4-15. Two ways of handling long file names in a directory.
 (a) In-line. (b) In a heap.

Shared Files (1)

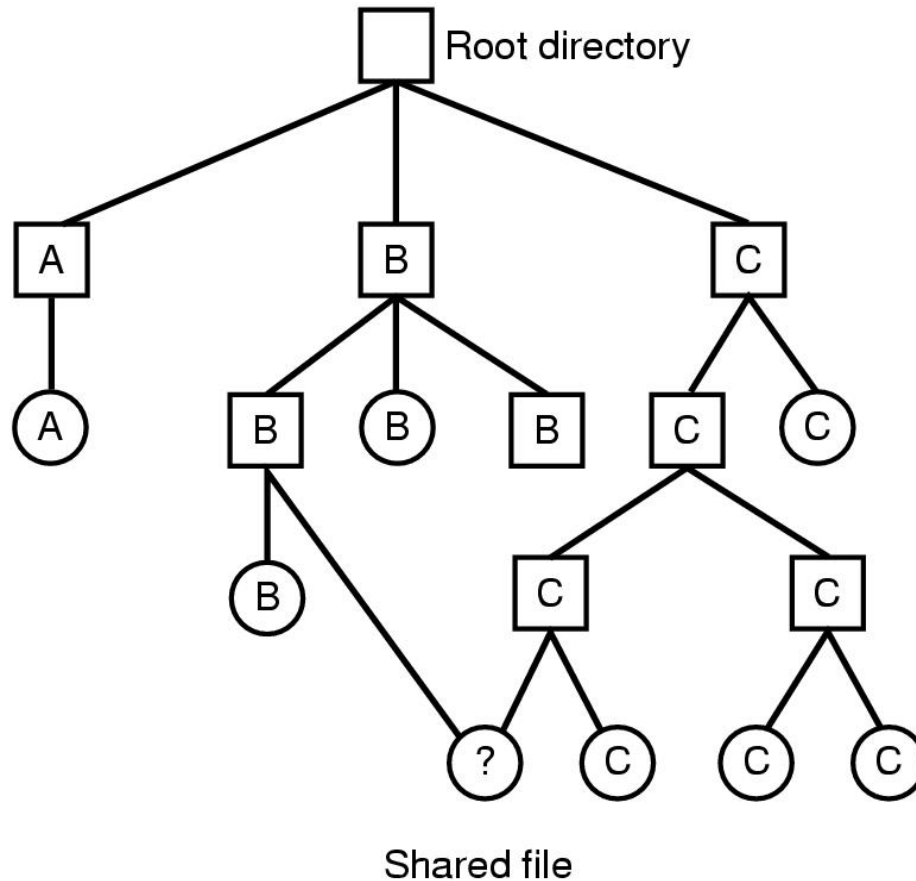


Figure 4-16. File system containing a shared file.

Shared Files (2)

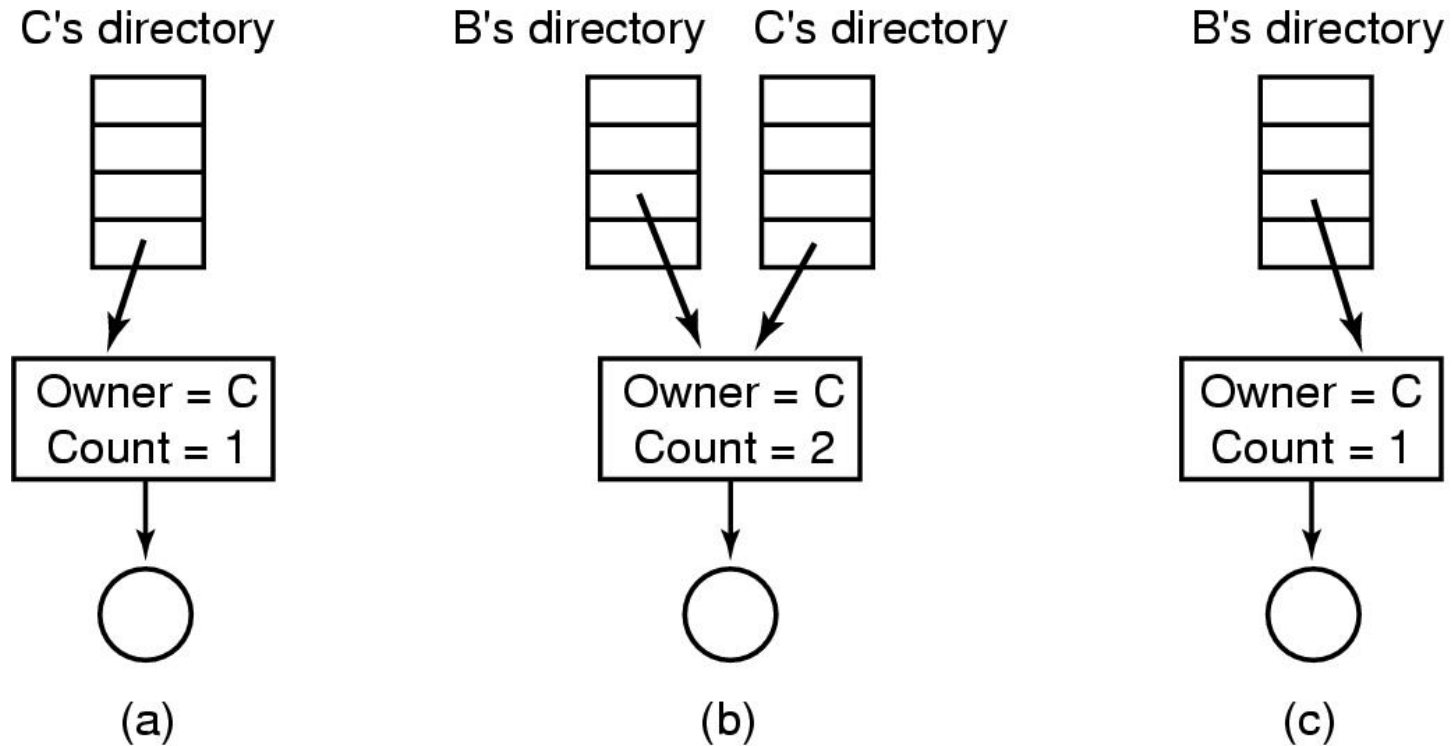


Figure 4-17. (a) Situation prior to linking. (b) After the link is created. (c) After the original owner removes the file.

Journaling File Systems

Operations required to remove a file in UNIX:

- Remove the file from its directory.
- Release the i-node to the pool of free i-nodes.
- Return all the disk blocks to the pool of free disk blocks.

Virtual File Systems (1)

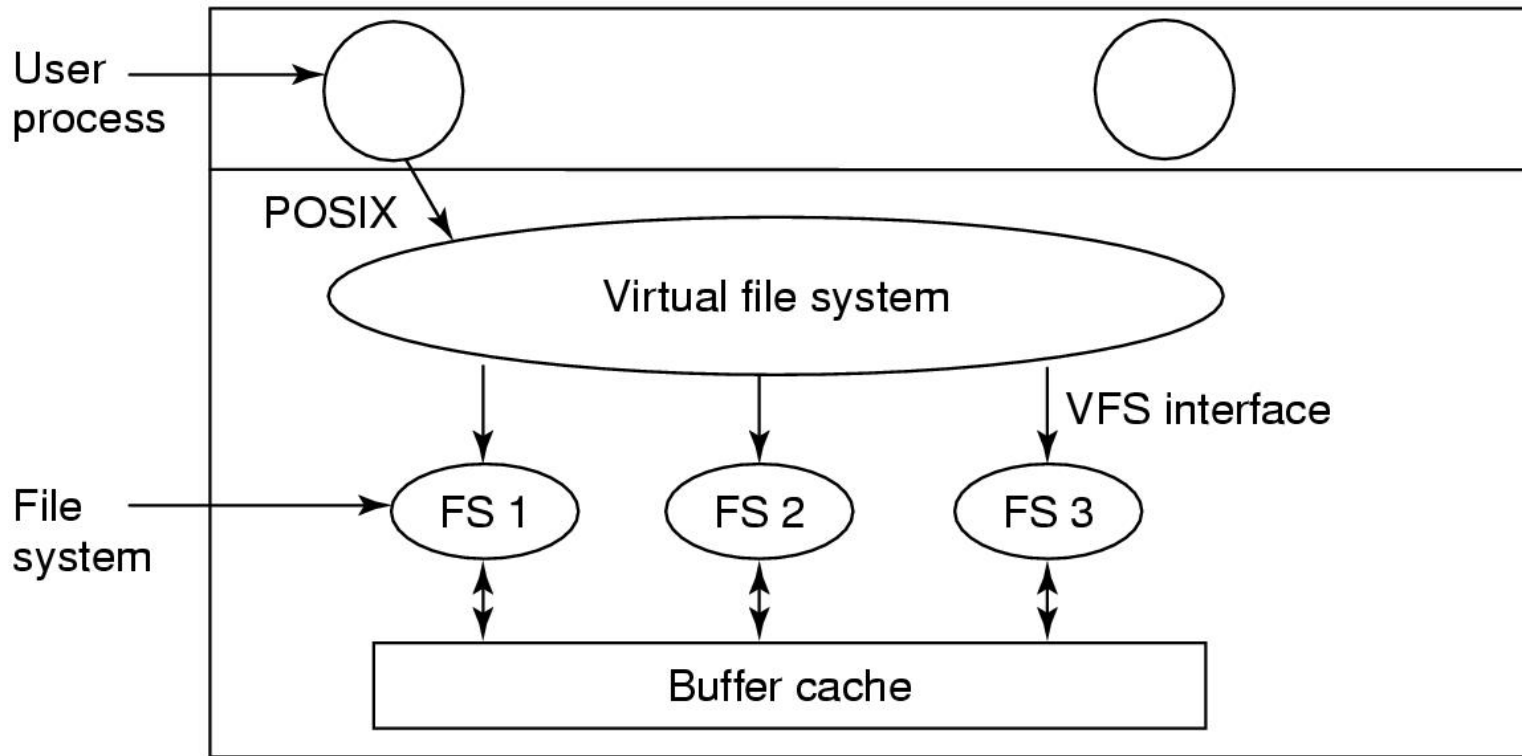


Figure 4-18. Position of the virtual file system.

Virtual File Systems (2)

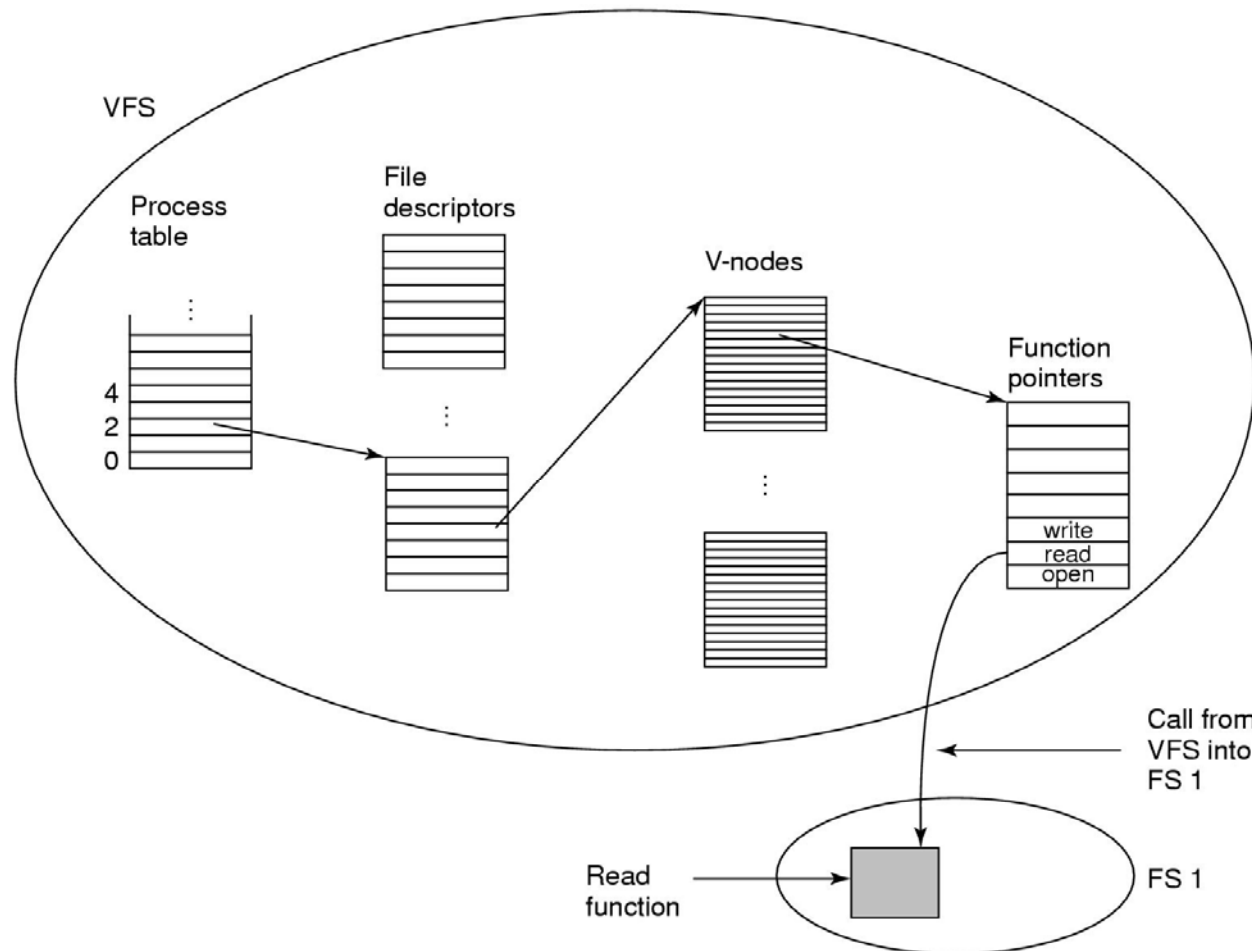


Figure 4-19. A simplified view of the data structures and code used by the VFS and concrete file system to do a read.

Disk Space Management Block Size (1)

Length	VU 1984	VU 2005	Web
1	1.79	1.38	6.67
2	1.88	1.53	7.67
4	2.01	1.65	8.33
8	2.31	1.80	11.30
16	3.32	2.15	11.46
32	5.13	3.15	12.33
64	8.71	4.98	26.10
128	14.73	8.03	28.49
256	23.09	13.29	32.10
512	34.44	20.62	39.94
1 KB	48.05	30.91	47.82
2 KB	60.87	46.09	59.44
4 KB	75.31	59.13	70.64
8 KB	84.97	69.96	79.69

Length	VU 1984	VU 2005	Web
16 KB	92.53	78.92	86.79
32 KB	97.21	85.87	91.65
64 KB	99.18	90.84	94.80
128 KB	99.84	93.73	96.93
256 KB	99.96	96.12	98.48
512 KB	100.00	97.73	98.99
1 MB	100.00	98.87	99.62
2 MB	100.00	99.44	99.80
4 MB	100.00	99.71	99.87
8 MB	100.00	99.86	99.94
16 MB	100.00	99.94	99.97
32 MB	100.00	99.97	99.99
64 MB	100.00	99.99	99.99
128 MB	100.00	99.99	100.00

Figure 4-20. Percentage of files smaller than a given size (in bytes).

Disk Space Management Block Size (2)

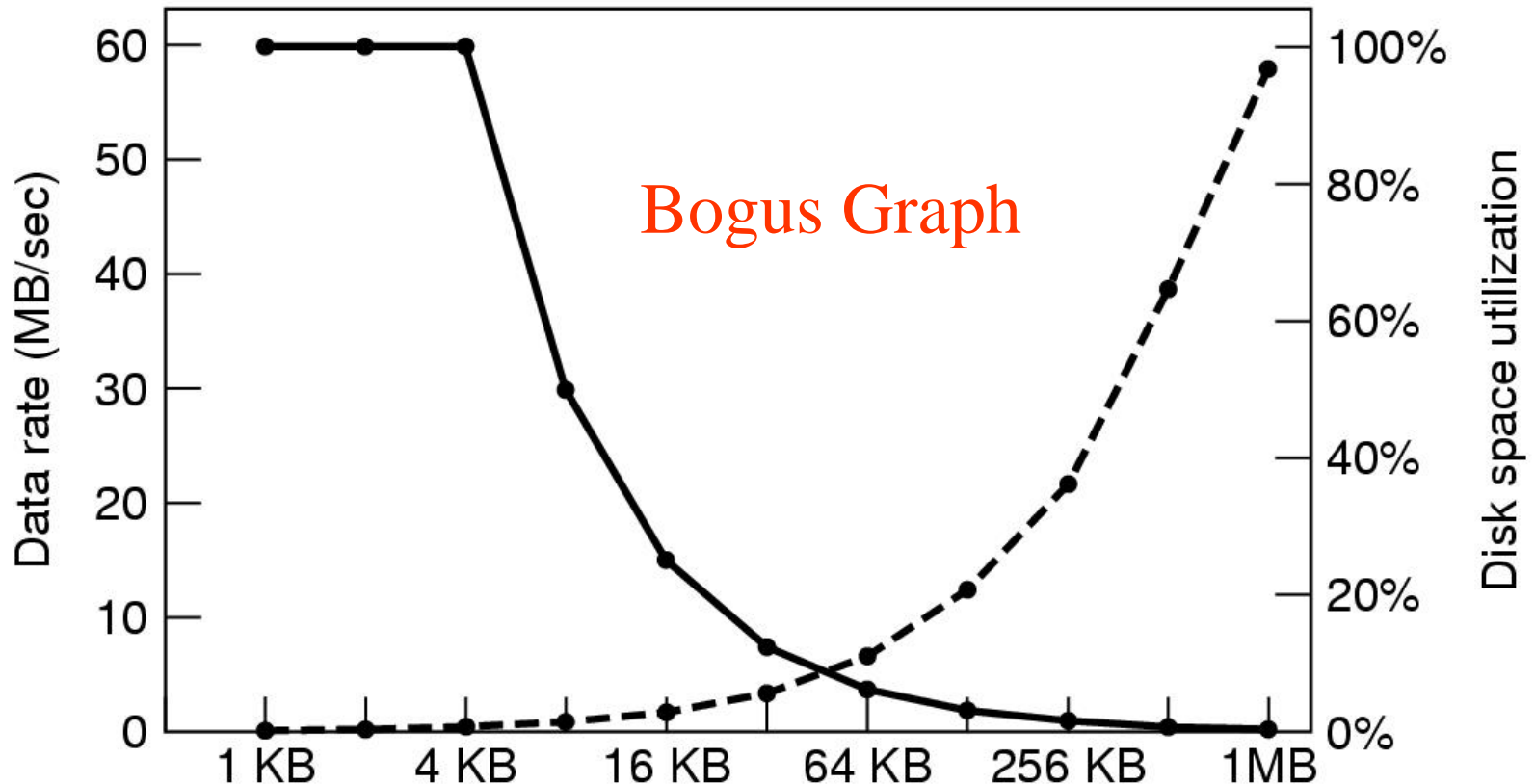


Figure 4-21. The solid curve (left-hand scale) gives the data rate of a disk. The dashed curve (right-hand scale) gives the disk space efficiency. All files are 4 KB.

Keeping Track of Free Blocks (1)

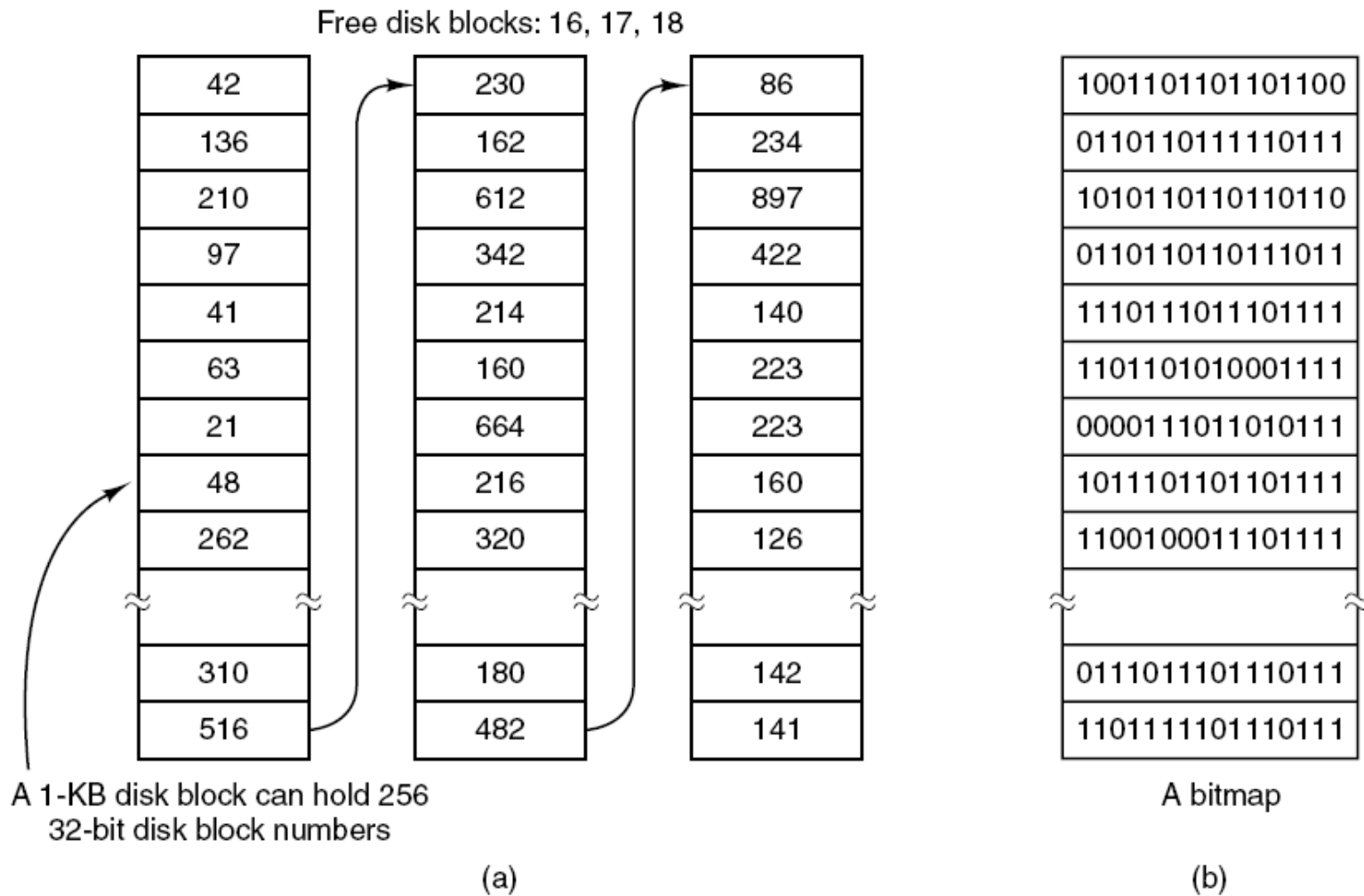


Figure 4-22. (a) Storing the free list on a linked list. (b) A bitmap.

Caching (1)

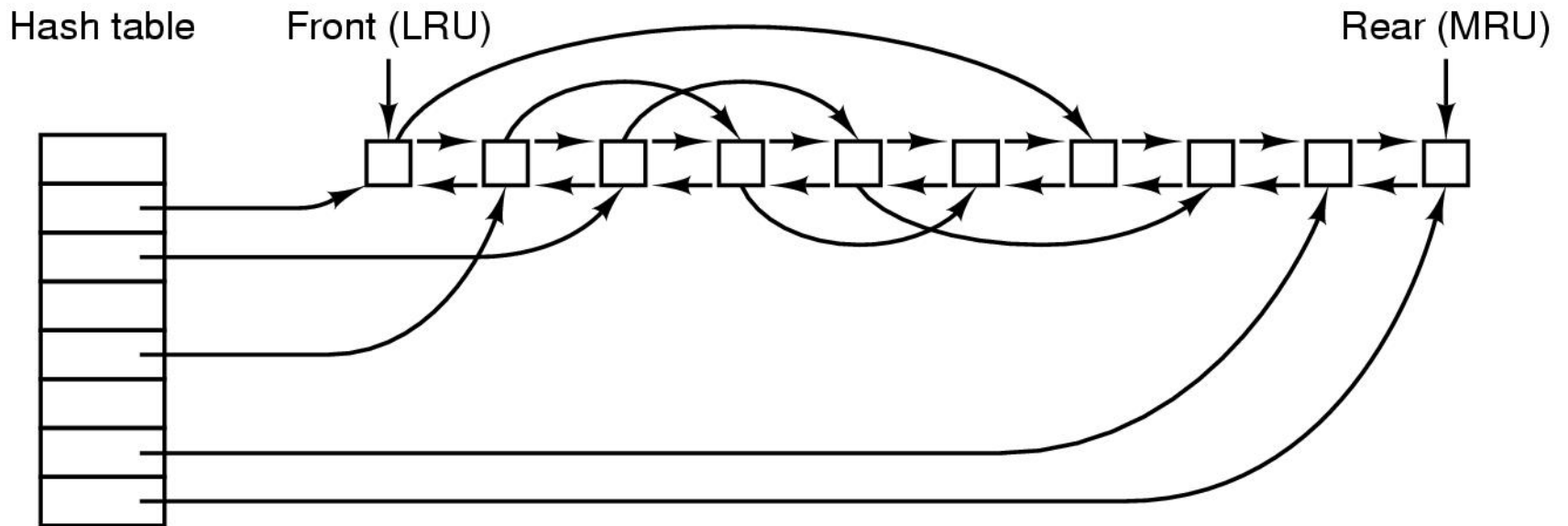


Figure 4-28. The buffer cache data structures.

Caching (2)

- Some blocks, such as i-node blocks, are rarely referenced two times within a short interval.
- Consider a modified LRU scheme, taking two factors into account:
 - Is the block likely to be needed again soon?
 - Is the block essential to the consistency of the file system?

Reducing Disk Arm Motion

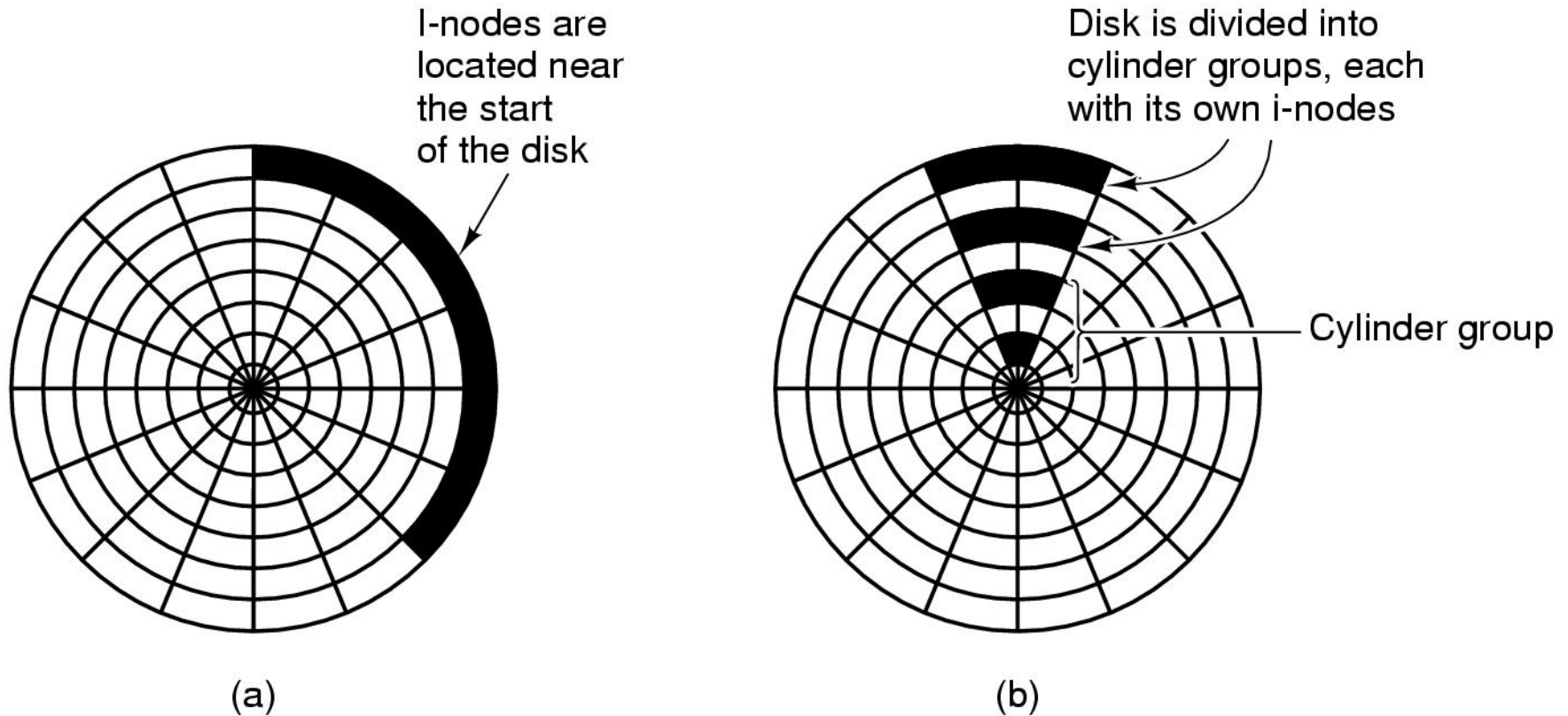


Figure 4-29. (a) I-nodes placed at the start of the disk. (b) Disk divided into cylinder groups, each with its own blocks and i-nodes.

The ISO 9660 File System

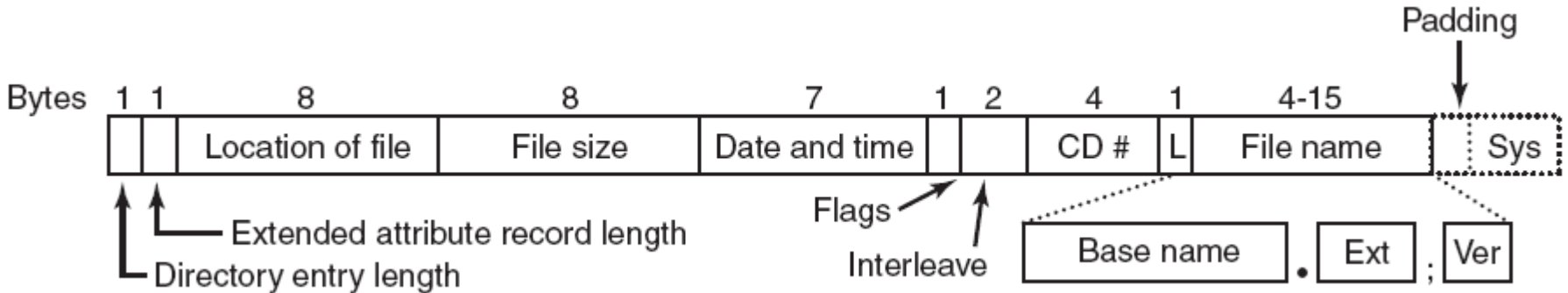


Figure 4-30. The ISO 9660 directory entry.

Rock Ridge Extensions

Rock Ridge extension fields:

- PX - POSIX attributes.
- PN - Major and minor device numbers.
- SL - Symbolic link.
- NM - Alternative name.
- CL - Child location.
- PL - Parent location.
- RE - Relocation.
- TF - Time stamps.

Joliet Extensions

Joliet extension fields:

- Long file names.
- Unicode character set.
- Directory nesting deeper than eight levels.
- Directory names with extensions

The MS-DOS File System (1)

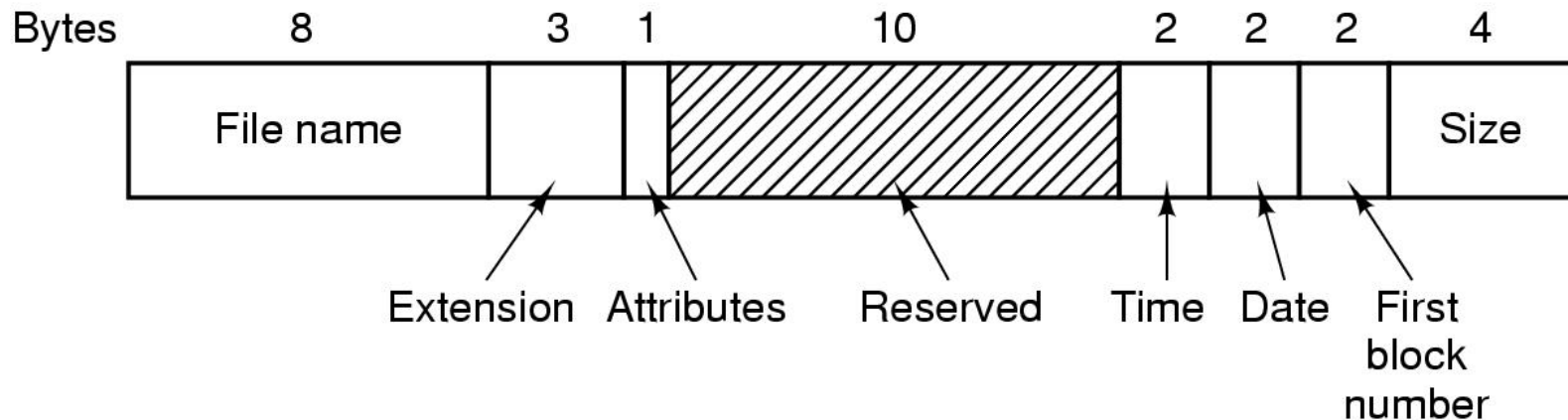


Figure 4-31. The MS-DOS directory entry.

The MS-DOS File System (2)

Block size	FAT-12	FAT-16	FAT-32
0.5 KB	2 MB		
1 KB	4 MB		
2 KB	8 MB	128 MB	
4 KB	16 MB	256 MB	1 TB
8 KB		512 MB	2 TB
16 KB		1024 MB	2 TB
32 KB		2048 MB	2 TB

Figure 4-32. Maximum partition size for different block sizes. The empty boxes represent forbidden combinations.

The UNIX V7 File System (1)

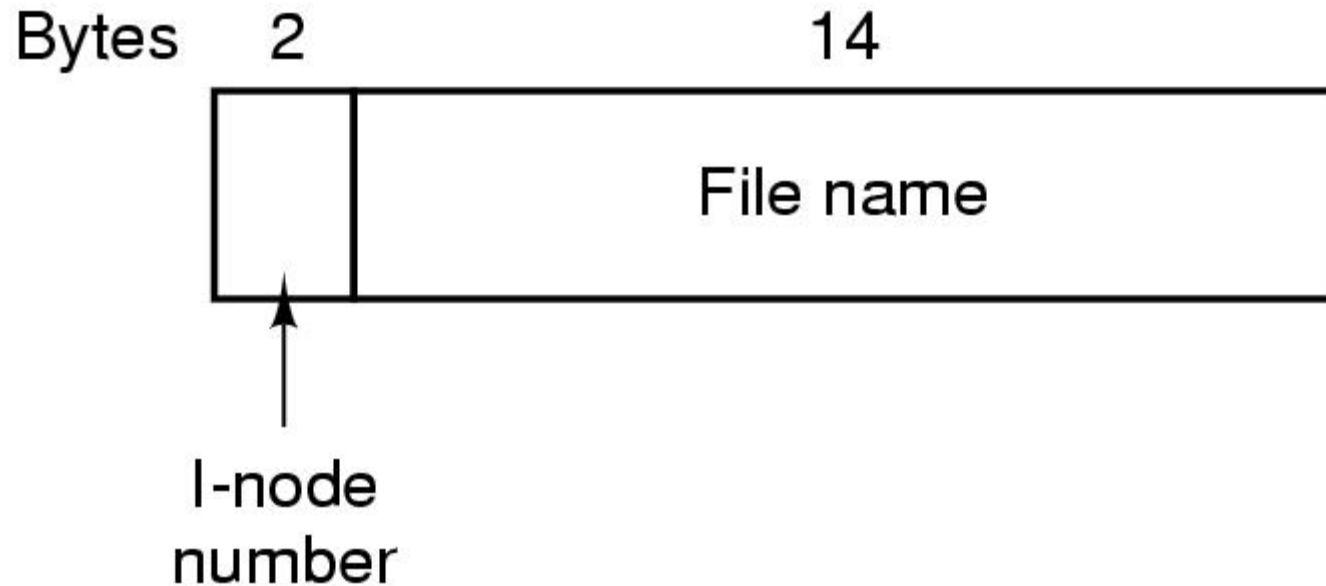


Figure 4-33. A UNIX V7 directory entry.

The UNIX V7 File System (2)

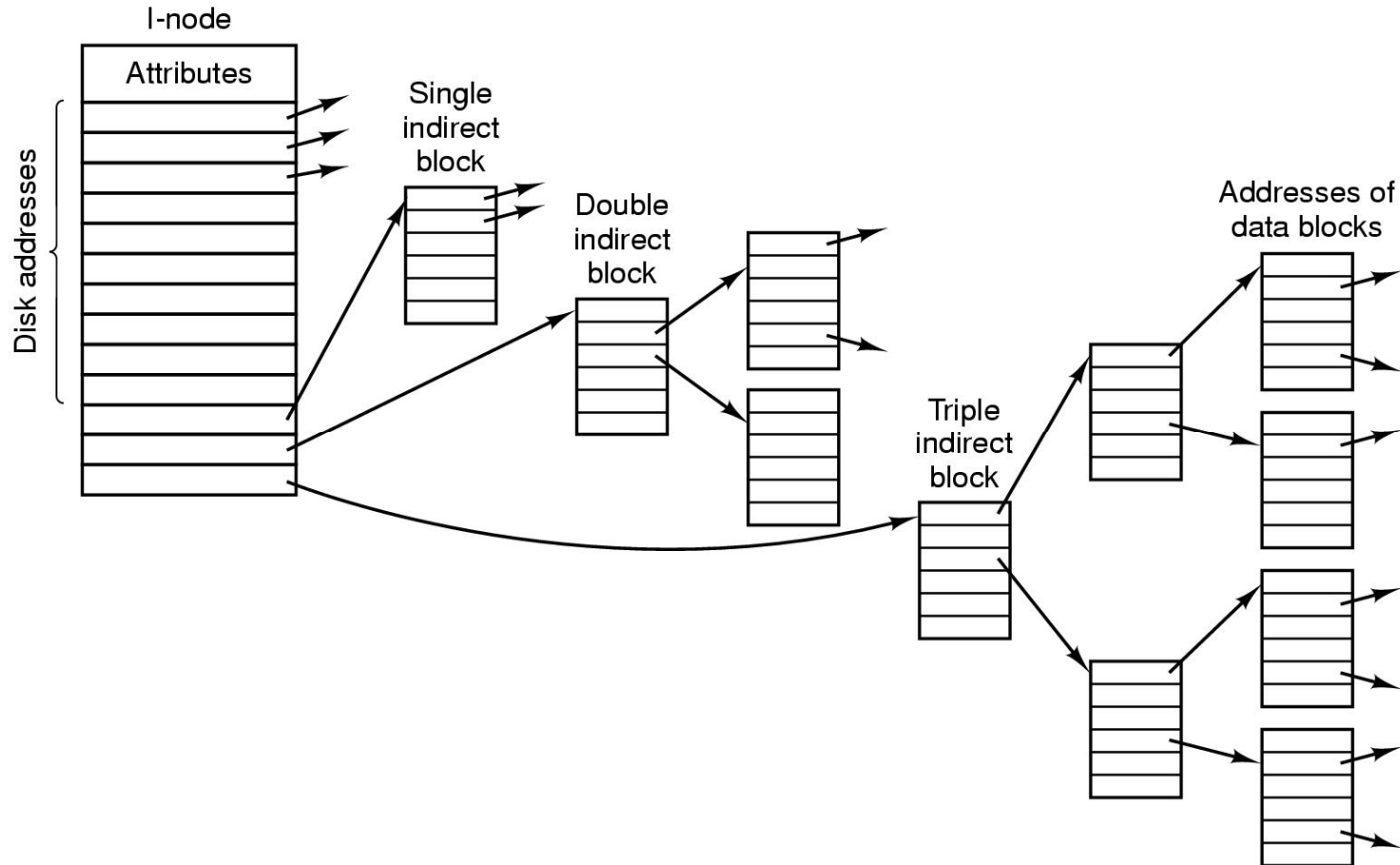


Figure 4-34. A UNIX i-node.

The UNIX V7 File System (3)

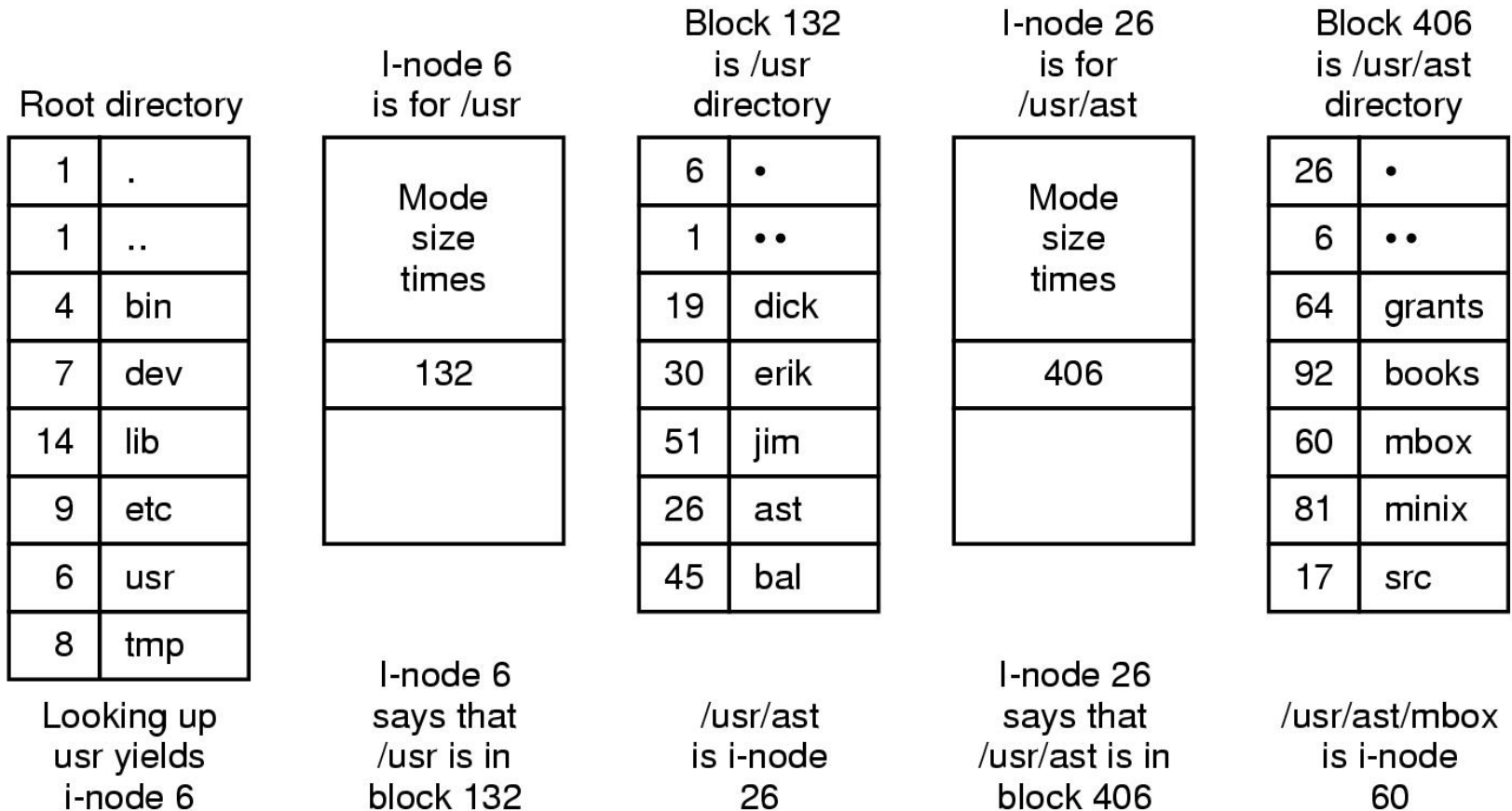


Figure 4-35. The steps in looking up `/usr/ast/mbox`.